CPSC 540: Machine Learning Exact Inference in Graphical Models

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Winter 2016

Admin

- Assignment 3:
 - Today is the final day to hand it in.
- Assignment 4:
 - Due on Tuesday.
 - Thursday is the last day to hand it in.
- Midterm:
 - March 17 in class.
 - Closed-book, two-page double-sided 'cheat sheet'.
 - Only covers topics from assignments A1-A4.
 - No requirement to pass.
 - Midterm from last year posted on Piazza.
 - Help session on March 16 from 3-5.
- Final Project:
 - Many of you are choosing project that are too big/hard.
 - In the project proposal, try to narrow down the scope:
 - Think of the final project as A6.
 - Main objective: show me you've learned something in this class, *and* explored a topic not covered in assignments.

DAGs vs. UGMs

General Graphs

Last Two Lectures: Directed and Undirected Graphical Models

• DAG models represent probability as ordered product of conditionals,

$$p(x) = \prod_{j=1}^d p(x_j | x_{\mathsf{pa}(j)}),$$

and are also known as "Bayesian networks" and "belief networks".

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and are also known as "Bayesian networks" and "belief networks".

• UGMs represent probability as product of non-negative potentials,

$$p(x) = \frac{1}{Z} \prod_{c \in \mathcal{C}} \phi_c(x_c),$$

and are also known as "Markov random fields" and "Markov networks".

• Models are useful fordensity estimation and structured prediction.

Markov Chains and Markov Property

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• A generalization of this property for general DAG models is:

$$p(x_j|x_{1:j-1}) = p(x_j|x_{\mathsf{pa}(j)}).$$

Markov Chains and Markov Property

• For chain-structured UGMs,

$$p(x) = \frac{1}{Z} \prod_{j=2}^{d} \phi_{j,j-1}(x_j, x_{j-1}),$$

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• But for UGMs we have the local Markov property,

$$p(x_j|x_{\{1:d\}\setminus j}) = p(x_j|x_{\mathsf{nei}(j)}),$$

where nei(j) are the neighbours in the graph (Markov blanket).

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• For chain-structured UGMs, we thus have

$$p(x_j | x_{\{1:d\}\setminus j}) = p(x_j | x_{j-1}, x_{j+1}),$$

that you're independent of the past/future given last/next time.

Markov Blanket

• Markov blanket in UGMs is all neighbours in the graphs:



• Markov blanket in DAGs is all parents, children, and co-parents:



Decoding, Inference, and Sampling in UGMs

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③ Sampling: Generate x according from the distribution:

$$x \sim p(x).$$

- All 3 are NP-hard in discrete UGMs.
- Even computing p(x) is NP-hard if we don't have Z.

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$$p(x_j = s) = \sum_{x_{\mathsf{pa}(j)}} p(x_j = s, x_{\mathsf{pa}(j)}) = \sum_{x_{\mathsf{pa}(j)}} \underbrace{p(x_j = s | x_{\mathsf{pa}(j)})}_{\text{given}} p(x_{\mathsf{pa}(j)}),$$

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and by independence of parents from unobserved children we have

$$p(x_{\mathsf{pa}}(j)) = \prod_{k \in \mathsf{pa}(j)} p(x_k | x_{\mathsf{pa}(k)}),$$

which is a product of marginals for k < j and conditionals that are given: (Sequentially compute $p(x_j = s)$ for each s from j = 1 to d.)

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• Ancestral sampling: If we want to sample from $p(x_1, x_2) = p(x_2|x_1)p(x_1)$,

Sample $x_1 \sim p(x_1)$, then sample $x_2 \sim p(x_2|x_1)$.

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• General DAGs: sample variables in order j = 1 to d, conditioning on earlier samples.

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- For example, still easy if condition on the first variables in the order:
 - Minor change to Chapman-Kolmogorov and ancestral sampling.



- NP-hard to condition on the last variables in the order:
 - Conditioning on descendent makes ancestors dependent.



Moralization: Converting DAGs to UGMs

- To address NP-hard problems, DAGs and UGMs use same methods.
- For DAGs, we typically just represent it as a UGM:

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• Graphically: we drop directions and "marry" parents (moralization).



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- This may lose some conditional independence information:
 - Models that be represented as DAGs or UGMs: "decomposable" and "triangulated".
 - Includes chain-structured and fully-connected graphs.

Empty Graphs

Chain-Structured Graphs

General Graphs



DAGs vs. UGMs

- 2 Empty Graphs
- Chain-Structured Graphs

④ General Graphs

Inference By Enumeration

• Last time, exact inference by table:

Cathy	Heather	Mark	Allison	np(1)	np(2)	np(3)	np(4)	ep(1)	ep(2)	ep(3)	prodPot	Probability
right	right	right	right	1	9	1	9	2	2	2	648	0.17
wrong	right	right	right	3	9	1	9	1	2	2	972	0.26
right	wrong	right	right	1	1	1	9	1	1	2	18	0.00
wrong	wrong	right	right	3	1	1	9	2	1	2	108	0.03
right	right	wrong	right	1	9	3	9	2	1	1	486	0.13
wrong	right	wrong	right	3	9	3	9	1	1	1	729	0.19
right	wrong	wrong	right	1	1	3	9	1	2	1	54	0.01
wrong	wrong	wrong	right	3	1	3	9	2	2	1	324	0.09
right	right	right	wrong	1	9	1	1	2	2	1	36	0.01
wrong	right	right	wrong	3	9	1	1	1	2	1	54	0.01
right	wrong	right	wrong	1	1	1	1	1	1	1	1	0.00
wrong	wrong	right	wrong	3	1	1	1	2	1	1	6	0.00
right	right	wrong	wrong	1	9	3	1	2	1	2	108	0.03
wrong	right	wrong	wrong	3	9	3	1	1	1	2	162	0.04
right	wrong	wrong	wrong	1	1	3	1	1	2	2	12	0.00
wrong	wrong	wrong	wrong	3	1	3	1	2	2	2	72	0.02

- Table is too expensive for decoding general UGMs.
 - We can't enumerate k^d possible configurations.

Inference without Edges

• To see idea behind more efficient methods, first let's consider empty graph:

$$p(x) = \frac{1}{Z} \prod_{j=1}^{d} \phi_j(x_j).$$

• If the x_i are binary, Z is sum of the 2^d products in the table:

$$Z = \sum_{x_1=0}^{1} \sum_{x_2=0}^{1} \cdots \sum_{x_d=0}^{1} \prod_{j=1}^{d} \phi_j(x_j).$$

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- If the x_j have k states, Z is the sum of k^d products over d variables.
- This looks hard, but independence lets us factorize into product of d simple sums.
 This trick was previously used in the EM notes.

Inference without Edges

• We can start by writing

$$Z = \sum_{x_1=0}^{1} \sum_{x_2=0}^{1} \cdots \sum_{x_{d-1}=0}^{1} \sum_{x_d=0}^{1} \left(\prod_{j=1}^{d-1} \phi_j(x_j) \right) \phi_d(x_d)$$

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• Now use $\sum_i ab_i = a \sum_i b_i$ to take terms not depending on x_d outside sum:

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$$= \sum_{x_1=0}^{1} \sum_{x_2=0}^{1} \cdots \sum_{x_{d-1}=0}^{1} \prod_{j=1}^{d-1} \phi_j(x_j) \underbrace{\sum_{x_d=0}^{1} \phi_d(x_d)}_{Z_d}$$

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• Now take the constant Z_d outside all the sums,

$$Z = Z_d \sum_{x_1=0}^{1} \sum_{x_2=0}^{1} \cdots \sum_{x_{d-1}=0}^{1} \prod_{j=1}^{d-1} \phi_j(x_j)$$

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• If we repeat these steps we obtain

$$Z = Z_d Z_{d-1} \sum_{x_1=0}^{1} \sum_{x_2=0}^{1} \cdots \sum_{x_{d-2}=0}^{1} \prod_{j=1}^{d-2} \phi_j(x_j),$$

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and if we keep going we get

$$Z = Z_d Z_{d-1} \dots Z_1 = \prod_{j=1}^d Z_j.$$

Inference without Edges

• Plugging in the definition of Z_j we get

$$Z = \prod_{j=1}^{d} \sum_{x_j=0}^{1} \phi_j(x_j),$$

so for independent variables Z is a product of d two-term sums.

• If each variable has k states, it costs O(dk) to compute.

Inference without Edges

 $\bullet\,$ Plugging in the definition of Z_{j} we get

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so for independent variables \boldsymbol{Z} is a product of \boldsymbol{d} two-term sums.

- If each variable has k states, it costs O(dk) to compute.
- By similar logic, we have $p(x_j) = \phi_j(x_j)/Z_j$ and can thus be computed in O(s).
- We could plug this back into the UGM to get

$$p(x) = \frac{1}{Z} \prod_{j=1}^{d} \phi_j(x_j) = \frac{1}{\prod_{j=1}^{d} Z_j} \prod_{j=1}^{d} \phi_j(x_j)$$
$$= \prod_{j=1}^{d} \frac{\phi_j(x_j)}{Z_j} = \prod_{j=1}^{d} p(x_j),$$

and this DAG representation allows ancenstral sampling in O(dk).

Decoding and Inference without Edges

• Since $\max_i \{ab_i\} = a \max_i \{b_i\}$ for $a \ge 0$, can use same logic for decoding:

$$\tilde{p}(x^*) = \max_{x} p(x)$$

$$= \max_{x_1} \max_{x_2} \cdots \max_{x_{d-1}} \max_{x_d} \prod_{j=1}^d \phi_j(x_j)$$

$$= \max_{x_1} \max_{x_2} \cdots \max_{x_{d-1}} \prod_{j=1}^{d-1} \phi_j(x_j) \max_{x_d} \phi_d(x_d)$$

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Tedious way of showing you set x_j to maximize its own potential.
"Generalized distributive law": work for many "+"' and "*"' operations:
E.g., commutative semi-rings (Gaussian elimination, fast Fourier transform).

Chain-Structured Graphs

General Graphs



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Computer Science Graduate Markov Model

- Computer Science Graduate Careers Markov chain:
 - Variable x_1 can be in one of three states:

State	Probability	Description
Industry	0.60	They work for a company or own their own company.
Grad School	0.30	They are trying to get a Masters or PhD degree.
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• Variable x_t only depends on x_{t-1} :

From\to	Video Games	Industry	Grad School	Video Games (with PhD)	Industry (with PhD)	Academia	Deceased
Video Games	0.08	0.90	0.01	0	0	0	0.01
Industry	0.03	0.95	0.01	0	0	0	0.01
Grad School	0.06	0.06	0.75	0.05	0.05	0.02	0.01
Video Games (with PhD)	0	0	0	0.30	0.60	0.09	0.01
Industry (with PhD)	0	0	0	0.02	0.95	0.02	0.01
Academia	0	0	0	0.01	0.01	0.97	0.01
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• So the probability of a sequence is

$$p(x_1, x_2, \dots, x_n) = p(x_1)p(x_2|x_1)p(x_3|x_2, x_1)\dots p(x_n|x_{n-1}, x_{n-2}, \dots, x_1)$$

= $p(x_1)p(x_2|x_1)p(x_3|x_2)\dots p(x_n|x_{n-1}).$

Markov Chain Models

• This is a special case of a UGM

$$p(x_1, x_2, \dots, x_n) = \phi_1(x_1) \prod_{i=2}^n \phi(x_i, x_{i-1}),$$

with a chain-structured dependency:

$$X_1 - X_2 - X_3 - X_4 - X_5 - X_6 - X_7$$

Markov Chain Models

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- Homogeneous chain: edge potentials are constant across time.
- Markov chains are ubiquitous in sequence/time-series models:

9 Applications 0.1 Disusion 9.2 Chemistry 9.3 Testina 9.4 Speech Recognition 0.5 Information principal 9.7 Internet applications 9.8 Statistics 9.9 Economics and finance 9.10 Social sciences 9.11 Mathematical biology 9.12 Genetics 9.13 Games Q 14 Muelo 9.15 Basebal 9.16 Markov text generators

General Chain-Structured UGM

• The general class of chain-structured UGMs is

$$p(x_1, x_2, \dots, x_n) \propto \prod_{i=1}^n \phi_i(x_i) \prod_{i=2}^n \phi_{i,i-1}(x_i, x_{i-1}),$$

 $(x_t \text{ could depend on future things that might happen})$

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• Includes hidden Markov models (discrete) and Kalman filters (Gaussian):



- O_i are observations (included in ϕ_i) and x_j are hidden states you want.
- Probably the most widely-used time-series models.

Applications of HMMs and Kalman Filters

Applications [edit]

HMMs can be applied in many fields where the goal is to recover a data sequence that is not immediately observable (but other data that depend on the sequence are). Applications include:

- . Single Molecule Kinetic analysis^[16]
- . Cryptanalysis
- . Speech recognition
- . Speech synthesis
- . Part-of-speech tagging
- . Document Separation in scanning solutions
- . Machine translation
- . Partial discharge
- . Gene prediction
- . Alignment of bio-sequences
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- . Protein folding^[17]
- . Metamorphic Virus Detection^[18]
- . DNA Motif Discovery^[19]

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Applications [edit]

- . Attitude and Heading Reference Systems
- . Autopilot
- . Battery state of charge (SoC) estimation^{[39][40]}
- . Brain-computer interface
- . Chaotic signals
- Tracking and Vertex Fitting of charged particles in Particle Detectors^[41]
- . Tracking of objects in computer vision
- . Dynamic positioning

- Economics, in particular macroeconomics, time series analysis, and econometrics^[42]
- . Inertial guidance system
- . Orbit Determination
- . Power system state estimation
- . Radar tracker
- . Satellite navigation systems
- . Seismology^[43]
- . Sensorless control of AC motor variable-frequency

- . Simultaneous localization and mapping
- . Speech enhancement
- . Visual odometry
- . Weather forecasting
- . Navigation system
- . 3D modeling
- . Structural health monitoring
- . Human sensorimotor processing^[44]

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- But variables are not independent:
 - We can't use our nice argument for empty graphs.
- But decoding in chains is not NP-hard:
 - Conditional independence structure yields efficient algorithms (Viterbi decoding).

• For Markov chains we have

$$p(x^*) = \max_{x} p(x_1) \prod_{j=2}^{d} p(x_j | x_{j-1})$$
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and again using that $\max_i ab_i = a \max_i b_i$ we get

$$p(x^*) = \max_{x_d} \max_{x_{d-1}} \dots \max_{x_2} \prod_{j=3}^d p(x_j | x_{j-1}) \underbrace{\max_{x_1} p(x_1) p(x_2 | x_1)}_{V(2, x_2)}.$$

• Not as nice as before: inner-most max is not a constant:

• It depends on x_2 so we can't take it outside sum over x_2 .

p

Decoding in Chain-Structured Models

• Let's just store the k values of $\{V(2,1),V(2,2),\ldots,V(2,k)\}$ and keep going,

$$(x^*) = \max_{x_d} \max_{x_{d-1}} \dots \max_{x_3} \max_{x_2} \prod_{j=3}^d p(x_j | x_{j-1}) V(2, x_2)$$

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$$p(x^*) = \max_{x_d} \max_{x_{d-1}} \dots \max_{x_4} \prod_{j=5}^d p(x_j | x_{j-1}) \underbrace{\max_{x_3} p(x_4 | x_3) V(3, x_3)}_{V(4, x_4)}$$
$$= \max_{x_d} V(d, x_d).$$

- $\bullet~{\rm The}~V$ functions summarize everything you need to know about the past.
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- A special case of dynamic programming:
 - Optimal solution is defined through recursive calls,

$$V(j, x_{j+1}) = \max_{x_j} p(x_{j+1}|x_j) V(j, x_j).$$

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so we can solve the problem by storing answers to recursive calls.

- Viterbi decoding algorithm for general chain-structured UGMs:
 - Forward phase:

$$V_{1,s} = \phi_1(s), \quad V_{i,s} = \max_{s'} \{ \phi_i(s) \phi_{i,i-1}(s,s') V_{i-1,s'} \},$$

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- For the CS grad student Markov model with n = 60:
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 - Optimal decoding might not look like 'typical' state.
 - Optimal decoding would be different with inhomogeneous chain.
 - Optimal decoding would be different if we changed n.

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• Marginals are given by $p(x_i = s) \propto V_{i,s} B_{i,s}$.
Marginals in CS Grad Markov Chain



Sampling in Chain-Structured Models

- Sampling in Markov chains by ancestral sampling:
 - Sample time 1 based on $p(x_1)$.
 - Sample time t based on time t-1 using $p(x_t|x_{t-1})$.
 - Simulates the process forward from the beginning.

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 - Simulates the process forward from the beginning.
- Forward-filter backward-sample algorithm for general case:
 - Forward phase (same as before):

$$V_{1,s} = \phi_1(s), \quad V_{i,s} = \sum_{s'} \phi_i(s)\phi_{i,i-1}(s,s')V_{i-1,s'}.$$

- Backward phase: sample x_n now that we have $p(x_n)$, then sample time (t-1) based on $V_{t-1,s}$ and x_t .
- Simulates the process backwards from the end.

Samples in CS Grad Markov Chain

Samples are more informative about what the model looks like:



Chain-Structured Graphs

General Graphs



DAGs vs. UGMs

- 2 Empty Graphs
- Chain-Structured Graphs



Decoding, Inference, and Sampling in General Graphs

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Decoding, Inference, and Sampling in General Graphs

- What if we want to go beyond chains? Can't we apply same logic?
 - Yes, but there is going to be a problem...
- Consider a simple 4-node grid-structure UGM:

 $p(x) \propto \phi_{12}(x_1, x_2)\phi_{23}(x_2, x_3)\phi_{34}(x_3, x_4)\phi_{14}(x_1, x_4).$



Variable Elimination in General Graphs

• We have that Z is defined by

$$Z = \sum_{x_4} \sum_{x_3} \sum_{x_2} \sum_{x_1} \phi_{12}(x_1, x_2) \phi_{23}(x_2, x_3) \phi_{34}(x_3, x_4) \phi_{14}(x_1, x_4)$$

=
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so now x_j our V_{24} function has k^2 values instead of k.

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so now x_j our V_{24} function has k^2 values instead of k.

• Continuing, we get

$$Z = \sum_{x_4} \sum_{x_3} \phi_{34}(x_3, x_4) \underbrace{\sum_{x_2} \phi_{23}(x_2, x_3) V_{24}(x_2, x_4)}_{V_{34}(x_3, x_4)},$$

and so on. The total cost will now be $O(dk^3)$.

• This strategy is called variable eliminiation.

Variable Elimination in General Graphs

 $\bullet~$ If we add the edge $(1,3)\mbox{,}$



we get

$$Z = \sum_{x_4} \sum_{x_3} \sum_{x_2} \sum_{x_1} \phi_{12}(x_1, x_2) \phi_{13}(x_1, x_3) \phi_{23}(x_2, x_3) \phi_{34}(x_3, x_4) \phi_{14}(x_1, x_4)$$

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so now we have a V_{234} function with k^3 possible values.

• Same $O(dk^4)$ cost of exhaustive enumeration.

Variable Elimination in General Graphs

• The cost also changes if we change the order of the sums.

Variable Elimination in General Graphs

- The cost also changes if we change the order of the sums.
- Consider chain-structured graph with sums in a different order:



• So even though we have a chain, we have a V with k^3 values instead of k.

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 - For a d_1 by d_2 grid, $\omega = \min\{d_1, d_2\}$.

Variable Elimination and Treewdith

• Trees have $\omega = 1$, decoding/inference/sampling costs $O(dk^2)$.



• A loop has $\omega = 2$, cost is $O(dk^3)$.



• A time-series with 3 variables has $\omega = 3$, cost is $O(dk^4)$.



Chain-Structured Graphs

General Graphs



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Summary

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Summary

- Markov blanket is set of nodes that make x_j independent of all others.
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- Decoding/inference/sampling with different graph structures:
 - Factorizing sum for independent distributions.
 - Viterbi decoding and forward-backward for chains.
 - Variable eliminiation for general graphs.
- I will be gone for the next 3 lectures:
 - Michael Gelbart will introduce deep learning and Bayesian stats.
 - Then we'll have the midterm.
 - Then I'll cover advanced topics in graphical models, deep learning, and Bayesian stats.