

BLC: BL Defined Using Compatible Reduction

CPSC 509: Programming Language Principles

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14 October 2015

Syntax

$$\begin{aligned}
 t &\in \text{TERM}, \quad b \in \mathbb{B}, \quad x \in \text{VAR}, \\
 t &::= \text{true} \mid \text{false} \mid \text{if } t \text{ then } t \text{ else } t \mid x \mid \text{let } x = t \text{ in } t \\
 b &::= \text{true} \mid \text{false}
 \end{aligned}$$

Semantics

$$\begin{aligned}
 v &\in \text{VALUE}, \quad r \in \text{REDEX}, \quad C \in \text{CTXT}, \quad p \in \text{PGM} = \{t \in \text{TERM} \mid FV(t) = \emptyset\}, \quad o \in \text{OBS} = \mathbb{B} \\
 v &::= b \mid x \\
 C &::= \square \mid C[\text{if } \square \text{ then } t \text{ else } t] \mid C[\text{let } x = \square \text{ in } t] \\
 &\quad \mid C[\text{if } t \text{ then } \square \text{ else } t] \mid C[\text{if } t \text{ then } t \text{ else } \square] \mid C[\text{let } x = t \text{ in } \square] \\
 r &::= \text{if } v \text{ then } t \text{ else } t \mid \text{let } x = v \text{ in } t
 \end{aligned}$$

Free Variables

$$\begin{aligned}
 FV : \text{TERM} &\rightarrow \mathcal{P}(\text{VAR}) \\
 FV(b) &= \emptyset \\
 FV(\text{if } t_1 \text{ then } t_2 \text{ else } t_3) &= FV(t_1) \cup FV(t_2) \cup FV(t_3) \\
 FV(x) &= \{x\} \\
 FV(\text{let } x = t_1 \text{ in } t_2) &= FV(t_1) \cup (FV(t_2) \setminus \{x\})
 \end{aligned}$$

All Variables References

$$\begin{aligned}
 Vars : \text{TERM} &\rightarrow \mathcal{P}(\text{VAR}) \\
 Vars(b) &= \emptyset \\
 Vars(\text{if } t_1 \text{ then } t_2 \text{ else } t_3) &= Vars(t_1) \cup Vars(t_2) \cup Vars(t_3) \\
 Vars(x) &= \{x\} \\
 Vars(\text{let } x = t_1 \text{ in } t_2) &= Vars(t_1) \cup Vars(t_2)
 \end{aligned}$$

Capture-Avoiding Substitution

$$\begin{aligned}
 [t/x] : \text{TERM} &\rightarrow \text{TERM} \\
 [t/x]b &= b \\
 [t/x]x &= t \\
 [t/x]x_0 &= x_0 \text{ if } x \neq x_0 \\
 [t/x](\text{if } t_1 \text{ then } t_2 \text{ else } t_3) &= \text{if } [t/x]t_1 \text{ then } [t/x]t_2 \text{ else } [t/x]t_3 \\
 [t/x](\text{let } x = t_1 \text{ in } t_2) &= \text{let } x = [t/x]t_1 \text{ in } t_2 \\
 [t/x](\text{let } x_0 = t_1 \text{ in } t_2) &= \text{let } x_1 = [t/x]t_1 \text{ in } [t/x][x_1/x_0]t_2 \\
 &\quad \text{if } x \neq x_0 \\
 \text{where } x_1 &= \begin{cases} x_0 & x \notin FV(t_2) \vee x_0 \notin FV(t) \\ \min_{x \in \text{VAR}} (x \notin Vars(t) \cup Vars(t_2)) & \text{otherwise} \end{cases}
 \end{aligned}$$

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$\rightsquigarrow \subseteq \text{REDEX} \times \text{TERM}$ **Notions of Reduction**

if true then t_2 else $t_3 \rightsquigarrow t_2$

if false then t_2 else $t_3 \rightsquigarrow t_3$

let $x = v$ in $t \rightsquigarrow [v/x]t$

$\longrightarrow \subseteq \text{TERM} \times \text{TERM}$ **Single-Step Reduction**

$$\frac{t \rightsquigarrow t'}{C[t] \longrightarrow C[t']}$$

$eval_c : \text{PGM} \rightarrow \text{OBS}$
 $eval_c(t) = b$ if $t \longrightarrow^* b$