CPSC 427 Video Game Programming



AI Reloaded

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Overview

- 1. Recap Behaviour trees
- 2. Shortest path and other search algorithms
- 3. Debugging (if time permits, likely next time)



Milestone 1

- Due on Friday
- Update the development plan based on proposal feedback
- Submit an individual progress & feedback report
 - on Canvas
 - Like the readme in the assignment
 - It should summarise your own contribution to the milestone
 - Did you go beyond a basic implementation?
 - Which of your features need an explanation?
 - Did you do more than expected/compared to your teammates?
 - Any issues with teamwork? react early: try to resolve internally, contact us if stuck



Tutorial Wednesday (tomorrow)

4-5 pm: implementation details

- Behaviour trees
- Advanced ECS
- Advanced OpenGL

5-6 pm: team-TA meetings



Recap: Finite State Machine (FSM)





Finite State Machines: States + Transitions





FSM Example: Pac-Man Ghosts



Behaviour Trees

- flow of decision making of an AI agent
- tree structured
- Each frame:
- Visit nodes from root to leaves
 - depth-first order
 - check currently running node
 - succeeds or fails:
 - return to parent node and evaluate its Success/Failure
 - the parent may call new branches in sequence or return Success/Failure
 - continues running: recursively return Running till root (usually)
 - Upon failure, return to the root of the behaviour tree! Start again!





Schematic examples



https://www.gamasutra.com/blogs/ChrisSimpson/20140717/2 21339/Behavior_trees_for_AI_How_they_work.php



Types





Behaviour Tree Elements

Leaf node

- A custom function, does the actual work
- Returns Running/Success/Failure

Decorator node

- has a single child
- Passes on Running/Success/Failure from child
- may invert Success/Failure

Composite node

- has one or more children
- returns 'Running' until children stopped running







Useful Composites

Sequence

- execute all children in order
- Success if all children succeed (= AND)

Selector

- execute all children in order
- return Success if any child succeeded (= OR)

Random Selectors / Sequences

Randomized order of above composites







Useful Decorators

Inverter

- Negates success/failure
 Succeeder
- always returns success
 Repeater
- Repeat child N times Repeat Until Fail
- Repeat until child fails



return "Success";





Leaf Nodes

Functionality

- *init(...)*
 - Called by parent to initialize
 - Sets state to Running
 - Not called gain before returning Success/Failure

process()

- Called every frame/tick the node is running
- Does internal processing, interacts with the world
- Returns Running/Success/Failure

Example: Walk to goal location

 Sets goal position for path finding

- Computes shortest path
- Sets character velocity
- Returns
 - success: Reached destination
 - failure: No path found
 - running: En route



Running

Leaf

Composite

Leaf

Leaf

Decorator Running

Runnina

Early exit?

- All parents of the currently running leaf node are running too
- A node early in the tree can return Success/Failure
 - Terminates children implicitly

Example



- abort sleeping
- init running node

- Trying again?
 - Re-initialize children with new parameters to init(...)

- try to sleep if alarm is off
- init sleeping node



Strategy

- Given current state, determine **BEST** next move
- Short term: best among immediate options
- Long term: what brings something closest to a goal
 - How?
 - Search for path to best outcome
 - Across states/state parameters



ivanhoe.pro



Pathfinding

• How do I get from point A to point B?





DFS: Depth First Search

Explore each path on the frontier until its end (or until a goal is found) before considering any other path.





Breadth-first search (BFS)

 Explore all paths of length L on the frontier, before looking at path of length L + 1





Breadth-first



https://en.wikipedia.org/wiki/Dijkstra%27s_algorithm



When to use BFS vs. DFS?

The search graph has cycles or is infinite

BFS

• We need the shortest path to a solution

BFS

There are only solutions at great depth

DFS

• There are some solutions at shallow depth

BFS

• No way the search graph will fit into memory



Search with Costs



Def.: The cost of a path is the sum of the costs of its arcs

$$\operatorname{cost}(\langle n_0, \dots, n_k \rangle) = \sum_{i=1}^k \operatorname{cost}(\langle n_{i-1}, n_i \rangle)$$

Want to find the solution that minimizes cost



Example: Tower Defence



Normal unit motion cost:

- Street: cost 1
- Other: cost infinity

Boss unit: which shortcuts will it take?

- Street: cost 1
- Dirt road: cost 5
- Grass: cost 50
- Purple stuff: cost infinity



Slide 24



Lowest-Cost-First Search (LCFS)

- Lowest-cost-first search finds the path with the lowest cost to a goal node
- At each stage, it selects the path with the lowest cost on the frontier.
- The frontier is implemented as a priority queue ordered by path cost.



Use of search

- Use search to determine next state (next state on shortest path to goal/best outcome)
- Measures:
 - Evaluate goal/best outcome
 - Evaluate distance (shortest path in what metric?)

Problems:

- Cost of full search (at every step) can be prohibitive
- Search in adversarial environment
 - Player will try to outsmart you



Heuristic Search

- Blind search algorithms do not take goal into account until they reach it
- We often have estimates of distance/cost from node n to a goal node
- Estimate = search heuristic
 - a scoring function h(x)



Best First Search (BestFS)

- Best First: always choose the path on the frontier with the smallest h value
 - Frontier = priority queue ordered by h
 - Once reach goal can discard most unexplored paths...
 - Why?
 - Worst case: still explore all/most space
 - Best case: very efficient
- Greedy: (only) expand path whose last node seems closest to the goal
 - Get solution that is **locally** best



A* search



A* Search



- A* search takes into account both
 - c(p) = cost of path p to current node
 - h(p) = heuristic value at node p (estimated "remaining" path cost)
- Let f(p) = c(p) + h(p).
 - f(p) is an estimate of the cost of a path from the start to a goal via p.
 actual estimate



A* always chooses the path on the frontier with the lowest estimated distance from the start to a goal node constrained to go via that path.



A* implementation

- 1. Initialize open and closed lists.
 - Put starting node on open list.
- 2. While open list is not empty:
 - Find node with smallest f on the list, call it q
 - Pop q off of open list
 - Find q's "successors", and set their parent nodes to q



A* implementation

- 1. Initialize open, closed lists. Put starting node on open list.
- 2. While open list is not empty:
 - Find node with smallest f on the list, call it q
 - Pop q off of open list
 - Find q's "successors", and set their parent nodes to q
 - For each successor:
 - If successor is the goal, done!
 - c(successor) = c(q) + d(q,successor)
 h(successor) = D(goal, successor)
 - If successor already exists in open list with lower
 f = c + h, skip it
 - If successor already exists in closed list with lower f, skip it
 - Otherwise, add successor to open list



A* implementation

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 - Find node with smallest f on the list, call it q
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 - For each successor:
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 - g(successor) = g(q) + d(q,successor)
 h(successor) = d(goal, successor)
 - If successor already exists in open list with lower f, skip it
 - If successor already exists in closed list with lower f, skip it
 - Otherwise, add successor to open list

Put q on closed list



A* search

Key idea: H is a heuristic, and not the real distance:

h(p,q) = |(p.x - q.x)| + |(p.y - q.y)|

- Manhattan distance

$$h(p,q) = sqrt((p.x - q.x)^2 + (p.y - q.y)^2)$$

- Euclidean distance

Conditions:

- a <u>heuristic function</u> is admissible if it never overestimates the cost of reaching the goal
- a <u>heuristic function</u> is said to be consistent, or monotone, if its estimate is always less than or equal to the estimated distance from any neighbouring vertex to the goal, plus the cost of reaching that neighbour



https://en.wikipedia.org/wiki/Taxicab_geometry



Two-player games



www.npr.org

Min-Max Trees

- Adversarial planning in a turn-taking environment
 - Algorithm seeks to maximize our success **F**
 - Adversary seeks to minimize **F**
- Key idea: at each step algorithm selects move that minimizes highest (estimated) value of F adversary can reach
 - Assume the opponent does what looks best

Example

(from uliana.lecturer.pens.ac.id/Kecerdasan%20Buatan/ppt/Game%20Playing/gametrees.ppt)

We are playing X, and it is now our turn.

Our options:

Number = position after each legal move

Opponent options

Here we are looking at all of the opponent responses to the first possible move we could make.

Opponent options

Opponent options after our second possibility. Not good again...

Opponent options

Opponent options => Our options

Now they don't have a way to win on their next move. So now we have to consider our responses to their responses.

Our options

We have a win for any move they make. Original position in purple is an X win.

Other options

They win again if we take our fifth move.

Summary of the Analysis

So which move should we make? ;-)

MinMax algorithm

- Traverse "game tree":
 - Enumerate all possible moves at each node.
 - The children of each node are the positions that result from making each move. A leaf is a position that is won or drawn for some side.
- Assume that we pick the best move for us, and the opponent picks the best move for him (causes most damage to us)
- Pick the move that maximizes the minimum amount of success for our side.

MinMax Algorithm

- Tic-Tac-Toe: three forms of success: Win, Tie, Lose.
 - If you have a move that leads to a Win make it.
 - If you have no such move, then make the move that gives the tie.
 - If not even this exists, then it doesn't matter what you do.

Extensions

- Challenges: In practice
 - Trees too deep/large to explore
 - Opponent not always makes the 'best' choice
 - Randomness
- Solution Heuristics
 - Rate nodes based on local information.
 - For example, in Chess "rate" a position by examining difference in number of pieces

Heuristics in MinMax

- Strategy that will let us cut off the game tree at fixed depth (layer)
- Apply heuristic scoring to bottom layer
 - instead of just Win, Loss, Tie, we have a score.
- For "our" level of the tree we want the move that yields the node (position) with highest score. For a "them" level "they" want the child with the lowest score.

Self stuy: Pseudocode

```
int Minimax(Board b, boolean myTurn, int depth) {
    if (depth==0)
        return b.Evaluate(); // Heuristic
    for(each possible move i)
        value[i] = Minimax(b.move(i), !myTurn,
    depth-1);
    if (myTurn)
        return array_max(value);
    else
        return array_min(value);
}
```

Note: we don't use an explicit tree structure. However, the pattern of recursive calls forms a tree on the call stack.

Real Minimax Example

Evaluation function applied to the leaves!

Alpha Beta Pruning

Idea: Track "window" of expectations. Use two variables

- α Best score so far at a **max** node: increases
 - At a child **min** node:
 - Parent wants **max**. To affect the parent's current α , our β cannot drop below α .
 - If β ever gets less:
 - Stop searching further subtrees *of that child*. They do not matter!
- β Best score so far at a **min** node: decreases
 - At a child **max** node.
 - Parent wants **min**. To affect the parent's current β , our α cannot get above the parent's β .
 - If α gets bigger than β :
 - Stop searching further subtrees of that child. They do not matter!

Start with an infinite window ($\alpha = -\infty, \beta = \infty$)

Alpha Beta Example I

Alpha Beta Example II

Self stuy: Pseudo Code

```
int AlphaBeta(Board b, boolean myTurn, int depth, int alpha, int beta) {
   if (depth==0)
      return b.Evaluate(); // Heuristic
   if (myTurn) {
      for(each possible move i && alpha < beta)</pre>
         alpha = max(alpha,AlphaBeta(b.move(i),!myTurn,depth-1,alpha,beta));
      return alpha;
  else {
      for (each possible move i && alpha < beta)
         beta = min(beta,AlphaBeta(b.move(i), !myTurn, depth-1,alpha,beta));
      return beta;
```


- There will be bugs...
- Strategies for Fixing?

• There will be bugs...

• Strategies for Fixing?

- Anticipate
- Reproduce
- Localize
- Use proper debugging tools

Debugging:Strategies for Fixing?

- Anticipate I
 - Unit tests
 - Logging
 - Explicit tests for "what can go wrong" (assert)
 - Anything that can go wrong will go wrong... at the worst possible time
 - State/play saving and loading speeds up debugging
 - Visual testing (early)
 - Avoid randomness (use seed for rnd)
- Reproduce
- Localize
- Use proper debugging tools

Debugging: Strategies for Fixing?

- Anticipate II: your compiler (with Wall enabled) is your friend
 - "This enables all the warnings about constructions that some users consider questionable, and that are easy to avoid"
- Reproduce
- Localize
- Use proper debugging tools

- Strategies for Fixing?
- Anticipate
- Reproduce
 - When does it happen?
 - Logging + unit tests
 - Record/load gameplay
- Localize
- Use proper debugging tools

- Strategies for Fixing?
- Anticipate
- Reproduce
- Localize
 - In time: version control
 - In place: logging
 - Divide and Conquer
 - Minimal trigger input
 - Don't guess; measure
- Use proper debugging tools

- Strategies for Fixing?
- Anticipate
- Reproduce
- Localize
- Use proper debugging tools
 - Run with debug settings on
 - Run within a debugger
 - Set breakpoints
 - Examine internal state
 - Learn debugger options

Debugging (From Waterloo ECE 155, Zarnett & Lam)

- Strategies for Fixing?
- Scientific method.
 - Observe a failure.
 - Invent a hypothesis.
 - 3 Make predictions.
 - 4 Test the predictions using experiments and observations.
- Correct? Refine the hypothesis.
- Wrong? Try again with a new hypothesis.
- Repeat

Debugging (From Waterloo ECE 155)

More (Human Factor) Strategies

- Take a Break/Sleep on it
- Code Review
 - Look through code
 - Walk someone through the code

More (Human Factor) Strategies

- Question assumptions
- Minimize randomness
 - Use same seed
- Check boundary conditions
- Disrupt parallel computations

Debugging (From Waterloo ECE 155)

More Strategies

- Know your enemy: Types of bugs
 - Standard bug (reproducible)
 - Sporadic (need to chase right input combo)
 - Heisenbug
 - Memory (not initialized or stepped on)
 - Parallel execution
 - Optimization

Hard Bugs (cheat sheet)

- Bug occurs in Release but not Debug
 - Uninitialized data or optimization issue
- Bug disappears when changing something innocuous
 - Timing or memory overwrite problem
- Intermittent problems
 - Record as much info when it does happen
- Unexplainable behavior
 - Retry, Rebuild, Reboot, Reinstall
- Internal compiler errors (not likely)
 - Full rebuild, divide and conquer, try other machines
- Suspect it's not your code (not likely)
 - Check for patches, updates, or reported bugs