

# Domination

ISCI 330 Lecture 9

February 6, 2007

# Lecture Overview

- 1 Recap
- 2 Fun Game
- 3 Domination

# Max-Min Strategies

- Player  $i$ 's **maxmin strategy** is a strategy that maximizes  $i$ 's worst-case payoff, in the situation where all the other players (whom we denote  $-i$ ) happen to play the strategies which cause the greatest harm to  $i$ .
- The **maxmin value** (or **safety level**) of the game for player  $i$  is that minimum amount of payoff guaranteed by a maxmin strategy.
- Why would  $i$  want to play a maxmin strategy?
  - a conservative agent maximizing worst-case payoff
  - a paranoid agent who believes everyone is out to get him

## Definition

The **maxmin strategy** for player  $i$  is  $\arg \max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$ , and the **maxmin value** for player  $i$  is  $\max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$ .

# Min-Max Strategies

- Player  $i$ 's **minmax strategy** in a 2-player game is a strategy that minimizes the other player  $-i$ 's best-case payoff.
- The **minmax value** of the 2-player game for player  $i$  is that maximum amount of payoff that  $-i$  could achieve under  $i$ 's minmax strategy.
- Why would  $i$  want to play a minmax strategy?
  - to punish the other agent as much as possible

## Definition

The **maxmin strategy** for player  $i$  is  $\arg \max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$ , and the **maxmin value** for player  $i$  is  $\max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$ .

## Definition

In a two-player game, the **minmax strategy** for player  $i$  is  $\arg \min_{s_i} \max_{s_{-i}} u_{-i}(s_1, s_2)$ , and the **minmax value** for player  $i$  is  $\min_{s_i} \max_{s_{-i}} u_{-i}(s_1, s_2)$ .

# Minmax Theorem

## Theorem (Minmax theorem (von Neumann, 1928))

*In any finite, two-player, zero-sum game, in any Nash equilibrium each player receives a payoff that is equal to both his maxmin value and his minmax value.*

- The maxmin value for one player is equal to the minmax value for the other player. By convention, the maxmin value for player 1 is called the **value of the game**.
- For both players, the set of maxmin strategies coincides with the set of minmax strategies.
- Any maxmin strategy profile (or, equivalently, minmax strategy profile) is a Nash equilibrium. Furthermore, these are all the Nash equilibria. Consequently, all Nash equilibria have the same payoff vector (namely, those in which player 1 gets the value of the game).

# How to find maxmin and minmax strategies

Consider maxmin strategies for player  $i$  in a 2-player game.

- Notice that  $i$ 's maxmin strategy depends only on  $i$ 's utilities
  - thus changes to  $-i$ 's utilities do not change  $i$ 's maxmin strategy
- Consider the game where player  $i$  has the same utilities as before, but player  $-i$ 's utilities are replaced with the negatives of  $i$ 's utilities
  - this is now a zero-sum game
- Because of the minmax theorem, we know that any Nash equilibrium strategy in this game is also a maxmin strategy
  - Thus, find player  $i$ 's equilibrium strategy in the new game and we have  $i$ 's maxmin strategy in the original game
- We can use a similar approach for minmax.

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# Traveler's Dilemma

*Two travelers purchase identical African masks while on a tropical vacation. Their luggage is lost on the return trip, and the airline asks them to make independent claims for compensation. In anticipation of excessive claims, the airline representative announces: "We know that the bags have identical contents, and we will entertain any claim between \$180 and \$300, but you will each be reimbursed at an amount that equals the minimum of the two claims submitted. If the two claims differ, we will also pay a reward  $R$  to the person making the smaller claim and we will deduct a penalty  $R$  from the reimbursement to the person making the larger claim."*



# Traveler's Dilemma

- Action: choose an integer between 180 and 300
- If both players pick the same number, they both get that amount as payoff
- If players pick a different number:
  - the low player gets his number ( $L$ ) plus some constant  $R$
  - the high player gets  $L - R$ .
- Play this game *once* with a partner; play with as many different partners as you like.
  - $R = 5$ .

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  - the high player gets  $L - R$ .
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  - $R = 5$ .
  - $R = 180$ .

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- What is the equilibrium?
  - $(180, 180)$  is the only equilibrium, for all  $R \geq 2$ .
- What happens?
  - with  $R = 5$  most people choose 295–300
  - with  $R = 180$  most people choose 180

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# Domination

- Let  $s_i$  and  $s'_i$  be two strategies for player  $i$ , and let  $S_{-i}$  be the set of all possible strategy profiles for the other players

## Definition

$s_i$  **strictly dominates**  $s'_i$  if  $\forall s_{-i} \in S_{-i}, u_i(s_i, s_{-i}) > u_i(s'_i, s_{-i})$

## Definition

$s_i$  **weakly dominates**  $s'_i$  if  $\forall s_{-i} \in S_{-i}, u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i})$  and  $\exists s_{-i} \in S_{-i}, u_i(s_i, s_{-i}) > u_i(s'_i, s_{-i})$

## Definition

$s_i$  **very weakly dominates**  $s'_i$  if  $\forall s_{-i} \in S_{-i}, u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i})$



# Equilibria and dominance

- If one strategy dominates all others, we say it is **dominant**.
- A strategy profile consisting of dominant strategies for every player must be a Nash equilibrium.
  - An equilibrium in strictly dominant strategies must be unique.

# Equilibria and dominance

- If one strategy dominates all others, we say it is **dominant**.
- A strategy profile consisting of dominant strategies for every player must be a Nash equilibrium.
  - An equilibrium in strictly dominant strategies must be unique.
- Consider Prisoner's Dilemma again
  - not only is the only equilibrium the only non-Pareto-optimal outcome, but it's also an equilibrium in strictly dominant strategies!