

Utility and Foundations (2)

Modeling Human Strategic Behavior

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Lecture Overview

Recap

Proof sketch

Fun Game!

Recap: Axioms

- Completeness

$$o_1 \succeq o_2 \text{ or } o_2 \succeq o_1$$

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- Continuity

$$o_1 \succ o_2 \succ o_3 \implies \exists p \in [0, 1] : o_2 \sim [p:o_1, (1-p):o_3]$$

Recap: Representation Theorem

Theorem [von Neumann & Morgenstern, 1944]

Suppose that a preference relation \succeq satisfies the axioms **Completeness**, **Transitivity**, **Monotonicity**, **Substitutability**, **Decomposability**, and **Continuity**.

Then there exists a function $u : O \rightarrow \mathbb{R}$ such that

1. $\forall o_1, o_2 \in O : o_1 \succeq o_2 \iff u(o_1) \geq u(o_2)$, and
2. $\forall [p_1 : o_1, \dots, p_k : o_k] \in O : u([p_1 : o_1, \dots, p_k : o_k]) = \sum_{j=1}^k p_j u(o_j)$.

That is, there exists a utility function u that represents \succeq .

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2. Construct $u(o) = p$ such that $o \sim [p: o^+, (1 - p): o^-]$
3. Substitutability lets us replace everything with these “canonical” lotteries; Monotonicity lets us assert the ordering between them.

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for all $b \in \mathbb{R}$ and $c > 0$

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Fun Game: Buying Lottery Tickets

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3. How much would you pay to play the lottery

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If you knew that the last seven draws had been 5, 5, 7, 5, 9, 9, 5?

Beyond von Neumann & Morgenstern

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- **Question:** If two rational agents have different prices for $[0.3:\$5, 0.3:\$7, 0.4:\$9]$, what does that suggest about their **preferences for money**?

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- von Neumann and Morgenstern's utility theory assumes **known, objective** probabilities.
- There are other representation theorems [e.g., Savage 1954] that state that rational agents must (a) have probabilistic beliefs, (b) update those beliefs as if by conditioning, (c) maximize the expected value of some utility function wrt them

Utility Summary

Utility theory proves that agents whose preferences obey certain simple axioms about preferences over lotteries must act as if they were maximizing the expected value of a scalar function.

- **“Rational”** agents are those whose behaviour satisfies the **axioms**
- ***If you don’t buy the axioms, then you shouldn’t buy that this theorem is about rational behavior.***
- Conversely, if you don’t buy that rational agents must behave in this way, then there must be at least one axiom that you disagree with.

This approach extends to “subjective” probabilities:

- Axioms about **preferences over uncertain “acts”** that do not describe how agents manipulate probabilities.

Game Representations

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Lecture Overview

Normal-Form

Repeated

Extensive Form

Bayesian Games

TCP Backoff Game



Should you send your packets using correctly-implemented TCP (which has a “backoff” mechanism) or using a defective implementation (which doesn’t)?

- Consider this situation as a two-player game:
 - **both use a correct implementation:** both get 1 ms delay
 - **one correct, one defective:** 4 ms delay for correct, 0 ms for defective
 - **both defective:** both get a 3 ms delay.

TCP Backoff Game

- Consider this situation as a two-player game:
 - **both use a correct implementation:** both get 1 ms delay
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 - **both defective:** both get a 3 ms delay.
- Go into a breakout room. Play once with each person.
- Questions:
 - What **action** should a player of the game take?
 - Would all users behave **the same** in this scenario?
 - What global **patterns of behaviour** should the system designer expect?
 - Under what **changes to the delay numbers** would behavior be the same?
 - What effect would **communication** have?
 - Does it matter if I believe that my opponent is **rational**?

Defining Games

- Finite, n -person game: $\langle N, A, u \rangle$:
 - N is a finite set of n **players**, indexed by i
 - $A = \langle A_1, \dots, A_n \rangle$ is a tuple of **action sets** for each player i
 - $a \in A$ is an **action profile**
 - $u = \langle u_1, \dots, u_n \rangle$, a **utility function** for each player, where $u_i : A \mapsto \mathbb{R}$
- Writing a 2-player game as a **matrix**:
 - row player is player 1, column player is player 2
 - rows are actions $a \in A_1$, columns are $a' \in A_2$
 - cells are outcomes, written as a tuple of utility values for each player

Games in Matrix Form

Here's the **TCP Backoff Game** written as a matrix ("normal form").

| | | |
|----------|----------|----------|
| | <i>C</i> | <i>D</i> |
| <i>C</i> | -1, -1 | -4, 0 |
| <i>D</i> | 0, -4 | -3, -3 |

More General Form

Prisoner's dilemma is any game

| | | |
|----------|----------|----------|
| | <i>C</i> | <i>D</i> |
| <i>C</i> | a, a | b, c |
| <i>D</i> | c, b | d, d |

with $c > a > d > b$.

Games of Pure Competition

Players have **exactly opposed** interests

- There must be precisely two players (*otherwise they can't have exactly opposed interests*)
- For all action profiles $a \in A$, $u_1(a) + u_2(a) = c$ for some constant c
 - Special case: zero sum
- Thus, we only need to store a utility function for one player
 - in a sense, it's a one-player game

Matching Pennies

One player wants to **match**; the other wants to **mismatch**.

| | Heads | Tails |
|-------|-------|-------|
| Heads | 1, -1 | -1, 1 |
| Tails | -1, 1 | 1, -1 |

Rock-Paper-Scissors

Generalized matching pennies.

| | Rock | Paper | Scissors |
|----------|-------|-------|----------|
| Rock | 0, 0 | -1, 1 | 1, -1 |
| Paper | 1, -1 | 0, 0 | -1, 1 |
| Scissors | -1, 1 | 1, -1 | 0, 0 |

...Believe it or not, there's an annual international competition for this game!

Games of Cooperation

Players have **exactly the same** interests.

- no conflict: all players want the same things
- $\forall a \in A, \forall i, j, u_i(a) = u_j(a)$
- we often write such games with a single payoff per cell
- why are these even still games?

Coordination Game

Which **side of the road** should you drive on?

| | Left | Right |
|-------|------|-------|
| Left | 1, 1 | 0, 0 |
| Right | 0, 0 | 1, 1 |

General Games: Battle of the Sexes

The most interesting games combine elements of **cooperation and competition**.

| | | |
|---|------|------|
| | B | F |
| B | 2, 1 | 0, 0 |
| F | 0, 0 | 1, 2 |

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Play this game in breakout rooms. Be fast!