

Auctions Introduction

CPSC 532A Lecture 20

November 21, 2006

Lecture Overview

- 1 Recap
- 2 VCG caveats
- 3 Auctions
- 4 Standard auctions
- 5 More exotic auctions

Groves Uniqueness

Theorem

An efficient social choice function $C : \mathbb{R}^{X^n} \rightarrow X \times \mathbb{R}^n$ can be implemented in dominant strategies for agents with unrestricted quasilinear utilities only if $p_i(v) = h(v_{-i}) - \sum_{j \neq i} v_j(\chi(v))$.

- it turns out that the same result also holds for the broader class of Bayes-Nash incentive-compatible efficient mechanisms.

VCG

Definition (Vickrey-Clarke-Groves (VCG) mechanism)

The *Vickrey-Clarke-Groves mechanism* is a direct quasilinear mechanism $(\mathbb{R}^{|X|n}, \chi, p)$, where

$$\chi(\hat{v}) = \arg \max_x \sum_i \hat{v}_i(x)$$

$$p_i(\hat{v}) = \sum_{j \neq i} \hat{v}_j(\chi(\hat{v}_{-i})) - \sum_{j \neq i} \hat{v}_j(\chi(\hat{v}))$$

Two definitions

Definition (Choice-set monotonicity)

An environment exhibits **choice-set monotonicity** if $\forall i, |X_{-i}| \leq |X|$.

- removing any agent weakly decreases—that is, never increases—the mechanism's set of possible choices X

Definition (No negative externalities)

An environment exhibits **no negative externalities** if $\forall i \forall x \in X_{-i}, v_i(x) \geq 0$.

- every agent has zero or positive utility for any choice that can be made without his participation

VCG Individual Rationality

Theorem

The VCG mechanism is ex-post individual rational when the choice set monotonicity and no negative externalities properties hold.

Another property

Definition (No single-agent effect)

An environment exhibits **no single-agent effect** if $\forall x, \forall i$ such that $\exists v_{-i}$ where $x \in \arg \max \sum_j v_j(x)$ there exists a choice x' that is feasible without i and that has $\sum_{j \neq i} v_j(x') \geq \sum_{j \neq i} v_j(x)$.

Theorem

The VCG mechanism is weakly budget-balanced when the no single-agent effect property holds.

Bad news

Theorem

No dominant strategy incentive-compatible mechanism is always both efficient and weakly budget balanced, even if agents are restricted to the simple exchange setting.

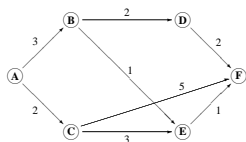
Theorem

No Bayes-Nash incentive-compatible mechanism is always simultaneously efficient, weakly budget balanced and ex-interim individual rational, even if agents are restricted to quasilinear utility functions.

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Frugality



- VCG can end up paying arbitrarily more than an agent is willing to accept (or equivalently charging arbitrarily less than an agent is willing to pay)
- Consider AC , which is not part of the shortest path.
 - If the cost of this edge increased to 8, our payment to AB would increase to $p_{AB} = (-12) - (-2) = -10$.
 - If the cost were any $x \geq 2$, we would select the path $ABEF$ and would have to make a payment to AB of $p_{AB} = (-4 - x) - (-2) = -(x + 2)$.
 - The gap between agents' true costs and the payments that they could receive under VCG is unbounded.

Privacy

- VCG requires agents to fully reveal their private information
- this private information may have value to agents that extends beyond the current interaction
 - for example, the agents may know that they will compete with each other again in the future
- it is often preferable to elicit only as much information from agents as is required to determine the social welfare maximizing outcome and compute the VCG payments.

Collusion

Example

Agent	$U(\text{build road})$	$U(\text{do not build road})$	Payment
1	200	0	150
2	100	0	50
3	0	250	0

- What happens if agents 1 and 2 both increase their declared valuations by \$50?

Collusion

Example

Agent	$U(\text{build road})$	$U(\text{do not build road})$	Payment
1	250	0	
2	150	0	
3	0	250	0

- What happens if agents 1 and 2 both increase their declared valuations by \$50?

Collusion

Example

Agent	$U(\text{build road})$	$U(\text{do not build road})$	Payment
1	250	0	100
2	150	0	0
3	0	250	0

- What happens if agents 1 and 2 both increase their declared valuations by \$50?
- The outcome is unchanged, but both of their payments are reduced.
- Thus, while no agent can gain by changing his declaration, groups *can*.

Returning profits to the agents

- we may want to use VCG to induce agents to report their valuations honestly, but may not want to make a profit by collecting money from the agents.
- Thus, we might want to find some way of returning the mechanism's profits back the agents.
- However, the possibility of receiving a rebate after the mechanism has been run changes the agents' incentives.
- In fact, even if profits are given to a charity that the agents care about, or spent in a way that benefits the local economy and hence benefits the agents, the VCG mechanism is undermined.
- Thus, burning the money collected by the mechanism is the only way ensuring that the agents' incentives are not altered!

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Motivation

- Auctions are any mechanisms for allocating resources among self-interested agents
- Very widely used
 - government sale of resources
 - privatization
 - stock market
 - request for quote
 - FCC spectrum
 - real estate sales
 - eBay

CS Motivation

- resource allocation is a fundamental problem in CS
- increasing importance of studying distributed systems with heterogeneous agents
- markets for:
 - computational resources (JINI, etc.)
 - SETI, etc.
 - autonomous agents
 - P2P systems
 - network bandwidth
- currency needn't be real money, just something scarce
 - that said, real money trading agents are also an important motivation

Formal Model

- while we think of auctions in terms of a guy with a gavel, “going-going-gone!”, they’re actually a much broader theoretical framework for resource allocation
- another way of thinking of an auction: any negotiation mechanism which is
 - mediated (auctioneer)
 - well-specified (follows rules)
 - market-based (determines an exchange in terms of currency)

Modeling Auctions

Every resource allocation mechanism in a setting with quasilinear utilities can be understood as an auction

- ascending auction: an extensive-form game with imperfect information
- sealed-bid auction: direct mechanism; a variety of payment functions are possible here
 - give the good to the person who says they need it the most
 - a non-incentive compatible mechanism with a payment function $p_i = 0$
 - charge a fixed price for the good, sell a unit of it to anyone who wants one
 - trivial allocation rule, constant payment function
- stock market
 - both buyers and sellers make bids
 - market-maker clears the market and keeps the spread between ask and buy

Auction Dimensions

- rules for bidding
 - who can bid, when
 - what is the form of a bid
 - restrictions on offers, as a function of:
 - bidder's own previous bid
 - auction state (others' bids)
 - eligibility (e.g., budget constraints)
 - expiration, withdrawal, replacement
- rules for what information is revealed
 - when to reveal what information to whom
- rules for clearing
 - when to clear
 - at intervals
 - on each bid
 - after a period of inactivity
 - allocation (who gets what)
 - payment (who pays what)

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Some popular auctions

- English
- Dutch
- First-Price
- Second-Price

English Auction

- auctioneer starts the bidding at some “reservation price”
- bidders then shout out ascending prices
- once bidders stop shouting, the high bidder gets the good at that price

Dutch Auction

- the auctioneer starts a clock at some high value; it descends
- at some point, a bidder shouts “mine!” and gets the good at the price shown on the clock

First-Price Auction

- bidders write down bids on pieces of paper
- auctioneer awards the good to the bidder with the highest bid
- that bidder pays the amount of his bid

Second-Price Auction

- bidders write down bids on pieces of paper
- auctioneer awards the good to the bidder with the highest bid
- that bidder pays the amount bid by the second-highest bidder

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Some more exotic auction types

- Japanese auction
- All-pay auction
- Continuous double auction
- Call market (“periodic clear”)

Japanese Auction

- Same as an English auction except that the auctioneer calls out the prices
- all bidders start out standing
- when the price reaches a level that a bidder is not willing to pay, that bidder sits down
 - once a bidder sits down, they can't get back up
- the last person standing gets the good
- analytically more tractable than English because jump bidding can't occur
 - consider the branching factor of the extensive form game...

All-pay auction

- sealed bid auction
- everyone pays the amount of their bid regardless of whether or not they win