

# Social Choice, Arrow's Theorem

CPSC 532A Lecture 15

October 26, 2006

# Lecture Overview

Recap

Voting Paradoxes

Fun Game

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Arrow's Theorem

# Ex-post expected utility

## Definition (*Ex-post* expected utility)

Agent  $i$ 's **ex-post expected utility** in a Bayesian game  $(N, A, \Theta, p, u)$ , where the agents' strategies are given by  $s$  and the agent' types are given by  $\theta$ , is defined as

$$EU_i(s, \theta) = \sum_{a \in A} \left( \prod_{j \in N} s_j(a_j | \theta_j) \right) u_i(a, \theta).$$

- ▶ The only uncertainty here concerns the other agents' mixed strategies, since  $i$  knows everyone's type.

# Ex-interim expected utility

## Definition (*Ex-interim* expected utility)

Agent  $i$ 's *ex-interim* expected utility in a Bayesian game  $(N, A, \Theta, p, u)$ , where  $i$ 's type is  $\theta_i$  and where the agents' strategies are given by the mixed strategy profile  $s$ , is defined as

$$EU_i(s|\theta_i) = \sum_{\theta_{-i} \in \Theta_{-i}} p(\theta_{-i}|\theta_i) \sum_{a \in A} \left( \prod_{j \in N} s_j(a_j|\theta_j) \right) u_i(a, \theta_{-i}, \theta_i).$$

- ▶  $i$  must consider every  $\theta_{-i}$  and every  $a$  in order to evaluate  $u_i(a, \theta_i, \theta_{-i})$ .
- ▶  $i$  must weight this utility value by:
  - ▶ the probability that  $a$  would be realized given all players' mixed strategies and types;
  - ▶ the probability that the other players' types would be  $\theta_{-i}$  given that his own type is  $\theta_i$ .

# Ex-ante expected utility

## Definition (*Ex-ante* expected utility)

Agent  $i$ 's **ex-ante expected utility** in a Bayesian game  $(N, A, \Theta, p, u)$ , where the agents' strategies are given by the mixed strategy profile  $s$ , is defined as

$$EU_i(s) = \sum_{\theta_i \in \Theta_i} p(\theta_i) EU_i(s|\theta_i)$$

or equivalently as

$$EU_i(s) = \sum_{\theta \in \Theta} p(\theta) \sum_{a \in A} \left( \prod_{j \in N} s_j(a_j|\theta_j) \right) u_i(a, \theta).$$

# Nash equilibrium

## Definition (Bayes-Nash equilibrium)

A **Bayes-Nash equilibrium** is a mixed strategy profile  $s$  that satisfies  $\forall i \ s_i \in BR_i(s_{-i})$ .

## Definition (*ex-post* Bayes-Nash equilibrium)

A ***ex-post* Bayes-Nash equilibrium** is a mixed strategy profile  $s$  that satisfies  $\forall \theta, \forall i, s_i \in \arg \max_{s'_i \in S_i} EU_i(s'_i, s_{-i}, \theta)$ .

- ▶ almost a **dominant strategy**, but not quite

# Social Choice

## Definition (Social choice function)

Assume a set of agents  $N = \{1, 2, \dots, n\}$ , and a set of outcomes (or alternatives, or candidates)  $O$ . Let  $L$  be the set of strict total orders on  $O$ . A **social choice function** (over  $N$  and  $O$ ) is a function  $C : L^n \rightarrow O$ .

## Definition (Social welfare function)

Let  $N, O, L$  be as above. A **social welfare function** (over  $N$  and  $O$ ) is a function  $W : L^n \rightarrow L^-$ , where  $L^-$  is the set of weak total orderings (that is, total preorders) on  $O$ .

# Some Voting Schemes

- ▶ **Plurality**
  - ▶ pick the outcome which is preferred by the most people
- ▶ **Plurality with elimination** (“instant runoff”)
  - ▶ everyone selects their favorite outcome
  - ▶ the outcome with the fewest votes is eliminated
  - ▶ repeat until one outcome remains
- ▶ **Borda**
  - ▶ assign each outcome a number.
  - ▶ The most preferred outcome gets a score of  $n - 1$ , the next most preferred gets  $n - 2$ , down to the  $n^{\text{th}}$  outcome which gets 0.
  - ▶ Then sum the numbers for each outcome, and choose the one that has the highest score
- ▶ **Pairwise elimination**
  - ▶ in advance, decide a schedule for the order in which pairs will be compared.
  - ▶ given two outcomes, have everyone determine the one that they prefer



# Condorcet Condition

- ▶ If there is a candidate who is preferred to every other candidate in pairwise runoffs, that candidate should be the winner
- ▶ While the Condorcet condition is considered an important property for a voting system to satisfy, there is not always a Condorcet winner
- ▶ sometimes, there's a cycle where  $A$  defeats  $B$ ,  $B$  defeats  $C$ , and  $C$  defeats  $A$  in their pairwise runoffs

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# Sensitivity to Losing Candidate

35 agents:  $A \succ C \succ B$

33 agents:  $B \succ A \succ C$

32 agents:  $C \succ B \succ A$

- ▶ What candidate wins under plurality voting?

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- ▶ What candidate wins under Borda voting?  $A$
- ▶ Now consider dropping  $C$ . Now what happens under both Borda and plurality?

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- ▶ What candidate wins under Borda voting?  $A$
- ▶ Now consider dropping  $C$ . Now what happens under both Borda and plurality?  $B$  wins.



# Sensitivity to Agenda Setter

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- ▶ Who wins with the ordering  $A, C, B$ ?  $B$
- ▶ Who wins with the ordering  $B, C, A$ ?  $A$

# Another Pairwise Elimination Problem

1 agent:  $B \succ D \succ C \succ A$

1 agent:  $A \succ B \succ D \succ C$

1 agent:  $C \succ A \succ B \succ D$

- ▶ Who wins under pairwise elimination with the ordering  $A, B, C, D$ ?

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1 agent:  $A \succ B \succ D \succ C$

1 agent:  $C \succ A \succ B \succ D$

- ▶ Who wins under pairwise elimination with the ordering  $A, B, C, D$ ?  $D$ .



## Another Pairwise Elimination Problem

1 agent:  $B \succ D \succ C \succ A$

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- ▶ Who wins under pairwise elimination with the ordering  $A, B, C, D$ ?  $D$ .
- ▶ What is the problem with this?

# Another Pairwise Elimination Problem

1 agent:  $B \succ D \succ C \succ A$

1 agent:  $A \succ B \succ D \succ C$

1 agent:  $C \succ A \succ B \succ D$

- ▶ Who wins under pairwise elimination with the ordering  $A, B, C, D$ ?  $D$ .
- ▶ What is the problem with this?
  - ▶ *all* of the agents prefer  $B$  to  $D$ —the selected candidate is Pareto-dominated!

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# Fun Game

- ▶ Imagine that there was an opportunity to take a one-week class trip at the end of term, to one of the following destinations:
  - ▶ (O) Orlando, FL
  - ▶ (L) London, England
  - ▶ (M) Moscow, Russia
  - ▶ (B) Beijing, China
- ▶ Construct your preference ordering

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- ▶ Construct your preference ordering
- ▶ Vote (truthfully) using each of the following schemes:
  - ▶ plurality (raise hands)

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  - ▶ plurality (raise hands)
  - ▶ plurality with elimination (raise hands)

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  - ▶ Borda (volunteer to tabulate)

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- ▶ Construct your preference ordering
- ▶ Vote (truthfully) using each of the following schemes:
  - ▶ plurality (raise hands)
  - ▶ plurality with elimination (raise hands)
  - ▶ Borda (volunteer to tabulate)
  - ▶ pairwise elimination (raise hands, I'll pick a schedule)



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# Notation

- ▶  $N$  is the set of agents
- ▶  $O$  is a finite set of outcomes with  $|O| \geq 3$
- ▶  $L$  the set of all possible preference orderings over  $O$ .
- ▶  $\succsim$  is an element of the set  $L^n$  (a preference ordering for every agent; the input to our social welfare function)
- ▶  $\succsim_W$  is the preference ordering selected by the social welfare function  $W$ .
  - ▶ When the input to  $W$  is ambiguous we write it in the subscript; thus, the social order selected by  $W$  given the input  $\succsim'$  is denoted as  $\succsim_{W(\succsim')}$ .

# Pareto Efficiency

## Definition (Pareto Efficiency (PE))

$W$  is **Pareto efficient** if for any  $o_1, o_2 \in O$ ,  $\forall i o_1 \succ_i o_2$  implies that  $o_1 \succ_W o_2$ .

- ▶ when all agents agree on the ordering of two outcomes, the social welfare function must select that ordering.

# Independence of Irrelevant Alternatives

## Definition (Independence of Irrelevant Alternatives (IIA))

$W$  is **independent of irrelevant alternatives** if, for any  $o_1, o_2 \in O$  and any two preference profiles  $\succ', \succ'' \in L^n$ ,

$\forall i (o_1 \succ'_i o_2 \leftrightarrow o_1 \succ''_i o_2)$  implies that

$o_1 \succ_{W(\succ')} o_2 \Leftrightarrow o_1 \succ_{W(\succ'')} o_2$ .

- ▶ the selected ordering between two outcomes should depend only on the relative orderings they are given by the agents.

# Nondictatorship

## Definition (Non-dictatorship)

$W$  does not have a **dictator** if  $\neg \exists i \forall o_1, o_2 (o_1 \succ_i o_2 \Rightarrow o_1 \succ_W o_2)$ .

- ▶ there does not exist a single agent whose preferences always determine the social ordering.
- ▶ We say that  $W$  is *dictatorial* if it fails to satisfy this property.

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# Arrow's Theorem

## Theorem (Arrow, 1951)

*Any social welfare function  $W$  that is Pareto efficient and independent of irrelevant alternatives is dictatorial.*

We will assume that  $W$  is both PE and IIA, and show that  $W$  must be dictatorial. The argument proceeds in four steps.

# Step 1

*If every voter puts an outcome  $b$  at either the very top or the very bottom of his preference list,  $b$  must be at either the very top or very bottom of  $\succ_W$  as well.*

Consider an arbitrary preference profile  $\succ$  in which every voter ranks some  $b \in O$  at either the very bottom or very top, and assume for contradiction that the above claim is not true. Then, there must exist some pair of distinct outcomes  $a, c \in O$  for which  $a \succ_W b$  and  $b \succ_W c$ .



# Step 1

*If every voter puts an outcome  $b$  at either the very top or the very bottom of his preference list,  $b$  must be at either the very top or very bottom of  $\succ_W$  as well.*

Now let's modify  $\succ$  so that every voter moves  $c$  just above  $a$  in his preference ranking, and otherwise leaves the ranking unchanged; let's call this new preference ordering  $\succ'$ . We know from IIA that for  $a \succ_W b$  or  $b \succ_W c$  to change, the pairwise relationship between  $a$  and  $b$  and/or the pairwise relationship between  $b$  and  $c$  would have to change. However, since  $b$  occupies an extremal position for all voters,  $c$  can be moved above  $a$  without changing either of these pairwise relationships. Thus in profile  $\succ'$  it is also the case that  $a \succ_W b$  and  $b \succ_W c$ . From this fact and from transitivity, we have that  $a \succ_W c$ . However, in  $\succ'$  every voter ranks  $c$  above  $a$  and so PE requires that  $c \succ_W a$ . We have a contradiction.