Local Search

CPSC 322 Lecture 13

Lecture Overview

- Recap solving CSP systematically
- Local search
- Constrained Optimization
- Greedy Descent / Hill Climbing: Problems

Systematically solving CSPs: Summary

• Apply Depth-First Search with Pruning

OR

- Build Constraint Network
- Apply Arc Consistency
 - At least one domain is empty \rightarrow **no solution**
 - Each domain has a single value \rightarrow **one solution**
 - Some domains have more than one value \rightarrow ??? SO
- Search by Domain Splitting
 - Split the problem into a number of disjoint cases
 - Apply Arc Consistency to each case
 - Repeat as needed

Domain Splitting in Action:



Learning Goals for this class

You can:

- Implement local search for a CSP.
 - Implement different ways to generate neighbors
 - Implement scoring functions to solve a CSP by local search through either greedy descent or hill-climbing.

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Local Search motivation: Scale

- Many CSPs (scheduling, DNA computing, etc.) are simply too big for systematic approaches
- If you have 10^5 vars with dom(var_i) = 10^4
 - Systematic Search
 Arc Consistency



- $O(b^m)$ so $(10^4)^{(10^5)}$ A. $10^5 * 10^4$ B. $10^{10} * 10^8$ C. $10^{10} * 10^{12}$
- but if solutions are densely distributed......

Local Search: General Method

Remember, for CSP a solution is a possible world, **not** a path

- Start from a possible world
- Generate some neighbors ("similar" possible worlds)
- Move from the current node to a neighbor, selected according to a particular strategy

Local Search: Selecting Neighbors

How do we determine the neighbors?

- Usually this is simple: some small incremental change to the variable assignment
 - a) assignments that differ in one variable's value, by (for instance) a value difference of +1
 - b) assignments that differ (by any amount) in one variable's value
 - C) assignments that differ in two variables' values, etc.

Iterative Best Improvement

- How to determine the neighbor node to be selected?
- Iterative Best Improvement:
 - select the neighbor that optimizes some evaluation function
- Which strategy would make sense? Select neighbor with ... (choose the **best** answer, that is applicable in every situation)
- A. Maximal number of constraint violations
- B. Similar number of constraint violations as current state
- C. No constraint violations
- D. Minimal number of constraint violations
- E. Vogon poetry (the 3rd worst in the universe)

Selecting the best neighbor

- Example: A,B,C same domain {1,2,3}, (A=B, A>1, C≠3)
- Suppose we start with {A=1, B=1, C=1}
 - What are the neighbors? (depends on neighbor function)
 - Which neighbor is the best?

Example: N-Queens

 Put n queens on an n × n board with no two queens on the same row, column, or diagonal (i.e attacking each other)

 Positions a queen can attack



Example: N-queen as a local search problem CSP: N-queen CSP

- One variable per column; domains {1,...,N} => row where the queen in the ith column sits;
- Constraints: no two queens in the same row, column or diagonal

Neighbour relation: value of a single column differs Scoring function: number of attacks



How many neighbors ? A. 100

- B. 80
- C. 56
- D. 8 E. 42



Example: Local Search for N-Queens

For each column, assign randomly each queen to a row

(a number between 1 and N)

Repeat

- For each column & each number: Evaluate how many constraint violations changing the assignment would yield
- Choose the column and number that leads to the fewest violated constraints; change the assignment

Until solved



n-queens, Why?

Why this problem?

Lots of research in the 90's on local search for CSP was generated by the observation that the run-time of local search on n-queens problems is essentially independent of problem size!

Given random initial state, can solve *n*-queens in almost constant time for arbitrary n with high probability (e.g., n = 10,000,000)

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Constrained Optimization Problems

- So far we have assumed that we just want to find a possible world that satisfies all the constraints.
- But sometimes solutions may have different values / costs
- We want to find the optimal solution that
 - maximizes the value or
 - minimizes the cost

Hill Climbing means selecting the neighbor which best improves a (value-based) scoring function.

Greedy Descent means selecting the neighbor which minimizes a (cost-based) scoring function.

Hill Climbing

NOTE: Everything that will be said for Hill Climbing is also true for Greedy Descent

evaluation



Slide 20

Constrained Optimization Example

Example: A,B,C same domain {1,2,3}, (A=B, A>1, C≠3)

• Value = (C+A) so we want a solution that maximizes that

The scoring function we'd like to maximize might be: f(n) = (C + A) + #-of-satisfied-const

If we're doing Greedy Descent, then what we want to minimize is cost + #-of-conflicts

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Problems with Hill Climbing

- Local Maxima
- Plateaus and Shoulders



More problems in higher dimensions

E.g., Ridges – sequence of local maxima not directly connected to each other From each local maximum you can only go downhill



Corresponding problem for GreedyDescent Local minimum example: 8-queens problem



A local minimum with h = 1(all neighbors have h > 1)

Local Search: Summary

- A useful method for large CSPs
 - Start from a possible world (often randomly chosen)
 - Generate some neighbors ("similar" possible worlds)
 - Move from current node to a neighbor, selected to minimize/maximize a scoring function which combines:

 Information about how many constraints are violated
 Information about the cost/quality of the solution (you want the best solution, not just a solution)

Next Class

 How to address problems with Greedy Descent / Hill Climbing?

Stochastic Local Search (SLS)

