

Sept 15, 2017

①

- LP without LP
- Branch & Bound

TODAY

① LP w/o LP



Use that fact
that simplex
method work...

② Branch & Bound

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Summarized last time's example

$$(y_1) \text{ same} \rightarrow x_1 + 2x_2 \leq 8$$

$$(y_2) \text{ same} \rightarrow x_1 + x_2 \leq 5$$

$$(y_3) \text{ same} \rightarrow 2x_1 + x_2 \leq 8$$

$$\max (4x_1 + 5x_2) = z = z(\vec{x}) \quad (2)$$
$$x_1, x_2 \geq 0$$

add to get upper bound on $4x_1 + 5x_2$

$$(1) \quad 5(x_1 + x_2 \leq 5) \Rightarrow z(\vec{x}) \leq 25$$
$$4x_1 + 5x_2 \leq 5x_1 + 5x_2 \leq 25$$

(2) 3 (first ineq) + 1 (second) \rightarrow better

$$y_1 (x_1 + 2x_2 \leq 8)$$

$$y_2 (x_1 + x_2 \leq 5)$$

$$y_3 (2x_1 + x_2 \leq 8)$$

$$(y_1, y_2, y_3 \geq 0)$$

$$(y_1 + y_2 + 2y_3)x_1 + (2y_1 + y_2 + y_3)x_2 \leq 8y_1 + 5y_2 + 8y_3$$

So if $\underbrace{\hspace{10em}}$ at least 4 $\underbrace{\hspace{10em}}$ at least 5 $\underbrace{\hspace{10em}}$ bound on LP

Dual LP: minimize $8y_1 + 5y_2 + 8y_3$

$$\text{sit. } y_1 + y_2 + 2y_3 \geq 4, \quad 2y_1 + y_2 + y_3 \geq 5$$

$$\text{row} \Rightarrow A\vec{x} \leq \vec{b}, \quad \vec{x} \geq 0$$

$$[y_1 \ y_2 \ y_3] (A\vec{x} \leq \vec{b})$$

$$\vec{y}^T (A\vec{x} \leq \vec{b}) \quad (\text{if } \vec{y} \geq 0)$$

$$\Rightarrow (\vec{y}^T A) \vec{x} \leq \vec{y}^T \vec{b}$$

3

4

$$\begin{aligned} \max & 4x_1 + 5x_2 \quad \text{s.t.} \\ & x_1 + 2x_2 \leq 8 \\ & x_1 + x_2 \leq 6 \quad (\text{more coffee :}) \\ & 2x_1 + x_2 \leq 8 \end{aligned}$$

but x_1, x_2 have to be integers...

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Branch & bound:

$$x_1 = 0, 1, \dots, 6$$

$$x_1 + x_2 \leq 6 \text{ means } x_2 = 0, 1, \dots, 6$$

could try out $7 \cdot 7 = 49$ possibilities

NetPlex

