

Math 340 (102) Midterm Brief Solutions

Fall 1999

$$1. \quad x_3 = 6 - 3x_1 - 2x_2$$

$$x_4 = -4 + x_1 + x_2$$

$$z = x_1 + 2x_2$$

Infeasible, so add x_0 's

$$x_3 = 6 - 3x_1 - 2x_2 + x_0$$

$$x_4 = -4 + x_1 + x_2 + x_0$$

$$\omega = -x_0$$

x_0 enters, x_4 leaves

$$x_0 = 4 - x_1 - x_2 + x_4$$

$$x_3 = 10 - 4x_1 - 3x_2 + x_4$$

$$\omega = -4 + x_1 + x_2 - x_4$$

x_1 enters, x_3 leaves

$$x_1 = \frac{5}{2} - \frac{1}{4}x_3 - \frac{3}{4}x_2 + \frac{1}{4}x_4$$

$$x_0 = \frac{3}{2} + \frac{1}{4}x_3 - \frac{1}{4}x_2 + \frac{3}{4}x_4$$

ω = minus last line

x_2 enters, x_1 leaves

$$x_2 = \frac{10}{3} - \frac{1}{3}x_3 - \frac{4}{3}x_1 + \frac{1}{3}x_4$$

$$x_0 = \frac{2}{3} + \frac{1}{3}x_3 + \frac{1}{3}x_2 + \frac{2}{3}x_4$$

$$\omega = -\frac{2}{3} - \frac{1}{3}x_3 - \frac{1}{3}x_2 - \frac{2}{3}x_4$$

So optimal ω is $-3/2$ so

original problem is infeasible

Check: $\frac{1}{3}(x_3 \text{ const}) + \frac{2}{3}(x_4 \text{ constraint})$

$$\text{is } \frac{1}{3}(3x_1 + 2x_2 \leq 6)$$

$$+ \frac{2}{3}(-x_1 - x_2 \leq -4)$$

$$\underline{\frac{1}{3}x_1 \leq -\frac{2}{3}}$$

which is
infeasible

$$2. \quad x_1^* = 1 \rightarrow y_1^* \text{ must} = 0$$

$$x_2^* = 3 \rightarrow y_2^*$$

$$x_3^* = 0 \rightarrow y_3^*$$

$$x_4^* = 1 \rightarrow y_4^* \text{ must} = 0$$

$$x_5^* = 0 \rightarrow y_5^*$$

$$x_6^* = 0 \rightarrow y_6^* \text{ must} = 0$$

So original solution is feasible, and for dual:

$$0 = y_4^* = -7 + y_1^* + y_2^* + 3y_3^*$$

$$0 = y_5^* = -5 + y_1^* + 3y_2^* + y_3^*$$

and $y_1^* = 0$ so

$$y_2^* + 3y_3^* = 7 \quad \text{so } y_2^* = 1, y_3^* = 2$$

$$3y_2^* + y_3^* = 5$$

Check all y^* 's are ≥ 0 :

$$y_1^* = 0, y_2^* = 1, y_3^* = 2, y_4^* = y_5^* = 0,$$

$$y_6^* = -2 + y_1^* + y_2^* + y_3^* = 1 \text{ so OK.}$$

ALTERNATIVELY: Check dual feasibility from inequalities:

$$y_1 + y_2 + 3y_3 \geq 7$$

$$y_1 + 3y_2 + y_3 \geq 5$$

$$y_1 + y_2 + y_3 \geq 2$$

checks out with $(y_1, y_2, y_3) = (0, 1, 2)$

so OK.

Since derived dual solution is feasible, original solution is optimal.

$$3. \quad x_3 = 2 + \varepsilon - x_1 - x_2$$

$$x_4 = 2 + \varepsilon^2 - x_1 + 3x_2$$

$$z = 2x_1 + x_2$$

x_1 enters, x_4 leaves

$$x_1 = 2 + \varepsilon^2 - x_4 + 3x_2$$

$$x_3 = \varepsilon - \varepsilon^2 + x_4 - 4x_2$$

$$z = 4 + 2\varepsilon^2 - 2x_4 + 7x_2$$

x_2 enters, x_3 leaves

$$x_2 = \frac{\varepsilon - \varepsilon^2}{4} + \frac{1}{4}x_4 - \frac{1}{4}x_3$$

$$x_1 = \cancel{2 + \varepsilon^2} - 2 + \varepsilon^2 + \frac{3}{4}(\varepsilon - \varepsilon^2) + \text{whatever}$$

$$z = 4 + 2\varepsilon^2 + \frac{7}{4}(\varepsilon - \varepsilon^2) - \frac{1}{4}x_4 - \frac{7}{4}x_3$$

So optimal is $(x_1^*, x_2^*) = (2 + \varepsilon^2 + \frac{3}{4}(\varepsilon - \varepsilon^2), \frac{\varepsilon - \varepsilon^2}{4})$

with $z^* = 4 + \frac{7}{4}\varepsilon + \frac{1}{4}\varepsilon^2$. Letting $\varepsilon \rightarrow 0$ gives

$$(x_1^*, x_2^*) = (2, 0), \quad z^* = 4.$$

4. The constraint corresponds to a slack variable, say x_i .

If x_i is basic in final dictionary, there is a surplus of cocoa ingredients and such ingredients have no marginal value. So we won't buy.

If x_i is non-basic in final dictionary, look at its coefficient in the objective (i.e. z) row. If the coefficient is > 5 we buy, since its marginal value to us is $> \$5$ per kilo. If it is < 5 we don't buy.

If it = 5 we don't care.