

$$PSPACE = \bigcup_k SPACE(n^k)$$

$$NPSPACE = \bigcup_k NSPACE(n^k)$$

=
Savitch's Thm:

$$NSPACE(f(n)) \subseteq SPACE(f^2(n))$$

=
Pf: Let $L \in NSPACE(f(n))$

Claim
 $L \in TIME(2^{C f(n)})$

for some C.

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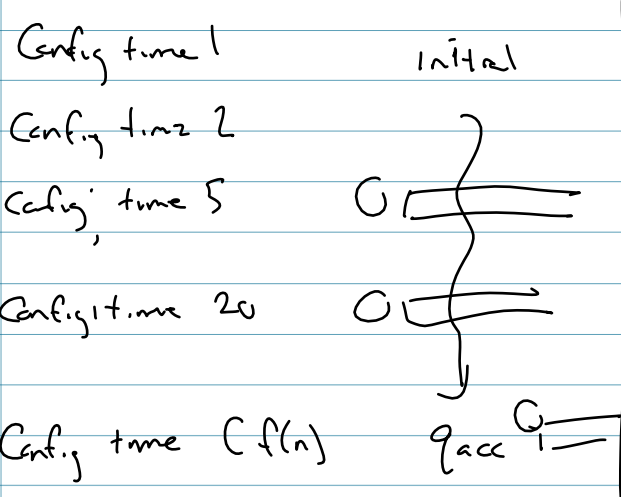
Ch 8: Space

8.1 Savitch's Theorem

[Allusions to 8.2, ...]

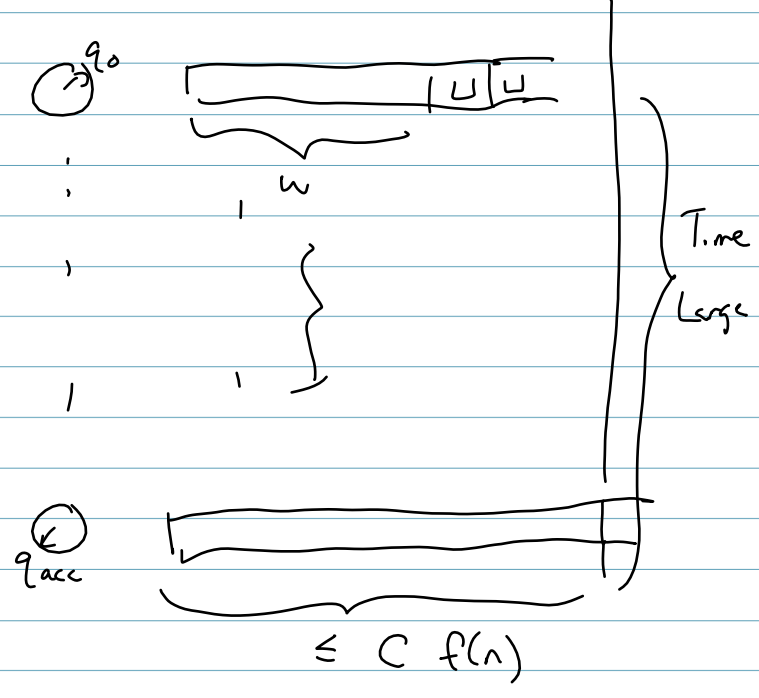
$$SPACE(f(n)) = \{ L \mid L \text{ can be decided in space } \leq C f(n) \}$$

$$NSPACE(\quad) = \{ \text{--- non-deterministic Turing machine ---} \}$$



If config time 5 =
... = 20
we can delete part between
config 5 to config 20

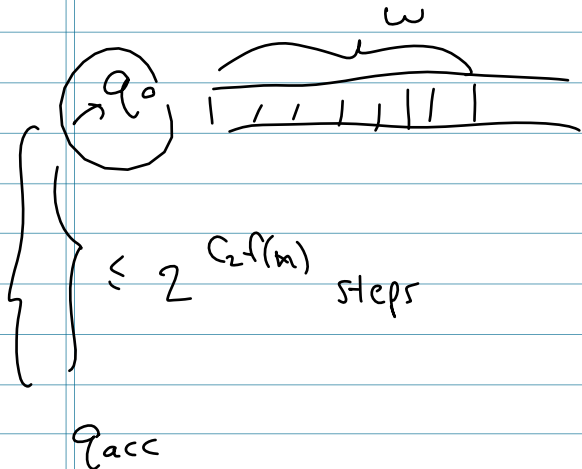
We know \exists non-det T.M., M ,
such that



$\leq C_1 f(n)$ or $2^{C_2 f(n)}$

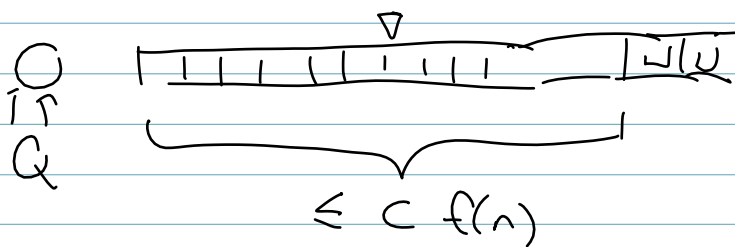
Now say # config in a computation reaching q_{acc}

$\leq 2^{C_2 f(n)}$



If w can be decided and reaches q_{acc} , it can be done without reaching the same configuration twice.

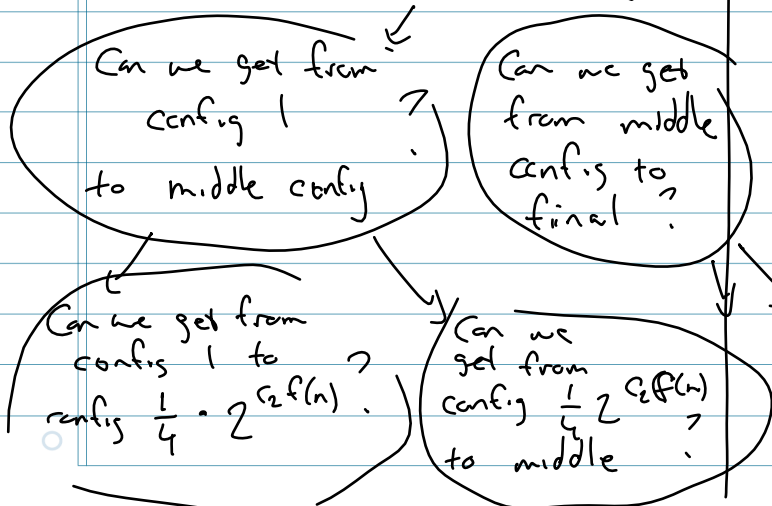
How many config?



$\leq |Q| \times |\Gamma|^{C f(n)} \times C f(n)$
 states what written tape head

First! What is config 1?
 What is config 2?
 write in space $C f_2(n)$

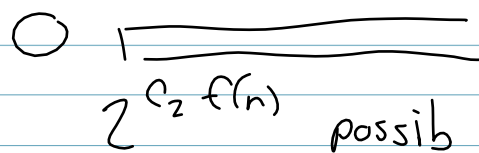
Second! What is in the middle?



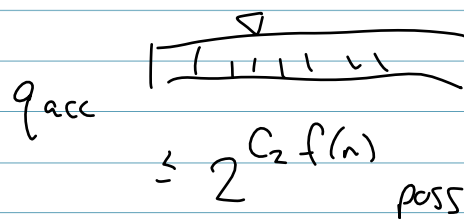
Algorithm:

initial only one possibility

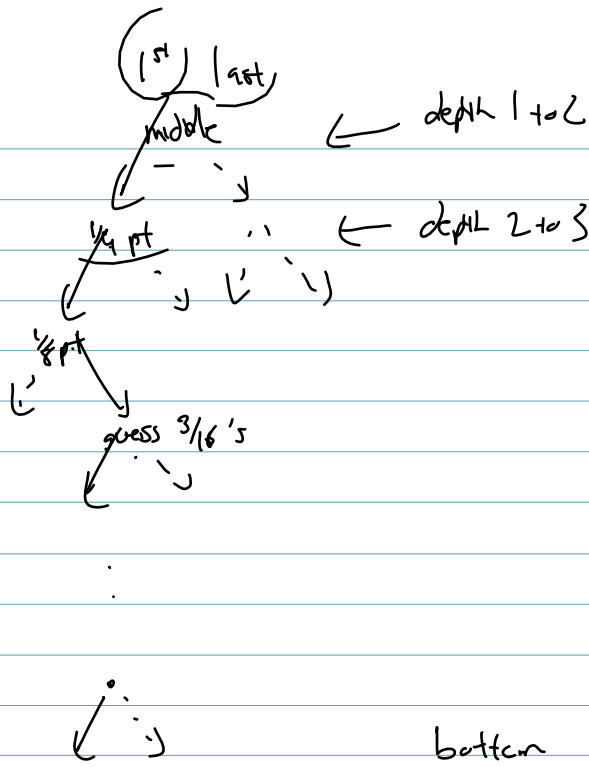
middle config



final:



recurse



depth d such that

$$\frac{1}{2^d} (\text{time}) \leq 1$$

Need to remember where you are in tree:

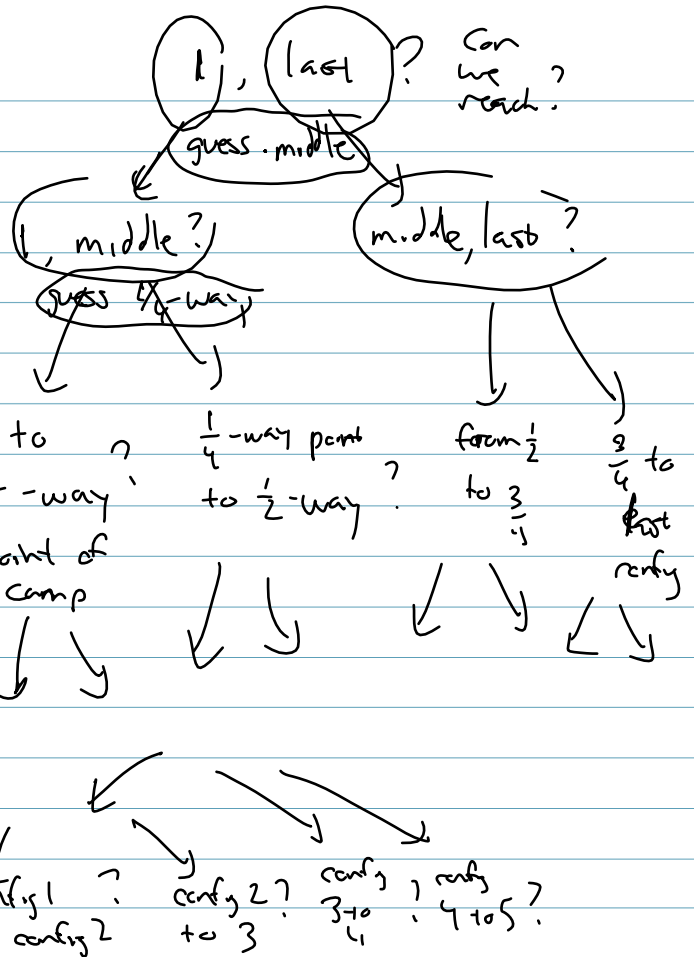
L/R, L/R, ..., L/R
 └───┬───┘
 d

Space: $\text{order}(d)$

Total Space:

$(d \text{ or } d+1) (C_2 f(n) + \text{small})$
 configs per config

+ $\text{order}(d)$ where we are in tree $\leq \text{order}(f(n)^2)$



If w accepted, it is accepted by a path of time $\leq 2^{C_2 f(n)}$

So depth need:

$$\frac{1}{2^d} 2^{C_2 f(n)} \leq 1$$

i.e. $d = C_2 f(n)$

To write down:

1st config, last, middle, any

config: space $C_2 f(n)$ per config