## Question 1

First suppose $S \in S$. Then by the definition of $S$, since $S \in S$ we have that $S \notin S$. This is a clear contradiction.

Now suppose $S \notin S$. By the definition of $S$ then, $S$ belongs in $S$, contradicting that $S$ is not in $S$.

I think no "reasonable" set theory can allow $S$ to exist, since no matter whether $S$ contains itself or not there is some logical inconsitency.

## Question 2

Let $S=\{0,1\}$. Note the DFAs are included at the end of the document.
(b) Regular expressions describing $\{w \mid w$ containing three 1's $\}$.

$$
S * 1 S * 1 S * 1 S * \text {, or } 0 * 10 * 10 * 1 S * \text {, or } S * 10 * 10 * 1 S * \text {. }
$$

One incorrect answer which was submitted is $S * 111 S$. This requires three 1 s to be consecutive, and thus is too restrictive.
(f) Regular expressions describing $\{w \mid w$ doesnt contain the substring 110$\}$.
$0 * \cup[(0 * 10) * 1 *$, or $0 *(100 *) * 1 *$, or $0 *(10 \cup 0) * 1 *$, or $(10 \cup 0) * 1 *$. Incorrect answers that were submitted were: $0 *(10 *) * 1 *$ : note that the middle $(10 *) *$ matches 1110, for example. $(0 * 1) * 0 * 1 *$ : note that the string 110 matches this expression.
(l) Regular expression describing $\{w \mid w$ contains an even number of 0 s, or exactly two 1 's $\}$. $(1 * 01 * 0) * 1 * \cup 0 * 10 * 10 *$

## Question 3

Let $R^{\prime}$ be the regular expression of problem 3. It is important to include "both directions" of the explanation that $L\left(R^{\prime}\right)$ describes the set of strings with an even number of 0 s and an odd number of 1 s . (Several solutions only presented the direction that shows $L\left(R^{\prime}\right)$ is a subset of $L$, where $L$ is language of string with an odd number of 1 s and an even number of 0 s , but omitted to show that $L$ is a subset of $L\left(R^{\prime}\right)$.
$L\left(R^{\prime}\right)$ is a subset of $L$ : Suppose that $w$ is in $L\left(R^{\prime}\right)$. We show that $w$ must have an even number of 0 s and an odd number of 1 s . Note that $w$ must be the concatenation of three strings, say $w=x y z$, where both $x$ and $z$ are in $L(R)$ and $y$ is in $L(1 \cup 01(11) * 0)$. Hence, both $x$ and $z$ must have an even number of 0 s and $y$ must have either zero or two 0 s , (depending on whether $y$ is in $L(1)$ or $y$ is in $L(01(11) * 0)$ ). Hence since all of $x, y$, and $z$ have an even number of 0 s , so must $w$. Also, both $x$ and $z$ must have an even number of 1 s , but $y$ must have an odd number of 1 s . Since two even numbers plus one odd number is an odd number, clearly $w$ must have an odd number of 1 s .
$L$ is a subset of $L\left(R^{\prime}\right)$ : (This is the harder direction, kudos to those of you who understood how to proceed with this one.) Suppose that $w$ has an even number of 0s and an odd number of 1 s . We show that $w$ is in $L\left(R^{\prime}\right)$. Let $x$ be the longest prefix of $w$ that has
an even number of 0 s and an even number of 1 s . Note that $x$ may be the empty string but $x$ cannot be the whole string $w$; that is, $x$ is a proper prefix of $w$. We now consider two cases.

1. The first case is that $x 1$ is a prefix of $w$. Then $w=x 1 z$ for some $z$. In this case, since $x$ has an even number of both 0 s and 1 s , so must $z$. Hence $w$ is in $L(R 1 R)$ and therefore in $L(R(1 \cup 01(11) * 0) R)$.
2. The second case is that $x 0$ is a prefix of $w$. Now, $x 00$ cannot be a prefix of $w$, since $x 00$ has an even number of both 0 s and 1 s , but we know that $x$ is the longest prefix of $w$ with an even number of both 0 s and 1 s . Therefore, $x 01$ must be a prefix of $w$. The string $x 01$ has an odd number of 0 s and an odd number of 1 s ; hence another 0 must occur in $w$ after the prefix $x 01$. Therefore, $w=x y z$ where $y$ is of the form $011 * 0$. But in fact, $y$ must contain an odd number of 1 s : if this were not true then $x y$ would be a prefix of $w$ containing an even number of 0 s and 1 s and would be longer than $x$, but we know that $x$ is the longest prefix of $w$ containing an even number of 1 s . Thus, $y$ is a string with a 0 at each end and an odd number of 1 s between these two 0 s, which means that $y$ must be in $L(01(11) * 0)$. Therefore, $y$ is also in $L(1 \cup 01(11) * 0)$. We have now shown that $w=x y z$ where $x$ has an even number of both 0 s and 1 s and $y$ is in $L(01(11) * 0)$. Since the string $x y$ has an even number of 0 s and an odd number of 1 s , it must be that $z$ has both an even number of 0 s and an even number of 1 s . Thus, $w=x y z$ where $x$ is in $L(R), y$ is in $L(1 \cup 01(11) * 1)$ and $z$ is in $L(R)$. Therefore, $w=x y z$ is in $L(R(1 \cup 01(11) * 1) R)$ and we are done.


Figure 1: DFA for problem 2.b.


Figure 2: DFA for this problem 2.f.


Figure 3: DFA for problem 2.l.

