

CPSC 421/501

Oct 10, 2025

- NFAs $\rightarrow$ DFAs (Oct 8)

- Regular Expressions:

[Sip]: $\emptyset, \varepsilon$, elements of Σ
plus $\cup, \circ, *$

[Could add: $+, \neg, \wedge, \dots$]

- Regular expressions describe regular languages (with proof)

- Vice versa: a cautionary tale...

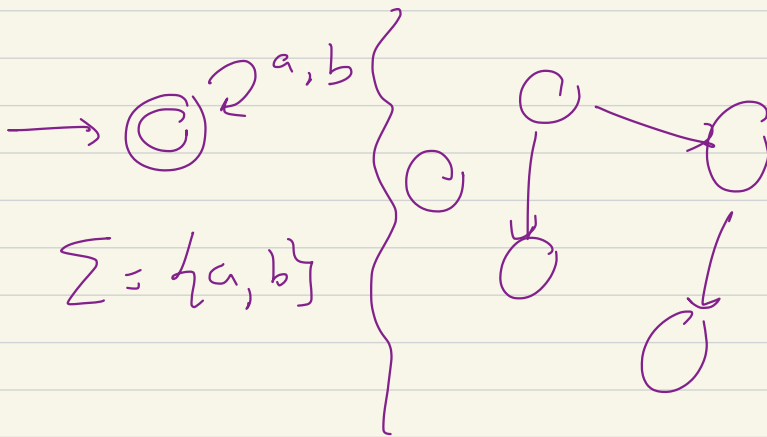
$\rightarrow$ in §1.3 [Sip]

Say you have an NFA with

$M = (Q, \Sigma, \delta, q_0, F)$ and

$|Q| = k$.

How large — how many states
could an equivalent DFA have?



at most?

not needed

$$\Sigma = \{a, b\}$$

$$(\Sigma^*) \circ (\{a\}) \circ \Sigma^4$$

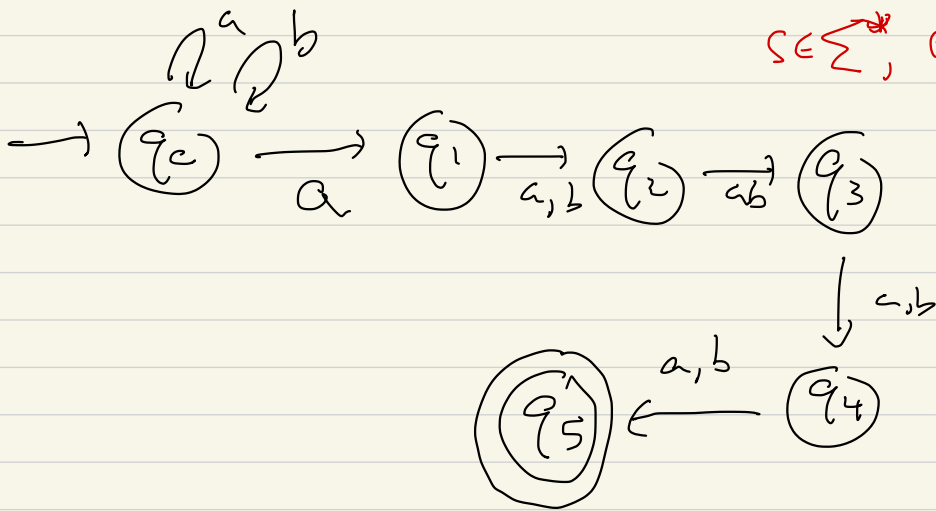
$$\Sigma^* a \Sigma^4 \leftarrow \text{regular exp.}$$

$$L = \{ s \mid \text{5th to last letter/symbol is } a \}$$

NFA

$$= \{ s a \sigma_1 \sigma_2 \sigma_3 \sigma_4 \mid$$

$$s \in \Sigma^*, \sigma_1, \dots, \sigma_4 \in \Sigma \}$$



blah
 blah $a \sigma_1 \sigma_2 \sigma_3 \sigma_4$
 blah

How large is a min size DFA

$$\text{for } L = \sum^* a \sum^4 = \sum^* a \circ \sum \circ \sum \circ \sum \circ \sum$$

idea

32 {

a a a a a
a a a a b
a a a b a
:
b b b b b

if you have
state for each
of these

i

31 states, then 2 of these have to
wind up in the same state

NFA $\rightarrow$ 6 states

DFA $\rightarrow 2^5$ states

for any $L \subseteq \Sigma^*$

$$L^{\text{rev}} = \{ s^{\text{rev}} \mid s \in L \}$$

and

s^{rev} ! means

$$(\sigma_1 \dots \sigma_n)^{\text{rev}} = \sigma_n \sigma_{n-1} \dots \sigma_1$$

$$(aab)^{\text{rev}} = baa$$

$$(abba)^{\text{rev}} = abba \text{ (palindrome)}$$

$$(\Sigma^* a \Sigma^*)^{\text{rev}} = (\Sigma^*)^{\text{rev}} (a)^{\text{rev}} (\Sigma^*)^{\text{rev}}$$

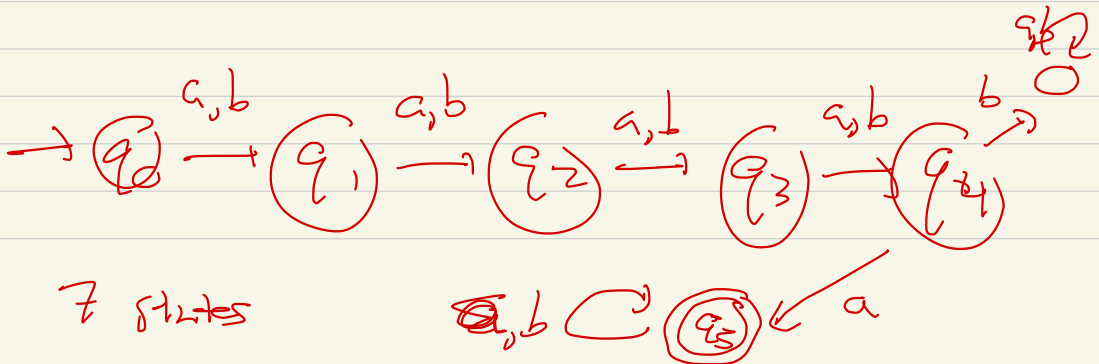
$$(L_1 L_2 L_3)^{\text{rev}} = L_3^{\text{rev}} L_2^{\text{rev}} L_1^{\text{rev}}$$

$$(L_1 \circ L_2 \circ L_3)^{\text{rev}} = L_3^{\text{rev}} \circ L_2^{\text{rev}} \circ L_1^{\text{rev}}$$

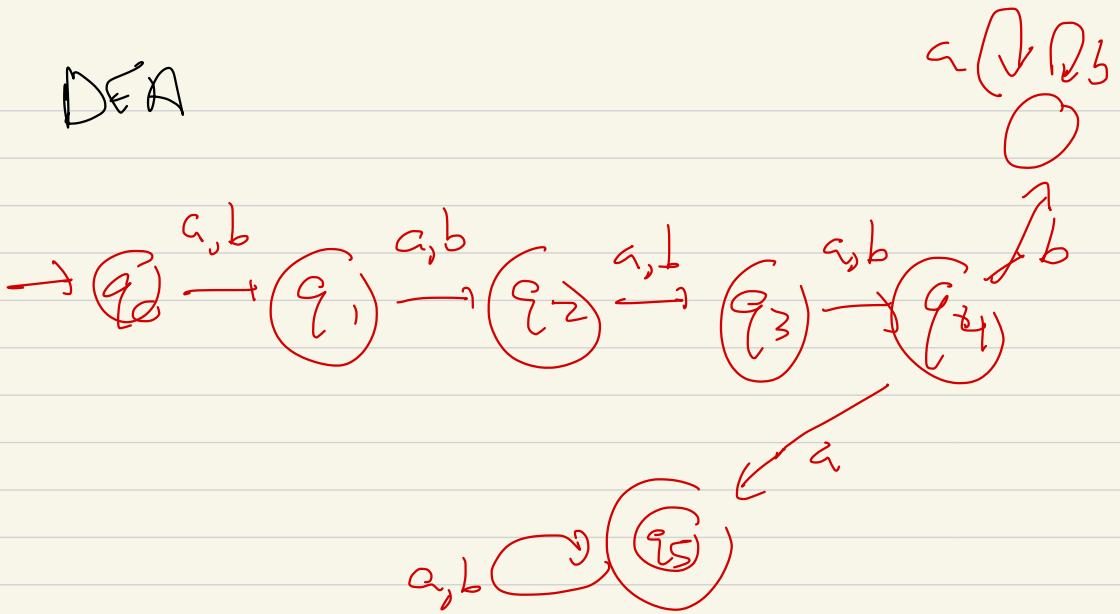
Example

$$L = \{ s \mid \text{5th to last letter/symbol is an "a"} \}$$

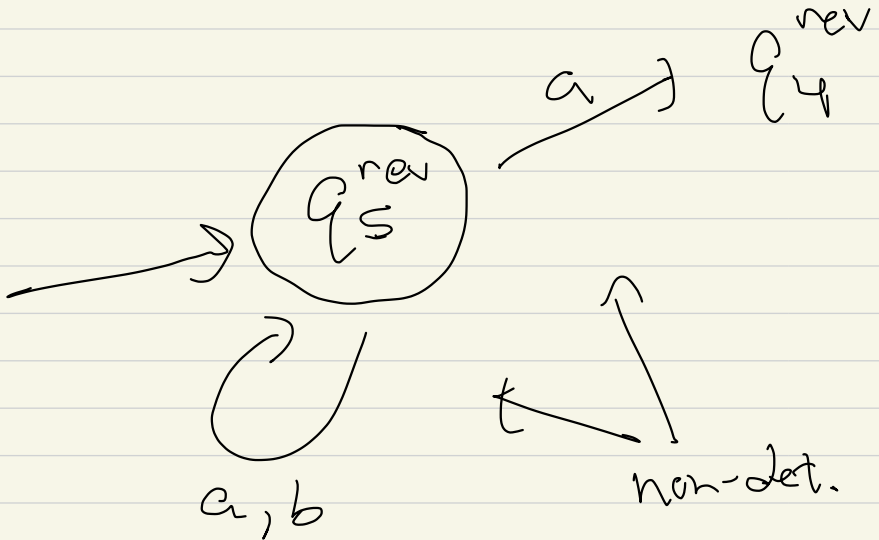
$$L^{\text{rev}} = \{ s \mid \text{5th (from first) letter/symbol is an "a"} \}$$



DEFA



rev



(NFA)

Regular Expression!

$\Sigma \leftarrow \text{alphabet}$

Rules:

Primitives: (1) $\emptyset$, (2) $\{\sigma\}$ for any $\sigma \in \Sigma$
 σ, σ
(3) $\{\epsilon\}$ or ϵ

Building Rules: if R_1, R_2 regular expressions, then so are

(4) $(R_1) \cup (R_2)$ (5) $(R_1) \cap (R_2)$ (6) $(R_1)^*$

$$\Sigma = \{a, b\}$$

e.g.,

$$\left(\{a\} \cup \{b\} \right)^* = \{a\} \circ$$

$$\left(\{a\} \cup \{b\} \right) \circ \left(\{a\} \cup \{b\} \right) \circ$$

$$\circ \left(\{a\} \cup \{b\} \right) \circ \left(\{a\} \cup \{b\} \right)$$

$$\{a\} \cup \{b\} \xleftarrow{\text{write}} \Sigma, L \circ L \circ L \circ L = L^4 \xrightarrow{\text{write}}$$

$$\Sigma^* \circ \{a\} \circ \Sigma \circ \Sigma \circ \Sigma \circ \Sigma$$

$$\Sigma^* \circ \{a\} \circ \Sigma^4 = \Sigma^* \circ (a) \circ \Sigma^4$$

$$= \Sigma^* a \Sigma^4$$