

CPSC 421/501

Oct. 8, 2025

- NFAs $\rightarrow$ DFAs

- Regular Expressions:

[Sip]: $\emptyset, \varepsilon$, elements of Σ

plus $\cup, \circ, *$

[Could add: $+$, $\neg$, $\wedge$, ...]

- Regular expressions describe regular languages (with proof)

- Vice versa: a cautionary tale...

Idea: L_1 regular DFA M_1
 L_2 " " M_2

$L_1 \circ L_2$ is regular

View M_1 as NFA
" M_2 " " "

Trick (uses NFA)

- build DFA for $L_1 \circ L_2$

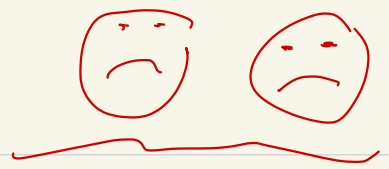
- NFA implementation of $M_1 \circ M_2$

Caution: # states in " $M_1 \circ M_2$ "

= (# states in M_1) + (# states in M_2)

NFA $\rightarrow$ DFA

~~~~~ # states in resulting DFA 



Trick:

$$NFA = (Q, \dots)$$



$$DFA = (Power(Q), \dots)$$

Idea:

State set of DFA is

$Power(Q)$ , intuition is

$$t \in Power(Q), \quad t \subseteq Q$$

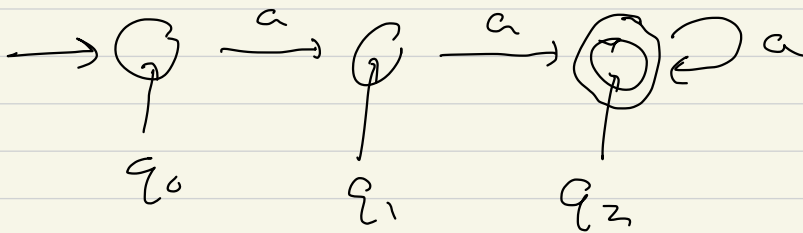
In state  $t$  of the DFA means

in the NFA we could get to all elements of  $t$

Rem:  $|Power(Q)| = 2^{|Q|}$

=

Say NEA:



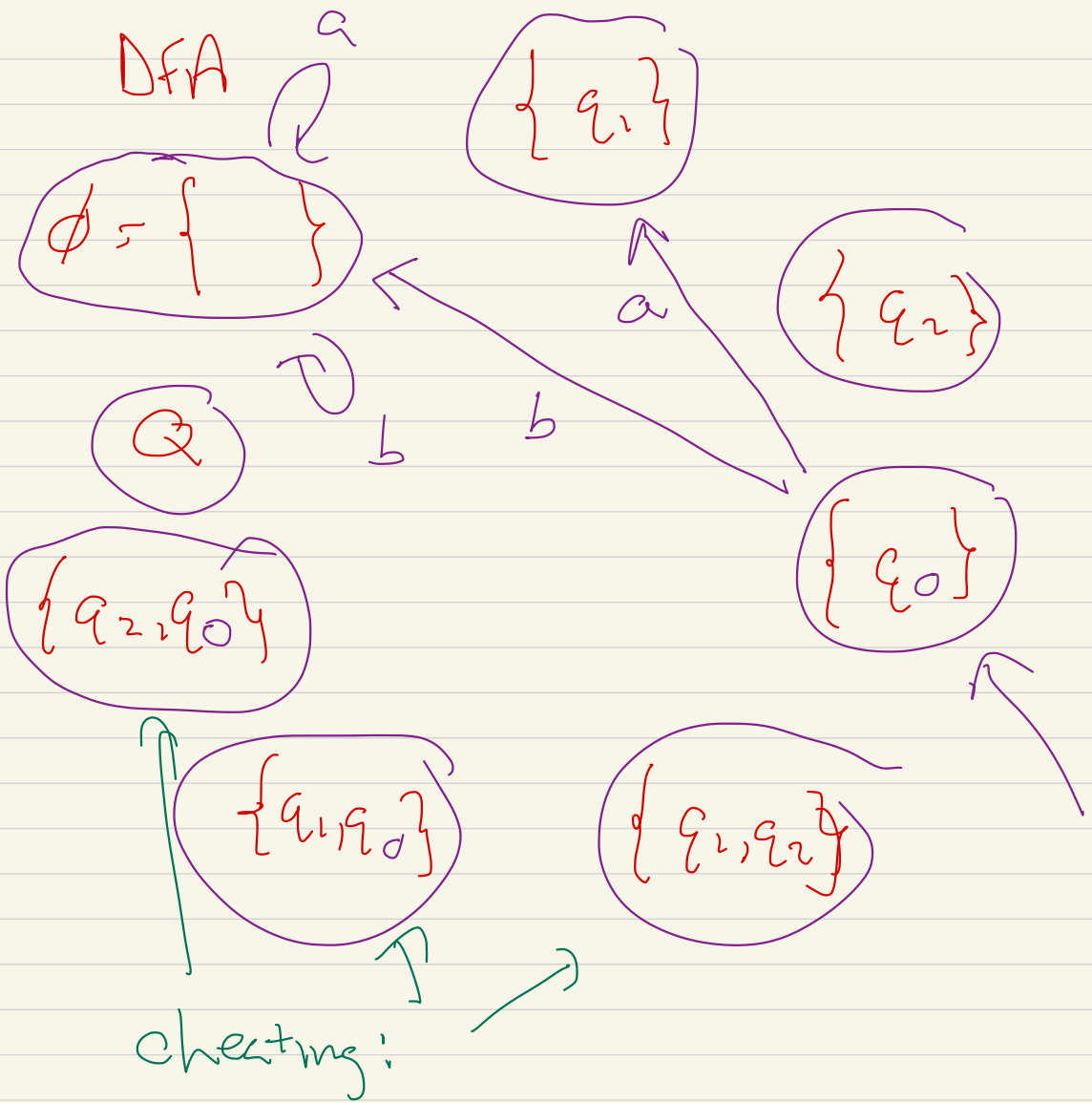
← NEA

$$\Sigma = \{a, b\}$$

In NEA

$aabbaa$   
 $\uparrow$   
 $q_0$

Initial, state before  
reading anything, is  $\{q_0\}$



Initial state:  $\{q_0\}$

$\delta(\{q_0\}, b) =$  set of all states we can reach:  $\emptyset$

$$\delta(\emptyset, \sigma) = \emptyset$$

any  
 $\sigma \in \Sigma^*$

if we can't  
 be in any  
 state

input ab

$\{q_0\}$

$\downarrow a$

$\{q_1\}$

$\xrightarrow{a}$

no  $\downarrow b$  arrows  
 $\emptyset$

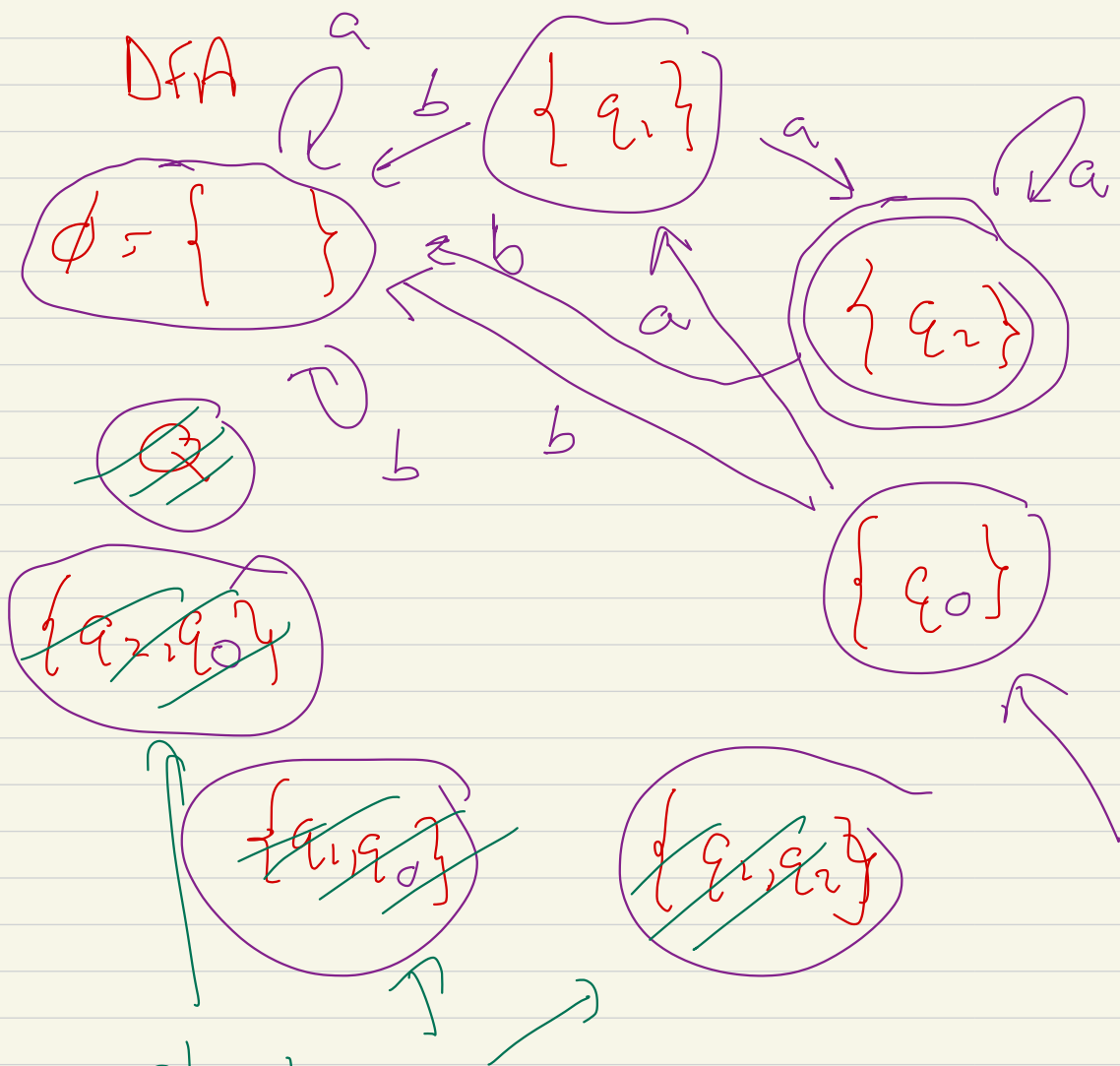
$\emptyset \in$

$\text{Power}(\{q_0, q_1, q_2\})$

$\emptyset \subset \{q_0, q_1, q_2\}$

$\{ \} \subset \{ \downarrow b \downarrow \}$

DFA



checking:

we don't really use

2-element subsets  
3-element subset

Warning:

$\{q_1, q_2\}$  would be  
accepting,  
if we needed it

$$t \subseteq Q, \quad t \in \text{Power}(Q)$$

$t$  is accepting iff  $t$  contains  
an accepting state  
of the NFA

$$\text{NFA} = (Q, \Sigma, \delta, q_0, F)$$

$$\delta: Q \times \Sigma_{\varepsilon} \longrightarrow Q$$

$$\underbrace{\Sigma}_{\Sigma \cup \{\varepsilon\}}$$

Let

$$L = \begin{cases} \emptyset & \text{if } P = NP \\ \Sigma^* & \text{if } P \neq NP \end{cases}$$

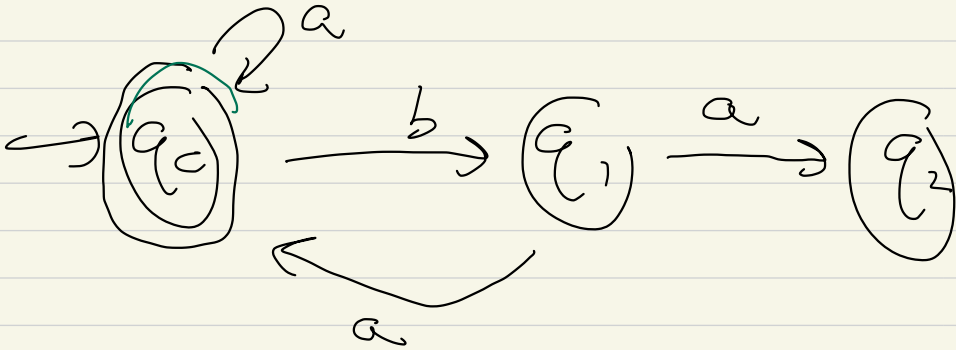
$$\emptyset : \rightarrow \textcircled{q_0} \rightarrow \Sigma$$

$$\Sigma^* : \rightarrow \textcircled{\textcircled{q_0}} \rightarrow \Sigma$$

then  $L$  has a 1-state DFA  
that recognizes it.

How to build it --- ???

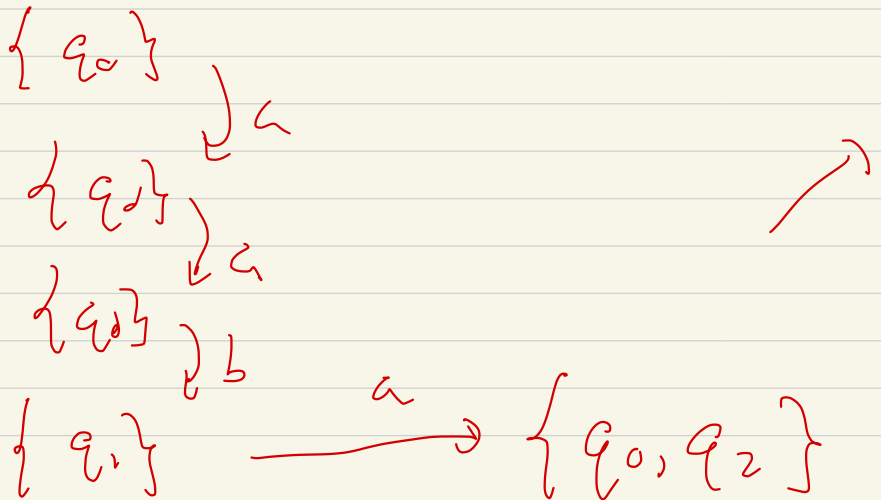
NFA:



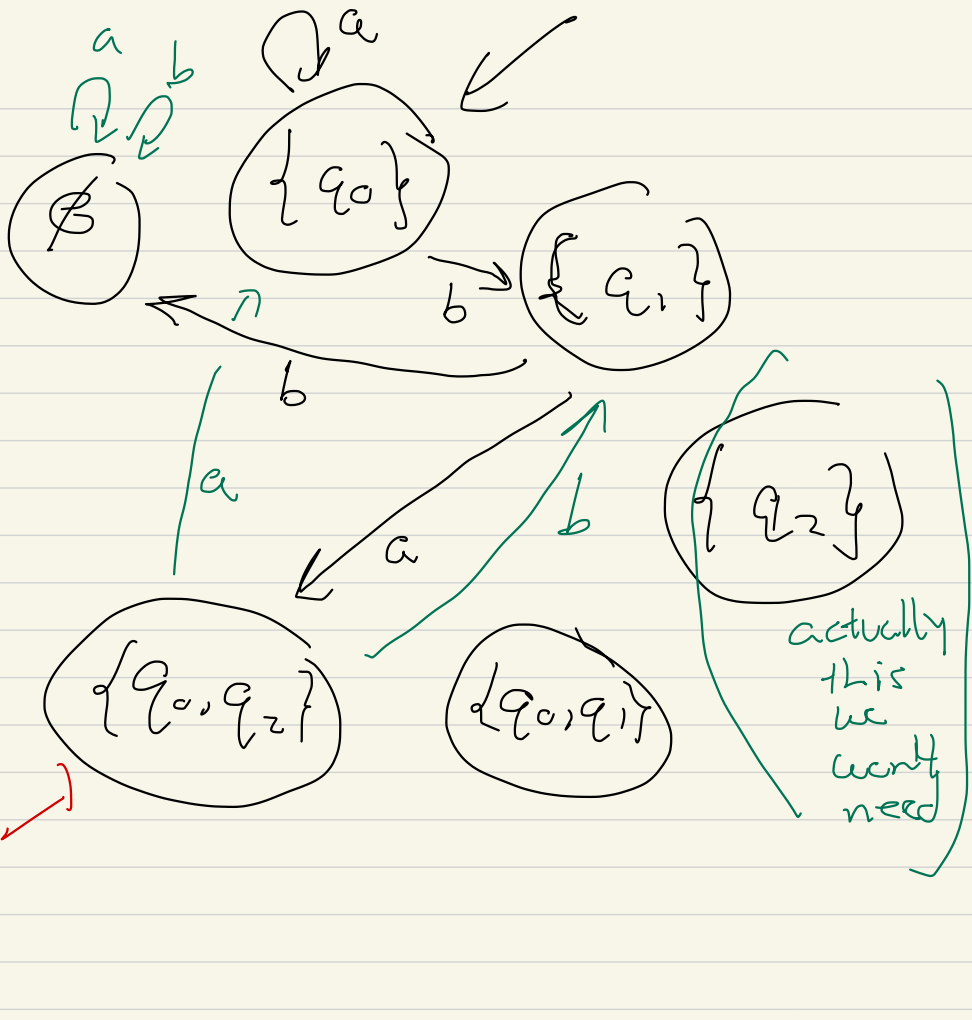
$(Q, \Sigma, \delta, q_0, F)$

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Sample input aaba---



DEFA



$$( \text{Power}(Q), \Sigma, \delta, \{q_0\}, \hat{F} )$$

$$\hat{F} = ( \{q_0\}, \{q_0, q_2\} )$$