

Last time:

A language $L \subseteq \text{ASCII}^*$ is

recognizable if

$$L \in \text{Image}(\text{LanguageRecBy}),$$

i.e.

$$\exists p \in \text{ASCII}^* \text{ s.t.}$$

$$L = \{ i \in \text{ASCII}^* \mid p \text{ accepts } i \}$$

return("yes")
↓

ACCEPTANCE

$$= \{ (p, i) \mid p \text{ accepts } i \}$$

is recognizable

$$T = \{p \mid p \notin \text{LanguageRecBy}(p)\}$$

is not recognizable.

So what?

Problem 3 on Homework 2

replaced with a bonus question,

due Monday, Sept 22, 11:59 pm

Homework 3 still due on

Thursday, Sept. 25, 11:59 pm

Technicalities:

- $\text{LanguageRecBy}(p)$ if p is not a valid python program --- ?

ACCEPTANCE

$$= \{ s \in \sum_{\text{ASCII}}^* \mid s = p \sigma_0 i \}$$

s.t. p is a valid Python program
accepting i }

REJECTION

LOOPING

NON-PYTHON-INPUT

LanguageRecBy(p)

defined $\left\{ \begin{array}{l} i \in \text{ASCII}^* \text{ s.t.} \\ p \text{ accepts } i \end{array} \right\}$

p returns "yes" on input i , viewing
 p as a Python program and
 i as an input to Python program

What if p is not a valid

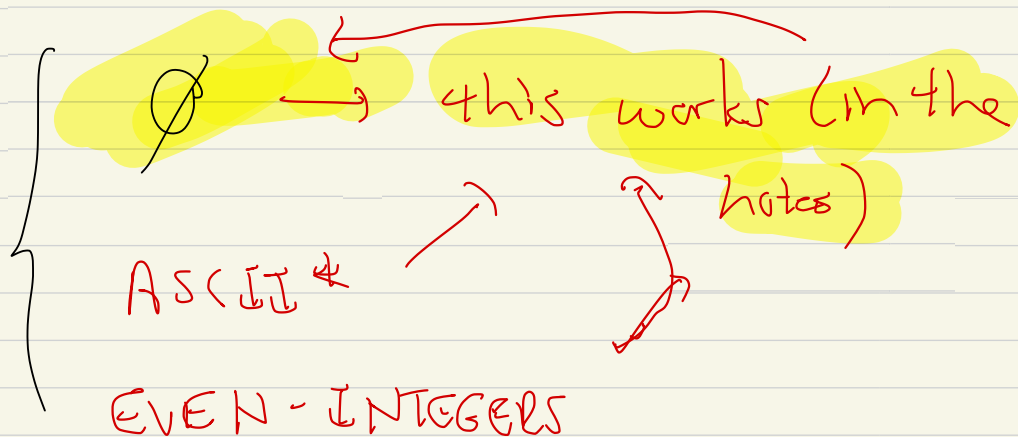
Python program? $p = "421"$?

$\text{LanguageRecBy}(p)$

where $p \in \text{ASCII}^*$ but p is
not a valid Python program?

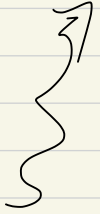
In this we define

$\text{LanguageRecBy}(p)$ to be



Warning! You don't want this to be
something you can't "compute"

VALID-PYTHON-PROGRAMS



$\subset \text{ASCII}^*$

we've defined

LangRecBy

LanguageRecBy:

$\text{ASCII}^* \rightarrow \text{Power}(\text{ASCII}^*)$

What happens when you feed
your program (p, i) ,

$$p \in \text{ASCII}^*$$

$$i \in \text{ASCII}^*$$

Let's say that $\sigma \in \text{ASCII}$
that isn't allowed anywhere is a
Python program

$$p \sigma i \text{ s.t. } p \in \text{VALID-PYTHON-PROG}$$

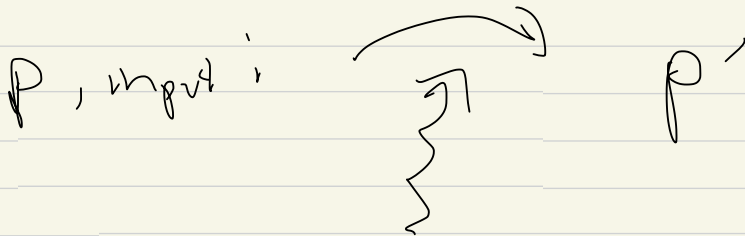
then $p \sigma i$ can be "parsed
unambiguously" as p , then i

Remark: $\forall p \in \text{VALID-PYTHON-PROG}$,
and input $i \in \text{ASCII}^*$

$\exists p' \in \text{VALID-PYTHON-PROG}$ s.t.,

① p on input i
② p' } "do the same thing"

Going from



Can be done in a finite
amount of time

This is a comment

#

def isDivByThree():

print("Give me a positive integer.")

i = input("Your input: ")

while (i > 0) : i = i - 3

if (i == 0) :

return ("yes")

return ("no")

i = 4212025

ignoring the input

```
# This is a comment
```

```
#
```

```
def isDivByThree( ):
```

```
    print("Give me a positive integer.")
```

```
    i = input("Your input: ")
```

$i = 4212025$

```
    while ( i > 0 ) : i = i - 3
```

```
    if ( i == 0 ) :
```

```
        return ( "yes" )
```

```
    return ( "no" )
```

'Hardwiring $i = 4212025$ as the input to p ' Ignoring real input.

Prop 4.19: Let T be as above. Then

$$q \circ q \in \left\{ \begin{array}{c} \text{NON-PYTHON} \\ \vdots \end{array} \right\}$$

$$\Rightarrow q \left\{ \begin{array}{c} \in \\ \notin \end{array} \right\} T$$

VALID.PYTHON.PROGRAM $\subset$ ASCII^{*}

σ is not allowed to be in
any Python program

formally describe (p, i) :

$\mapsto p \sigma i$

ACCEPTANCE $\equiv$

$$\left\{ p \sigma i \mid \begin{array}{l} p \in \text{VALID-PROGRAMS} \\ p \text{ accepts } i \\ (p \text{ returns "yes"}) \end{array} \right\}$$

REJECTION

$$\left\{ p \sigma i \mid \begin{array}{l} p \in \text{VALID-PROGRAMS} \\ p \text{ rejects } i \\ (p \text{ returns "no"}) \end{array} \right\}$$

LOOPING

$$\left\{ p \sigma i \mid \begin{array}{l} \sim \sim \sim \\ p \text{ does not accept or reject } i \\ (p \text{ loops on } i) \end{array} \right\}$$

NOT-VALID-PYTHON-INPUT

$$\left\{ s \in \text{ASCII}^* \mid \text{st. } s \right.$$

is not in ACCEPTANCE

$\cup \text{REJECTION} \cup \text{LOOPING} \}$

HALTING $\exists$

ACCEPTANCE $\cup$ REJECTION

CLASS ENDS

$$\begin{array}{ccc} \text{ASCII}^* & \times & \text{ASCII}^* \rightarrow \text{ASCII}^* \\ p & & i \end{array}$$

$(p, i) \mapsto p'$ found
by finding
the input
statement
~ ~

$$\text{ASCII}^* \times \text{ASCII}^* \rightarrow \text{ASCII}^*$$

$$(p, i) \mapsto p \sigma_i$$

$$(p, i) \mapsto p' \quad \begin{array}{l} \text{on any} \\ \text{input} \\ \text{returns} \\ \text{what } p \\ \text{returns on } i \end{array}$$

Rem: M.V. points out: if

$\sigma \in i$, then $(p, i) \mapsto p'$ should
be done more delicately

Maybe instead of

$$\rho \sigma i,$$

$$\theta_{p_1} \theta_{p_2} \dots \theta_{p_n} \mid i_1 i_2 \dots i_m$$

$$\rho = (p_1, \dots, p_n)$$

$$i = (i_1, \dots, i_m)$$