

CPSC 421/501

Dec 1, 2023

- If 3COLOR (or any NP-complete problem) over $\Sigma = \{T, F\}$ has

Function $3\text{COLOR}_n(x_1, \dots, x_n)$

that require circuits of size $\geq Cn^k$

for all C, k , then $P \neq NP$.

- If so, then Function 3COLOR_n requires formulas of size $\geq Cn^k$.

- Let's start there. (1950's ... 1960's \rightarrow 2023)

- Let's start with any function $\{T, F\}^n \rightarrow \{T, F\}$

Today, Friday Dec 1:

Last day of new material!

Mon, Wed, Dec 4, 6

- 501 Presentations

- Exam practice

Dec 11 → Final Exam

Parity_n(x₁, ..., x_n) is defined as

$$(x_1 + \dots + x_n) \bmod 2 \quad \text{with } x_i \in \{0, 1\}$$

equivalently,

$$x_1 \text{ xor } x_2 \text{ xor } \dots \text{ xor } x_n \quad \text{with } x_i \in \{T, F\}.$$

Thm (Easy): Parity_n has formulas of

$$\text{size} \leq C n^2.$$

Thm (Subbotovskaya, 1961): Parity_n

requires formula size $\geq C n^{3/2}$

(where $C > 0$).

Fact! A function

$$f: \{T, F\}^n \rightarrow \{T, F\}$$

OR

$$\{0, 1\}^n \rightarrow \{0, 1\}$$

$$f = f(x_1, \dots, x_n)$$

(1) If f depends on all x_1, \dots, x_n

$$\text{Min Formula Size}(f) \geq n$$

(2) To prove

$$\text{Min Formula Size}(f) \geq n \quad | \text{BOCGO} |$$

is not easy ... began in

1961 ...

Any $f = f(x_1, \dots, x_n)$ can

be written in formula size

$$n \cdot 2^n \quad (\text{or } n2^n/2 \text{ or } \dots)$$

Most $2^{(2^n)}$ functions $\{T, F\}^n \rightarrow \{T, F\}$

require formulas size

$$2^n / (4 + \log_2 2^n)$$

$$\text{Parity}_n = \text{Parity}_n(x_1, x_2, \dots, x_n)$$

$$= (x_1 + \dots + x_n) \bmod 2 \quad \{0, 1\}$$

$$= x_1 \text{ XOR } x_2 \text{ XOR } \dots \text{ XOR } x_n \quad \{T, F\}$$

$$X_1 \oplus X_2 = \begin{cases} \top & \text{if } (X_1, X_2) \\ & = (T, F) \vee \\ & (F, T) \end{cases}$$

$$= (X_1 \wedge \neg X_2) \vee (\neg X_1 \wedge X_2)$$

So

$$f \oplus g =$$

$$(f \wedge \neg g) \vee (\neg f \wedge g)$$

$$X_1 \oplus \dots \oplus X_n = (X_1 \oplus X_2) \oplus (X_3 \oplus X_4)$$

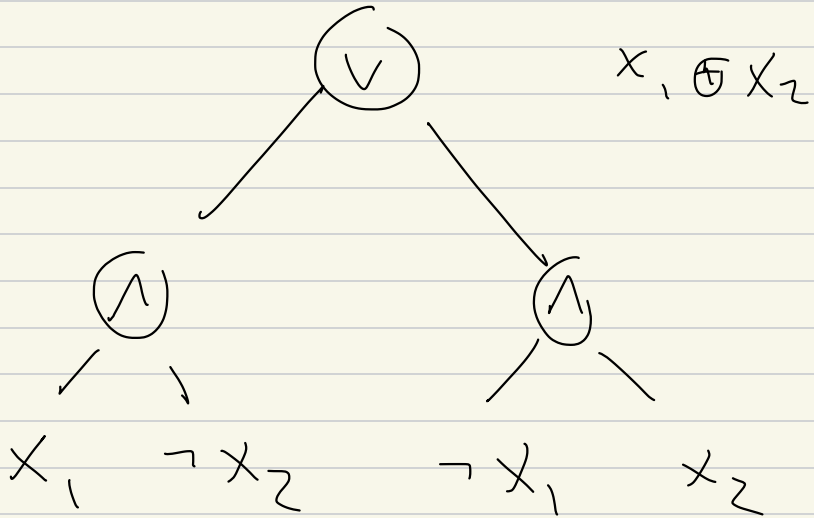
$$(f \wedge \neg g) \vee (\neg f \wedge g)$$

$$f = x_1 \oplus x_2$$

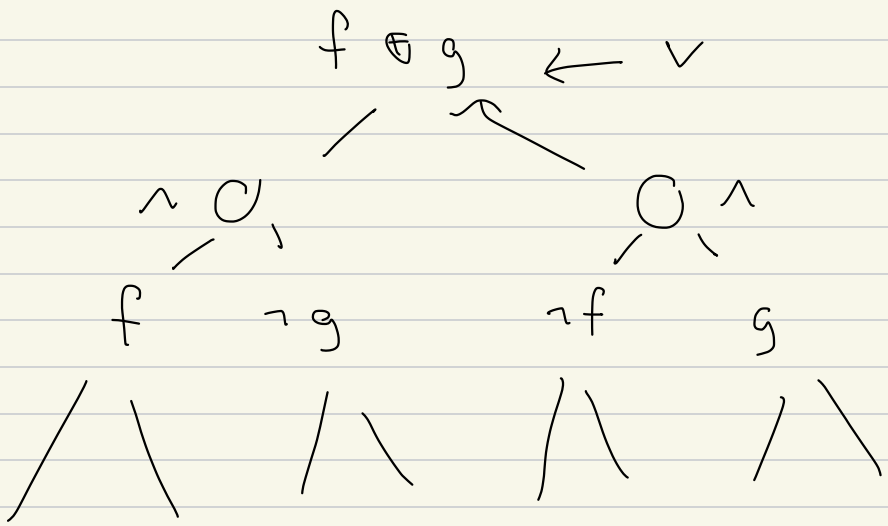
$$g = x_3 \oplus x_4$$

$$x_1 \oplus x_2 \oplus x_3 \oplus x_4$$

$$(x_1 \oplus x_2) \wedge \neg (x_3 \oplus x_4) \vee (\neg x_1 \oplus x_2 \wedge x_3 \oplus x_4)$$



Some x_3, x_4



Subbotovskaya (1961):

"restriction"

"random restrictions"

Let $L_n = \text{min formula size}$
to express $\text{Parity}_n(x_1, \dots, x_n)$.

Then:

$$(1) L_{n-1} \leq L_n \left(1 - \frac{3/2}{n}\right)$$

$$(2) L_n \geq \left(1 - \frac{3/2}{n}\right)^{-1} L_{n-1}$$
$$\geq \left(1 - \frac{3/2}{n}\right)^{-1} \left(1 - \frac{3/2}{n-1}\right)^{-1} L_{n-2}$$

$$\geq \underbrace{\left(1 - \frac{3/2}{n}\right)^{-1} \left(1 - \frac{3/2}{n-1}\right)^{-1} \dots \left(1 - \frac{3/2}{2}\right)^{-1}}_{\geq n^{3/2}} \cdot 1$$

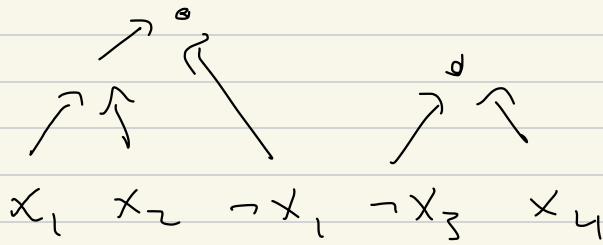
Rem: $\left[\left(1 - \frac{1}{n}\right) \left(1 - \frac{1}{n-1}\right) \dots \left(1 - \frac{1}{2}\right) \right]^{-1}$

$$\cancel{\left(\frac{n-1}{n}\right) \left(\frac{n-2}{n-1}\right) \dots \frac{1}{2}}$$

$$= \left[\frac{1}{n}\right]^{-1} = n$$

Pf: If Parity_n has formula size

L_n ,



L_n leaves,

then some x_i occurs $\geq \frac{L_n}{n}$ times

Say x_1 :

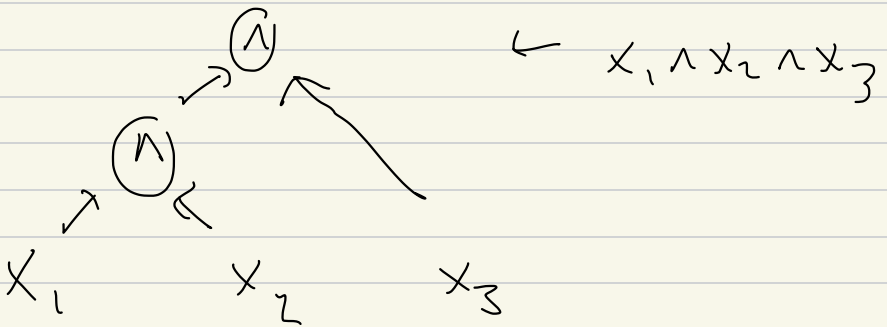
Set $x_1 = T$, see what happens

" $x_1 = F$, " " "

=

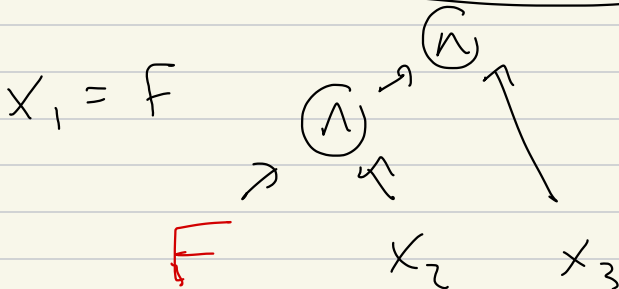
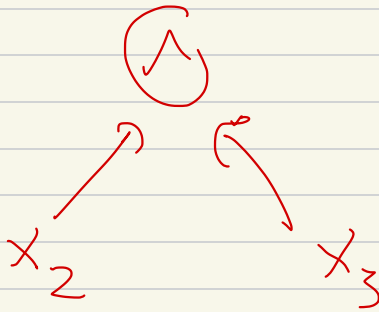
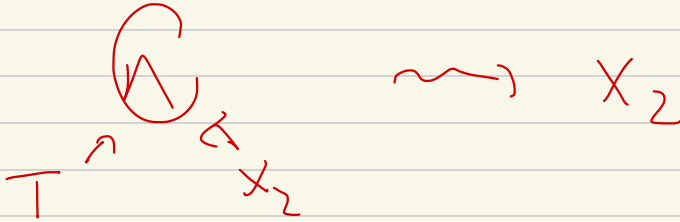
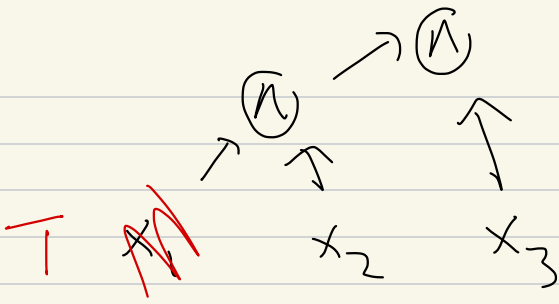
e.g.

$$x_1 \wedge x_2 \wedge x_3 :$$

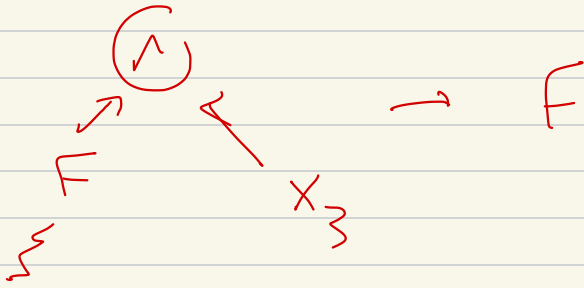
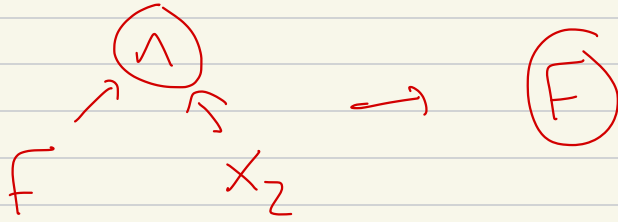


Restrict $x_1 = T$;

we get



$x_1 = F!$



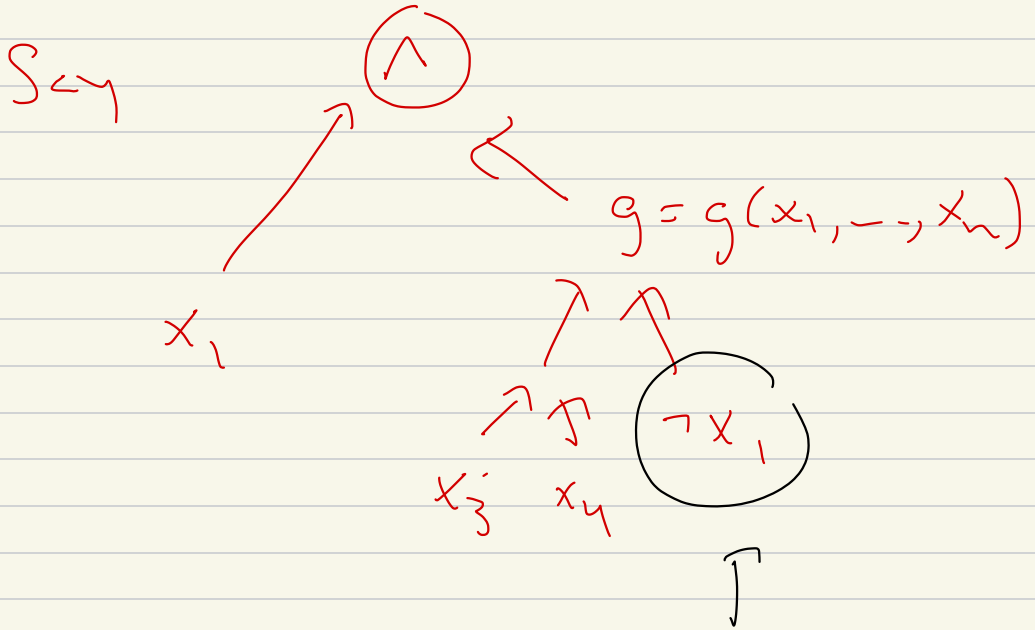
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Claim: Subbotovskaya:

by setting $x_1 = T, x_1 = f$

on average you eliminate $\geq 3/2$

(# of the x_i leaves)



Lemma:

if x_1 is on
other side, replace T

$$x_1 \wedge g(x_1, \dots, x_n) = T$$

\Leftrightarrow

$$x_1 \wedge g(T, x_2, \dots, x_n) = T$$

$$x_1 \wedge g = T$$

\Leftrightarrow

$$x_1 = T \text{ and } g = T$$

$$g = g(x_1, x_2, \dots)$$

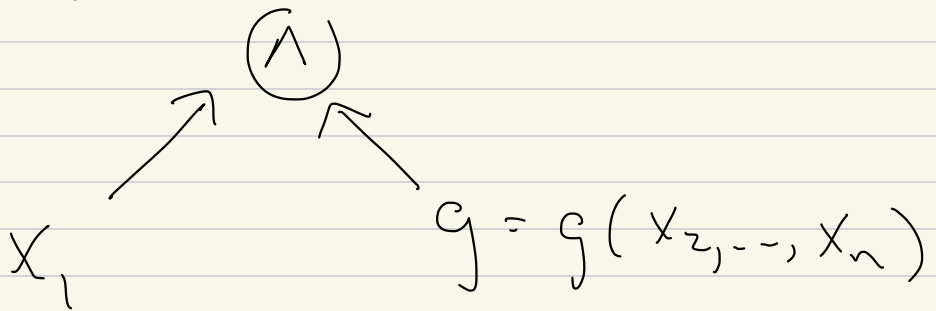
\Leftrightarrow

$$x_1 = T \text{ and } g(T, x_2, \dots, x_n) = T$$

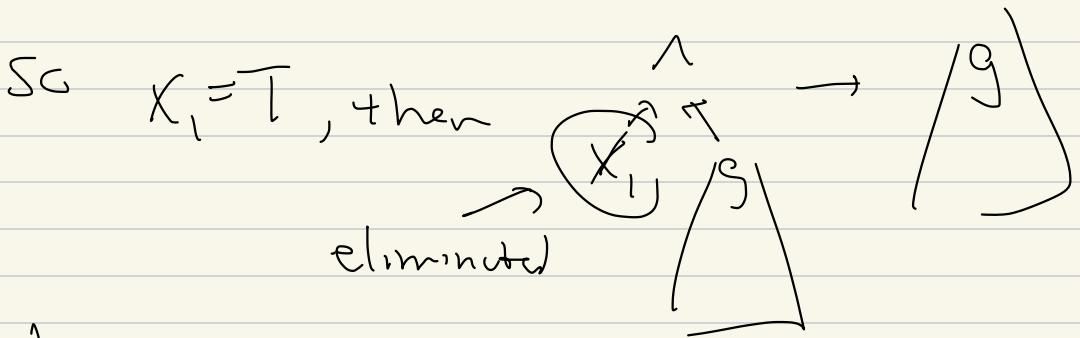
$$x_1 \vee g = T \Leftrightarrow$$

$$x_1 = T \text{ or } g = T \Leftrightarrow x_1 = T \text{ or } g = g(F, x_2, \dots, x_n)$$

So if



if min size formula



but

$x_1 = F$, then

$$x_1 \wedge g \rightarrow f \wedge g \rightarrow F$$

Hence

$X_1 \rightarrow T$ eliminates X_1

$X_1 \rightarrow F$ " X_1 and least
one other
leaf

$X_1 \rightarrow T$ eliminates 1

$X_1 \rightarrow F$ " 2

$\rightarrow \geq \frac{3}{2}$ on average!!