

GROUP HOMEWORK 10, CPSC 421/501, FALL 2021

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Please note:

- (1) You must justify all answers; no credit is given for a correct answer without justification.
- (2) Proofs should be written out formally.
- (3) Homework that is difficult to read may not be graded.
- (4) You may work together on homework in groups of up to four, **but you must submit a single homework as a group submission under Gradescope.**

THIS HOMEWORK IS NOT TO BE HANDED IN. IT IS QUITE SHORT.

Let $M = (Q, \Sigma, \Gamma, \delta, q_0, q_{\text{acc}}, q_{\text{rej}})$ be a one-tape Turing machine. Consider the configuration notation in class on November 30 and in [Sip], Chapter 3. For example: say that a one-tape TM is in state q_2 , and contents of the tape is

$$abba \sqcup xabba \sqcup \sqcup \sqcup \dots$$

where \sqcup is the blank symbol, and $x, a, b \in \Gamma$; for simplicity we assume that the sets Q and Γ have no intersection (otherwise we have to work with the *disjoint union*¹). Say that the tape head is on cell number two. Then we denote this configuration by

$$aq_2bba \sqcup xabba$$

having dropped the contiguous infinite string of \sqcup 's that extends infinitely to the right at some point, and having inserted q_2 in the second position to indicate the tape head position. Similarly the strings

$$abbq_2a \sqcup xabba, \quad abbaq_2 \sqcup xabba$$

denote the same configuration, except that the tape head location is, respectively, on cells numbers four and five.

Research supported in part by an NSERC grant.

¹The disjoint union is a limit, and defined only uniquely up to unique isomorphism, and is not defined in any “canonical” way. (The term “canonical” is not precise, but most humans get a feel for what this means.)

- (1) Say, with the above understanding, a one-tape TM is in state q_0 , and the contents of the tape is

$abba \square \square \square \dots$

What is the formal representation configuration of the Turing machine when:

- (a) the tape head is on cell one?
 - (b) the tape head is on cell two?
 - (c) the tape head is on cell four?
 - (d) the tape head is on cell eight?
- (2) Given a one-tape Turing machine, is the set of all possible configurations of the Turing machine countably infinite? Explain.
- (3) In what sense can one view a Turing machine transition from one step to another as a “local computation”? [Hint: consider the above two problems, Section 2 of the recently posted handout on the Cook-Levin theorem, class notes, and/or [Sip], Chapter 3.]

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