

- Set Theory Subtleties

Not subtle: If T is uncountable, and there is a surjection $S \rightarrow T$, then S is uncountable.

Subtle: If there is an injection $S \rightarrow T$, then there is a surjection $T \rightarrow S$.

- Russell's Paradox

- Related Paradoxes (Section 6 of handout)

- Start Finite Automata (§ 1.1 of Textbook)??

BREAKOUT ROOM PROBLEMS

① If S is countable, and there is a surjection $S \rightarrow T$, then T is countable. (Prove this)

related

② Is $\text{POWER}(\{a,b\}^*) \cup \mathbb{N}$ countable?

③ Is there a bijection

$$[2]^{\mathbb{N}} \rightarrow [3]^{\mathbb{N}}, \quad \leftarrow \text{new notation}$$

and can you describe one?

④ Let F be the set of functions $\mathbb{N} \rightarrow \mathbb{N}$ that can be described in English

(assuming a fixed, precise interpretation of English). Is F countable or uncountable?

⑤ Let $\mathbb{N}^{\mathbb{N}}$ be the set of functions $\mathbb{N} \rightarrow \mathbb{N}$. Is $\mathbb{N}^{\mathbb{N}}$ countable or uncountable?

related

Today: - Finish the handout

- Breakout rooms (students may be able to form their own breakout rooms, as of yesterday)

§5 of handout: Russell's Paradox & Set Theory Subtleties

OR (1) Why we consult with set theorists & logicians

(2) How to find more countable and uncountable sets

Last time:

Cantor's Thm: Any map $S \xrightarrow{f} \text{POWER}(S)$ is not

surjection. Idea $T = \{s \in S \mid s \notin f(s)\}$ ← kind of a self-ref. negation

is not in the image of f . (Main tool for uncountable sets.)

Implies: $\text{POWER}(\mathbb{N})$ is uncountable

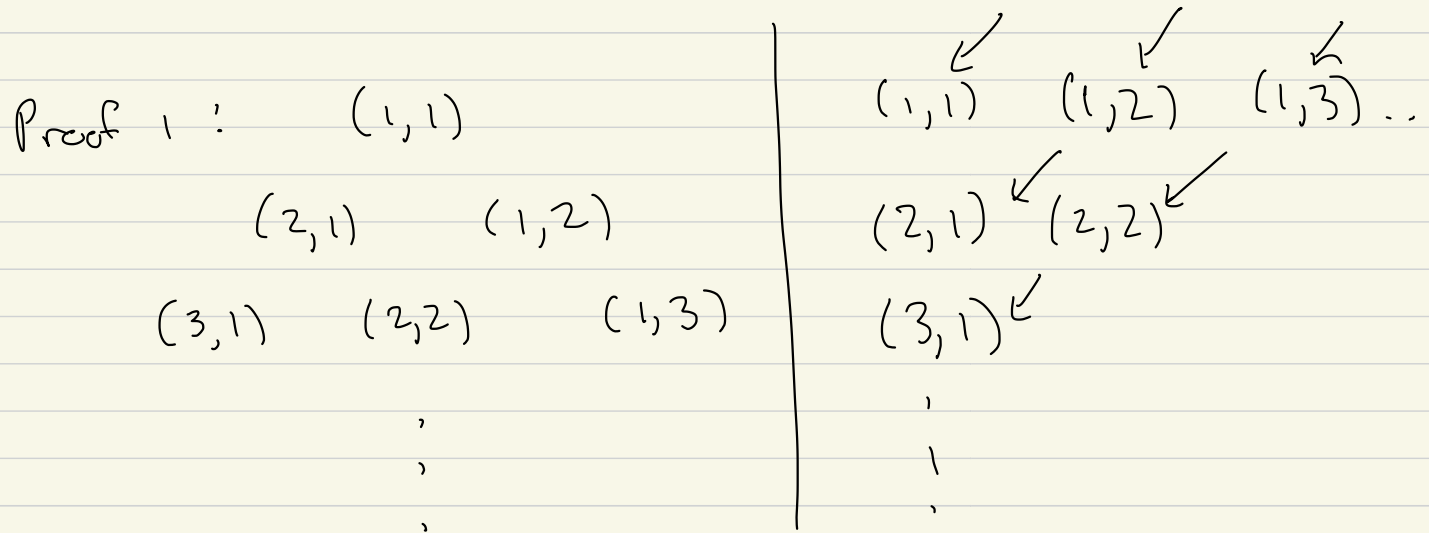
→ application $\text{POWER}(\Sigma^*)$ " " $\Sigma = \text{alphabet}$

Textbook, §4.2 proves that \mathbb{R} is uncountable, whereas

\mathbb{N} , \mathbb{Z} (integers), \mathbb{Q} (rational numbers) are countable

Last time $\mathbb{N} \times \mathbb{N} = \mathbb{N}^2 = \{(a,b) \mid a,b \in \mathbb{N}\}$

is countable



Idea

(3,5)	\rightsquigarrow	35
(17,1)	\rightsquigarrow	171
\mathbb{N}^2	\rightarrow	\mathbb{N}

Fix: (Gödel numbers)

(3,5)	\rightarrow	$2^3 3^5$
(17,1)	\rightarrow	$2^{17} 3^1$
:		:
(a,b)	\rightarrow	$2^a 3^b$

Given map $\mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ injection

Since $2^a 3^b = 2^{a'} 3^{b'} \Rightarrow \begin{matrix} a=a' \\ b=b' \end{matrix}$

Our definition! A is uncountable, if there exist no map $f: \mathbb{N} \rightarrow A$ that is surjective.

Here! $\mathbb{N}^2 \rightarrow \mathbb{N}$ is injective

Can you prove: if there is an injection $S \rightarrow \mathbb{N}$, is there a surjection $\mathbb{N} \rightarrow S$?????

This is true for finite sets

$$[3] = \{1, 2, 3\}, \quad [4] = \{1, 2, 3, 4\}$$

injection

$$[3] \rightarrow [4]$$

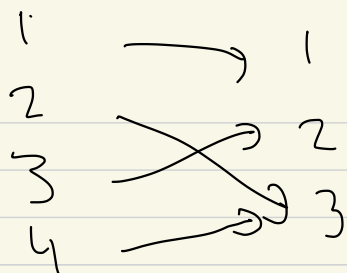
OR

$$[3] \rightarrow [4]$$

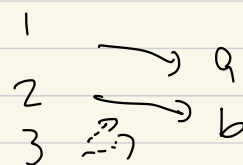
$$\begin{array}{ccc} 1 & \longrightarrow & 1 \\ 2 & \longrightarrow & 2 \\ 3 & \longrightarrow & 3 \\ & & 4 \end{array}$$

$$\begin{array}{ccc} 1 & & 1 \\ & \searrow & \\ 2 & & 2 \\ & \searrow & \\ 3 & & 3 \\ & \searrow & \\ & & 4 \end{array}$$

and surjection $[4] \rightarrow [3]$



Intuition! If there is an injection $S \rightarrow T$
 "size(S) \leq size(T)"



Intuition! If there is a surjection $T \rightarrow S$
 then "size(T) \geq size(S)"

Is it true that (1) \exists surjection $T \rightarrow S$ \Rightarrow
 \Leftrightarrow (2) \exists injection $S \rightarrow T$

For S, T finite, or $= \mathbb{N}$ yes

" $S, T \rightsquigarrow$ subtle (assumptions on "size" "cardinality")"

If we assume this, then

\exists surjection $\mathbb{N} \rightarrow S$
 "countable"

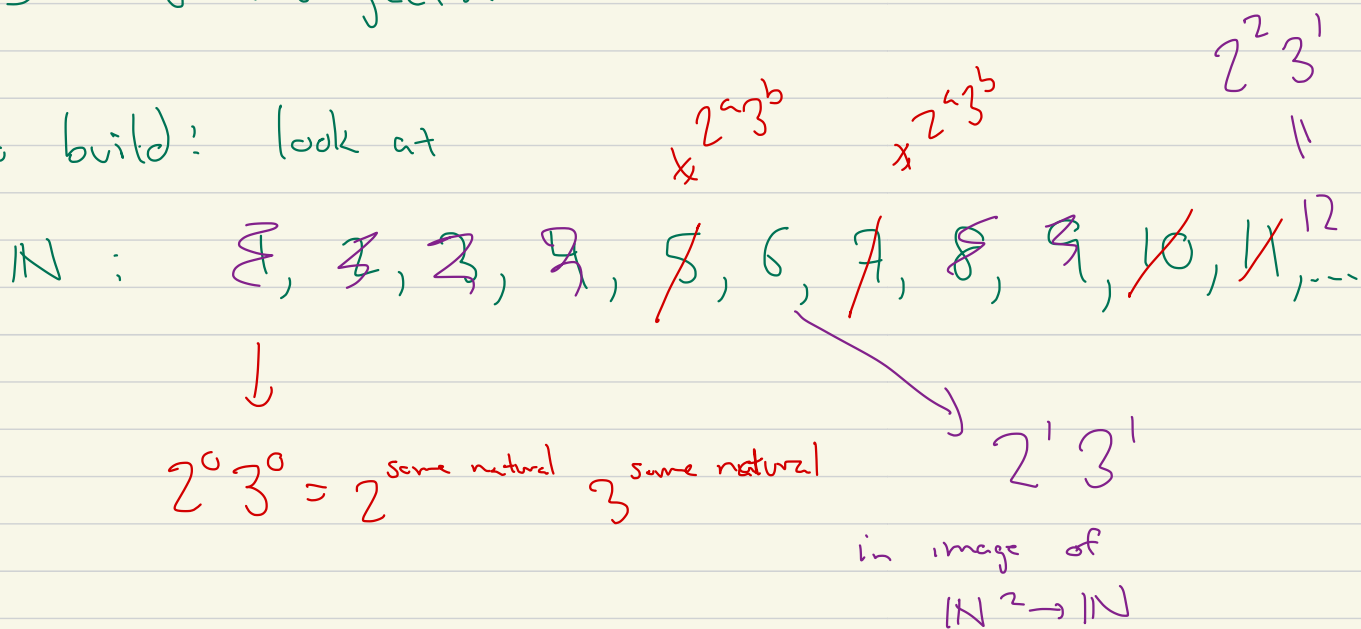
$\Leftrightarrow \exists$ injection $S \rightarrow \mathbb{N}$
 an equivalent form

Example: $\mathbb{N}^2 \rightarrow \mathbb{N}$

$(a, b) \mapsto 2^a 3^b$ is an injection

$\Rightarrow \exists$ a surjection $\mathbb{N} \rightarrow \mathbb{N}^2$

To build: look at



not easy to describe map $\mathbb{N} \rightarrow \mathbb{N}^2$ based on the map $\mathbb{N}^2 \rightarrow \mathbb{N}$

If S is countable, and there is a surjection $S \rightarrow T$, then T is countable

If S is countable, and there is a ~~surjection~~ ^{an injection} $S \rightarrow T$, then T is countable

$T \rightarrow S$

Claim: If T is uncountable and

are equivalent $\left\{ \begin{array}{l} \text{there is a surjection } S \rightarrow T \Rightarrow S \text{ is uncountable} \\ \text{" " an injection } T \rightarrow S \Rightarrow \text{" " "} \end{array} \right.$

Russell's Paradox: Let R be the set of sets that do not contain themselves.

Does $R \in R$?

Idea of Paradox: You have
① self-reference
② negation

$R = \{ S \text{ is a set s.t. } \underline{S \notin S} \}?$

Either $R \in R$ or $R \notin R$.

① $R \in R$ then R is a set s.t. $R \notin R$. \leftarrow contradiction

② $R \notin R$ then R is not a set s.t. $R \notin R$, i.e. $R \in R$
 \rightarrow also a contradiction \leftarrow

Related paradox: Does the set of all sets contain itself.

Next: we describe a few more $\left\{ \begin{array}{l} \text{"paradoxes"} \leftarrow \text{Russell's} \\ \text{"tlms"} \leftarrow \text{Cantor's thm.} \end{array} \right.$

$$[n] = \{1, 2, \dots, n\}$$

Notation: $\left\{ \begin{array}{l} \text{functions from } S \rightarrow T \end{array} \right\}$
 \uparrow set of function \uparrow

Notation T^S

E.g. $\text{PRIMES} = \{2, 3, 5, 7, 11, \dots\} \subseteq \{0, \dots, 9\}^*$
 \uparrow left out 02, 03

i.e. PRIMES is a language over $\{0, \dots, 9\} = \Sigma$
 alphabet

PRIMES

0	→	no
1	→	no
2	→	yes
3	→	yes
⋮		
9	→	no
00		
01		

02 → $\begin{cases} \text{no} \\ \text{yes} \end{cases}$

depends on if you allow leading zeros

More generally

equiv
① language over Σ
② elt of $\text{Power}(\Sigma^*)$

③ a decision $\Sigma^* \rightarrow \{\text{yes, no}\}$

④ an element of

$$\{\text{yes, no}\}^{\Sigma^*} = \{\text{functions } \Sigma^* \rightarrow \{\text{yes, no}\}\}$$

~~10:26~~ \leftrightarrow ~~10:43~~

My watch 10:28 am \rightarrow 10:43

15 minute breakout, or so

Breakout: ④ & ⑤

function $\mathbb{N} \rightarrow \mathbb{N}$ describe in English.

"The function taking 1 to 17, 2 to 23, ..."

allow infinite

But finite strings

"The function taking x to $\text{sqrt}(x)$."

↳ interpret

Strings in ASCII = ASCII* countable

⑤ { functions $\mathbb{N} \rightarrow \mathbb{N}$ } uncountable

functions $\mathbb{N} \rightarrow \{1, 2\}$ ↗ bijection
↔ $\mathbb{N} \rightarrow \{\text{yes, no}\}$

↔ POWER(\mathbb{N}) uncountable

① S countable, surj $S \rightarrow T$ then

② T is countable

⇒ surj ~~S~~ $\rightarrow T$ and ~~T~~ uncountable

then \mathcal{S} uncountable

Next time: finish this discussion

- Resolve Russell's paradox
- 1-2 paradox
- Textbook Ch.1, Regular Lang.

$$\Sigma = \{a, b, c\}$$

Σ

Σ^0

a b c

Σ^1

aa ab ac ... cc Σ^2

$$\{0, 1\}^* = \{ \mathbb{N} = \{ \textcircled{1}, \textcircled{2}, \textcircled{3}, \textcircled{4}, \textcircled{5}, \dots \} \}$$

$\Sigma, c, l, cc, cl, \dots$

Some
function
or meaningless

Some
function
or
meaningless