## INDIVIDUAL HOMEWORK 7, CPSC 421/501, FALL 2020

## JOEL FRIEDMAN

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Please note:

- (1) You must justify all answers; no credit is given for a correct answer without justification.
- (2) Proofs should be written out formally.
- (3) Homework that is difficult to read may not be graded.
- (4) You may work together on homework in groups of up to four, **but you must write up your own solutions individually and must acknowledge with whom you worked.** You must also acknowledge any sources you have used beyond the textbook and two articles on the class website.
- (1) Recall that

 $\text{HALT}_{\text{TM}} = \{ \langle M, w \rangle \mid M \text{ is a Turning machine that halts on input } w \}.$ 

Give a direct proof that HALT<sub>TM</sub> is undecidable (i.e., without using the fact that  $A_{TM}$  is undecidable); do so using the same strategy used to prove that  $A_{TM}$  is undecidable (see the proof Theorem 4.11 in [Sip], page 207, or class notes on 10\_22): assume that some Turing machine H decides HALT<sub>TM</sub>, and describe how to build a Turing machine D such that when you consider  $D(\langle D \rangle)$  (i.e., the outcome of D on input  $\langle D \rangle$ ), you get a contradiction.

(2) Is HALT<sub>TM</sub> (defined above) recognizable? Is the complement of HALT<sub>TM</sub> recognizable? Justify your answers.

DEPARTMENT OF COMPUTER SCIENCE, UNIVERSITY OF BRITISH COLUMBIA, VANCOUVER, BC V6T 1Z4, CANADA.

*E-mail address*: jf@cs.ubc.ca *URL*: http://www.cs.ubc.ca/~jf

Research supported in part by an NSERC grant.