

## GROUP HOMEWORK 4, CPSC 421/501, FALL 2020

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Please note:

- (1) You must justify all answers; no credit is given for a correct answer without justification.
- (2) Proofs should be written out formally.
- (3) Homework that is difficult to read may not be graded.
- (4) You may work together on homework in groups of up to four, **but you must submit a single homework as a group submission under Gradescope.**

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- (1) Let  $L$  be the language of strings that contain  $ab$  as a substring. Use the Myhill-Nerode theorem to show that any DFA accepting  $L$  has at least 3 states.
  - (2) (a) Let  $L$  be the language described by the regular expression  $(0 \cup 1)^* 1 (0 \cup 1)^3$ , i.e., the words of length at least 4 whose 4th last symbol is 1. Write an NFA with 5 states that recognizes  $L$ .  
(b) Generalize part (a) to the language  $(0 \cup 1)^* 1 (0 \cup 1)^k$  for any  $k \in \mathbb{N}$ .
  - (3) (a) Let  $L$  be the language described by the regular expression  $(0 \cup 1)^* 1 (0 \cup 1)^1$ , i.e., the words of length at least 2 whose 2nd last symbol is 1. Use the Myhill-Nerode to find minimum number of states needed in a DFA that recognizes  $L$ . [Hint: Show that  $\text{AccFut}_L(u)$  for  $u = 00, 01, 10, 11$  are all different. Are there any other possible values for  $\text{AccFut}_L(u)$  with  $u$  of length 3 or greater? What about for  $u$  of length 0 or 1?]  
(b) Give a DFA that recognizes  $L$  with fewest possible states, and explain how your DFA works.  
(c) Generalize parts (a,b) to the language  $(0 \cup 1)^* 1 (0 \cup 1)^k$  for any  $k \in \mathbb{N}$ .

**The following is a bonus question, worth an additional 10 points out of 100.**

- (4) Let  $L$  be a regular language over an alphabet  $\Sigma$ , and let  $L'$  be any language (not necessarily regular) over  $\Sigma$ . Is the language

$$L'' = \{w \in \Sigma^* \mid \text{for some } w' \in L' \text{ we have } ww' \in L\}$$

necessarily regular? [Hint: One way to do this is to write  $\text{AccFut}_{L''}(u)$  in terms of certain values of  $\text{AccFut}_L$ .]

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