

Make sure!

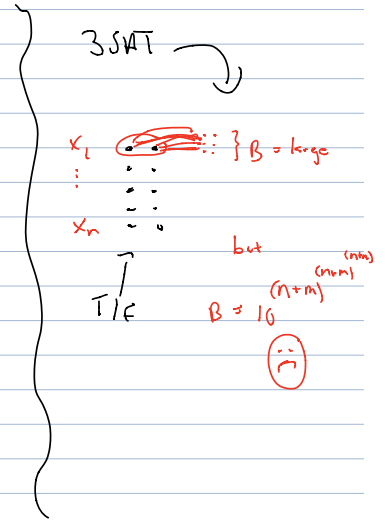
$$L_1 \leq p L_2$$

$$f: \Sigma_1^* \rightarrow \Sigma_2^* \text{ in } \text{poly time}$$

Last time

$$f: \left( \begin{array}{l} \text{string} \\ \text{rep} \\ \text{3CNF} \end{array} \right) \rightarrow \langle G, a, b \rangle$$

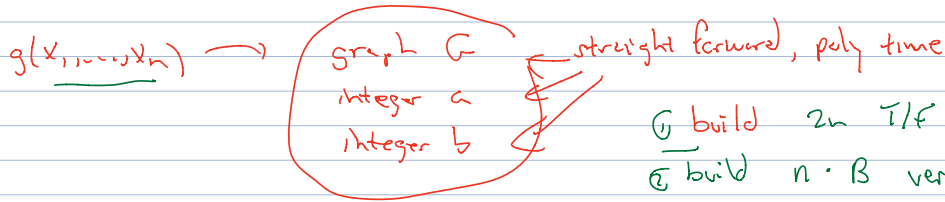
not :(



$$(1) S = 3SAT \Rightarrow f(S) \in \text{VERTEX-EXPANSION}$$

and (2)  $\Leftarrow$

or (2')  $S \notin 3SAT \Rightarrow f(S) \notin \text{VERTEX-EXPANSION}$



- ① build 2n T/F vertice
- ② build  $n \cdot B$  vertices

"T/F enforcement" vertices  
"dummy vertices"  
+ edges

$$3SAT = \left\{ \langle g \rangle \mid \begin{array}{l} \text{Boolean formula} \\ \text{in 3CNF form} \\ \text{that is satisfiable} \end{array} \right\}$$

- ③ build each of  $m$  clauses:

1 vertex  
3 edges

$$S \notin 3SAT \Leftrightarrow$$

$$\left\{ s \mid \begin{array}{l} \text{either } s \neq \langle g \rangle \text{ for } g \text{ in 3CNF} \\ \text{or } s = \langle g \rangle \text{ but } g \text{ is not} \\ \text{satisfiable} \end{array} \right\}$$

- ④ write  $a = n = \# \text{ Bool vars}$

$b =$  see homework

$$\left( \begin{array}{l} \text{maybe } nB + c, \\ B = n + m + 1 \end{array} \right)$$

Today: Start Ch 8: Space

=

Given  $\epsilon$  Tim.,  $M$ ,  $M$  runs in space  $f(n)$ ,

$f: \mathbb{N} \rightarrow \mathbb{N}$  if on any input,  $w$ ,  $M$  uses  
or  $\mathbb{Z}_2 \rightarrow \mathbb{D}_{2,0}$

at most  $O(f(|w|))$  space



SPACE( $f(n)$ )

space = maximum # of cells used in the entire algorithm (whether or not we actually change their content).

=

$$\text{PSPACE} = \text{Polynomial Space} = \bigcup_{k \in \mathbb{N}} \text{SPACE}(n^k)$$

Given  $\epsilon$  <sup>non-deterministic</sup> Tim.,  $M$ ,  $M$  runs in space  $f(n)$ ,

$f: \mathbb{N} \rightarrow \mathbb{N}$  if on any input,  $w$ ,  $M$  uses  
or  $\mathbb{Z}_2 \rightarrow \mathbb{D}_{2,0}$

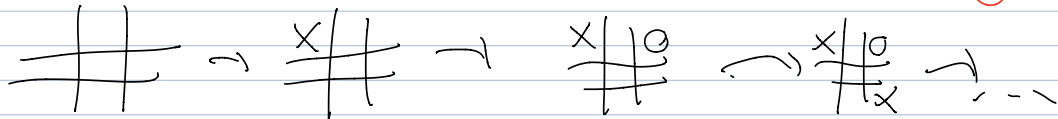
NSPACE( $f(n)$ ) at most  $O(f(|w|))$  space in any computation path.


$$\text{NPSPACE} = \text{Non-deterministic Polynomial Space} = \bigcup_{k \in \mathbb{N}} \text{NSPACE}(n^k)$$

Fact:  $NPSPACE = PSPACE$

Savitch's Thm:  $NSPACE(n^k) \subseteq SPACE(n^{2k})$

Break: How complicated is  $3 \times 3$  tic-tac-toe? 😊



Is it more complicated than  
- chess  $\lfloor$  😞  
- go  $19 \times 19$   😞  
 $3^{(19^2)}$

How complicated is  $4 \times 4 \times 4$  tic-tac-toe? 😞

$7 \times 7 \times 7 \times 7$  " ? 😞

How many configurations are there in  $3 \times 3$  tic-tac-toe? 😞

Upper bounds:  $\leq 3^9$  9 squares, X, O,  $\perp$ , alternate

How many configurations of a  $3 \times 3$  board are there  
with each square either  $\{X, O, \perp\}$ ?  $3^9$  😊

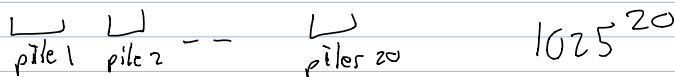
How many config ---  $7 \times 7 \times 7 \times 7$  board ---  
?  $3^{(7^4)}$

---  $7 \times 7 \times 7 \times 7 \times 7 \times 7 \times 7$  ---  
?

JF-game: 20 piles of chips, 1024 chips in each pile

- 2 players, alternate n moves,  
each move as many chips in any number of piles
- player to first clear out all piles wins...

# configs:

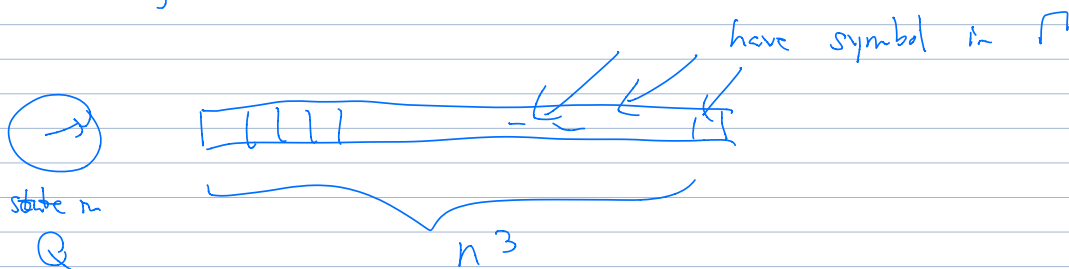



Game is not complex...

Back to TMs:

Say T.M., M, runs in SPACE  $\leq n^3$

How long could M take?



# poss configs = 

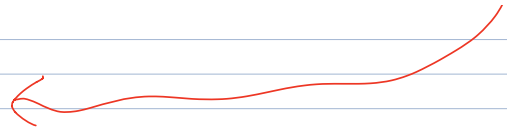
$$\leq |Q| \times |\Gamma|^{n^3} \times n^3$$

# states      tape      tape head

time 1  
time 2

← config      ← you loop!  
← same config as

time  $M$  takes  $\leq$



What if  $M$  is non-det??