

Section 1.3 Regular Expressions

- Regular Expressions (Limited Set)

- Equivalent to Regular Languages

=

Regular Expression: over alphabet Σ :

(1) $\{\sigma\}$ $\sigma \in \Sigma$

(2) ϵ $\epsilon =$ empty string

(3) \emptyset empty set

—

If R_1, R_2 regular expressions, then

(4) $(R_1 \cup R_2)$ is a regular expression

(5) $(R_1 \circ R_2)$, or simply $(R_1 R_2)$ is a regular expression

(6) (R_1^*) is a regular expression

=

e.g. $\Sigma = \{a, b\}$

- Then $(a \cup b)^* b$ is a regular expression

- Express two-letter words: $(a \cup b) \circ (a \cup b)$

$\Sigma = \{a, b\}$

$\underbrace{\hspace{2cm}}_{\Sigma} \quad \underbrace{\hspace{2cm}}_{\Sigma}$

← shorthand

$\Sigma \circ \Sigma$


$\Sigma^2 = \Sigma \circ \Sigma$ ← shorthand

- Express single word a^5

$((a \circ a) \circ (a)) \circ (a \circ a)$ write as a^5

$a \circ a \circ a \circ a \circ a$ write as a^5

$$\Sigma = \{a, 1, \dots, 9\}$$

Refer to $\{1, 3, 5, 7, 9\} \leftrightarrow (((1 \cup 3) \cup 5) \cup 7) \cup 9$ 

$1 \cup 3 \cup 5 \cup 7 \cup 9$
 $1, 3, 5, 7, 9 \quad \cup \leftrightarrow ,$

Given regular expression, R , there is an associated language $L(R)$.

$$(a^5 \cup a^7)^* , \quad L((a^5 \cup a^7)^*) = \{\epsilon, a^5, a^7, a^{10}, a^{12}, a^{14}, \dots\}$$

Sometimes

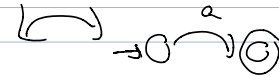
$$L^+ = \underline{L^1 \cup L^2 \cup L^3 \cup \dots} = \underline{L L^*} = \left\{ \begin{array}{l} \text{one or more elements of } L \\ \text{concatenated together} \end{array} \right\}$$

So we allow $L((abb)^+) = \{abb, abbb, \dots\}$

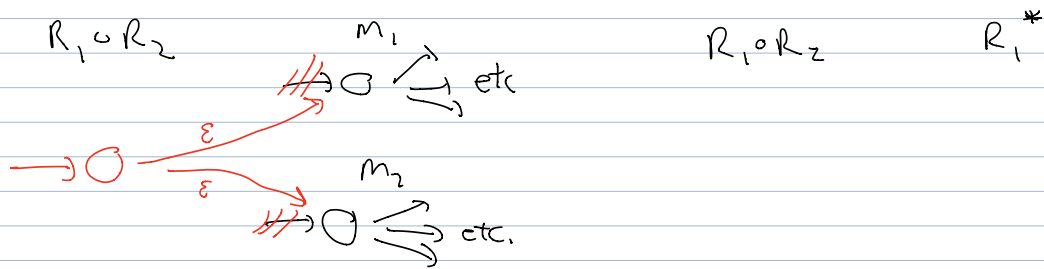
Thm: L is regular \iff there is a regular expression describing L

Textbook: $\emptyset^* \stackrel{?}{=} \{\epsilon\} = \emptyset^0$
 $= \emptyset^0 \cup \emptyset^1 \cup \dots$

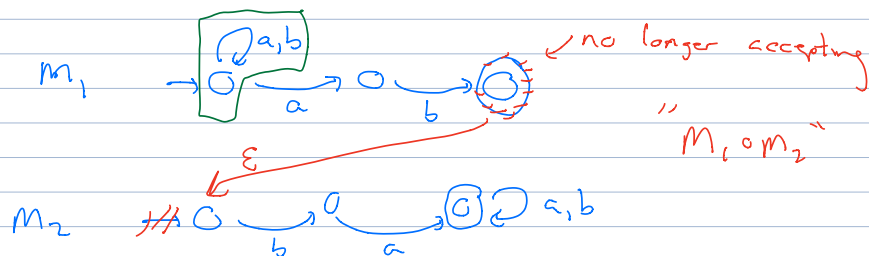
Easy! Given a regular expression, R , we can draw an NFA that recognizes the language R describes.

Why? $\Sigma = \{a, b\}$ regular expression a 

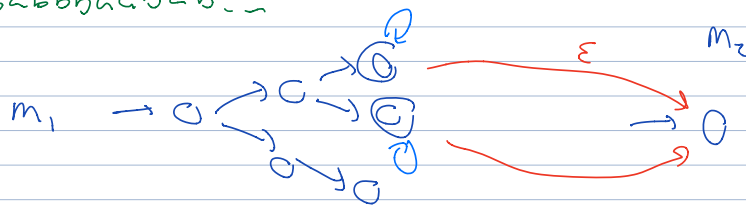
Given R_1, R_2 to each NFA, M_1, M_2
 ϵ
 \emptyset



$R_1 \circ R_2$: $R_1 = \Sigma^* ab$ $R_2 = ba \Sigma^*$, $R_1 \circ R_2 = \Sigma^* abba \Sigma^*$
 $\Sigma = \{a, b\}$



input $abbbbabbbbaabab...$



=

Part 2: Given a DFA/NFA \rightsquigarrow reg expression

e.g. $DIV_BY_3 = \{0, 3, 6, 9, 12, 15, \dots\}$

$DIV_BY_2 = \{0, 2, 4, 6, 8, 10, 12, \dots\}$

$((0 \cup 2 \cup 4 \cup 6 \cup 8) \cup ((1 \cup 2 \cup 3 \cup 4 \cup 5 \cup 6 \cup 7 \cup 8 \cup 9) \Sigma^* (0 \cup 2 \cup 4 \cup 6 \cup 8)))$

≤ 1 digit ≥ 2 digits could be ϵ \rightarrow 2 digit number

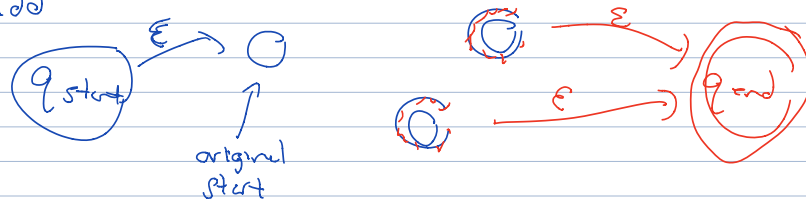
DIV-BY-2 - LEADING ZEROS - OK - EMPTY STRING - NOT OK

$$\Leftrightarrow \Sigma^*(0 \cup 2 \cup 4 \cup 6 \cup 8)$$

There is a procedure: take NFA: insert!



could add



intermediate state

