

Local Search

Computer Science cpssc322, Lecture 14
(Textbook Chpt 4.8)

Oct, 5, 2012

Announcements

- Assignment1 due now!
- Assignment2 out next week

Lecture Overview

- Recap solving CSP systematically
- Local search
- Constrained Optimization
- Greedy Descent / Hill Climbing:
Problems

Systematically solving CSPs: Summary

- Build Constraint Network

- Apply Arc Consistency

- One domain is empty → *no sol*
- Each domain has a single value → *unique sol*
- Some domains have more than one value → *?!*
may or maynot have a solution

- Apply Depth-First Search with Pruning

- Split the problem in a number of disjoint cases
 - Apply Arc Consistency to each case

Lecture Overview

- Recap
- **Local search**
- Constrained Optimization
- Greedy Descent / Hill Climbing:
Problems

Local Search motivation: Scale

- Many CSPs (scheduling, DNA computing, more later) are simply too big for systematic approaches
- If you have 10^5 vars with $\text{dom}(\text{var}_i) = 10^4$
- Systematic Search
- Constraint Network

$$10^5 * 10^4$$

$$10^{10} * 10^8$$

$$10^{10} * 10^{12}$$

- but if solutions are densely distributed.....

Local Search motivation: Scale

- Many CSPs (scheduling, DNA computing, more later) are simply too big for systematic approaches
- If you have 10^5 vars with $\text{dom}(\text{var}_i) = 10^4$

- Systematic Search

$$\left\{ \begin{array}{l} b = 10^4 \\ d = 10^5 \end{array} \right. \quad (10^4)^{10^5}$$

branching factor depth

- Constraint Network

$$\underbrace{10^5}_{\text{var nodes}} + \underbrace{10^5 \times 10^5}_{\text{constraint nodes}}$$

10^{10} max # of nodes

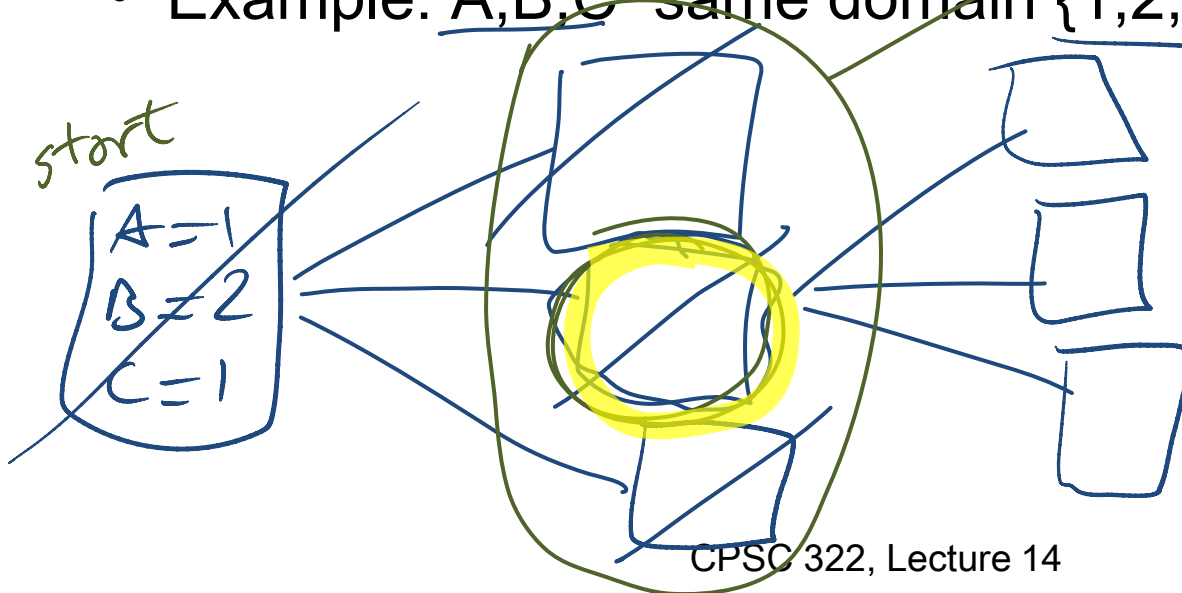
- but if solutions are densely distributed.....

Local Search: General Method

Remember , for CSP a solution is...?... possible world

(not a path)

- Start from a possible world
- Generate some neighbors (“similar” possible worlds)
- Move from the current node to a neighbor, selected according to a particular strategy
 - Example: A,B,C same domain {1,2,3}



Local Search: Selecting Neighbors

How do we determine the **neighbors**?

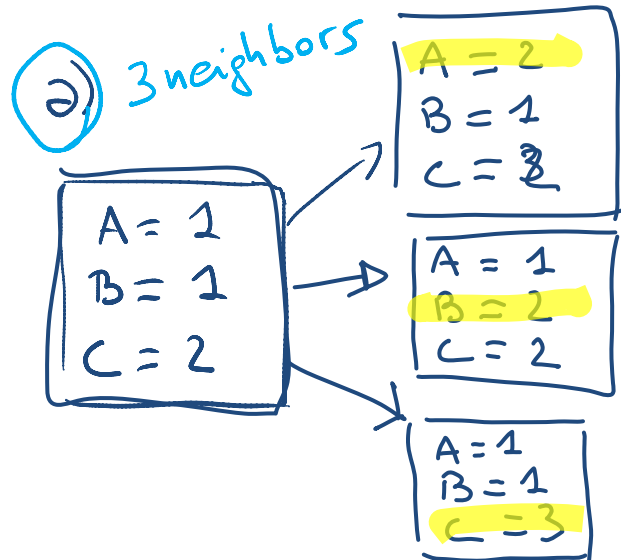
- Usually this is simple: some small incremental change to the variable assignment

a) assignments that differ in one variable's value, by (for instance) a value difference of +1

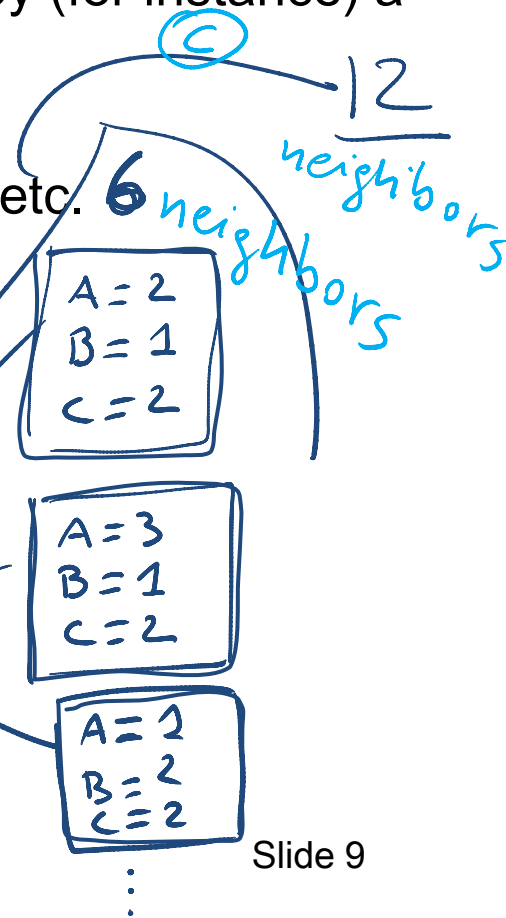
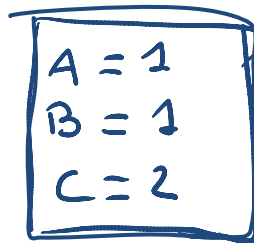
b) assignments that differ in one variable's value

c) assignments that differ in two variables' values, etc.

- Example: A, B, C same domain {1, 2, 3}



b)



Iterative Best Improvement

- How to determine the neighbor node to be selected?
- **Iterative Best Improvement:**
 - select the neighbor that optimizes some evaluation function
- Which strategy would make sense? Select neighbor with ...

Maximal number of constraint violations

Similar number of constraint violations as current state

No constraint violations

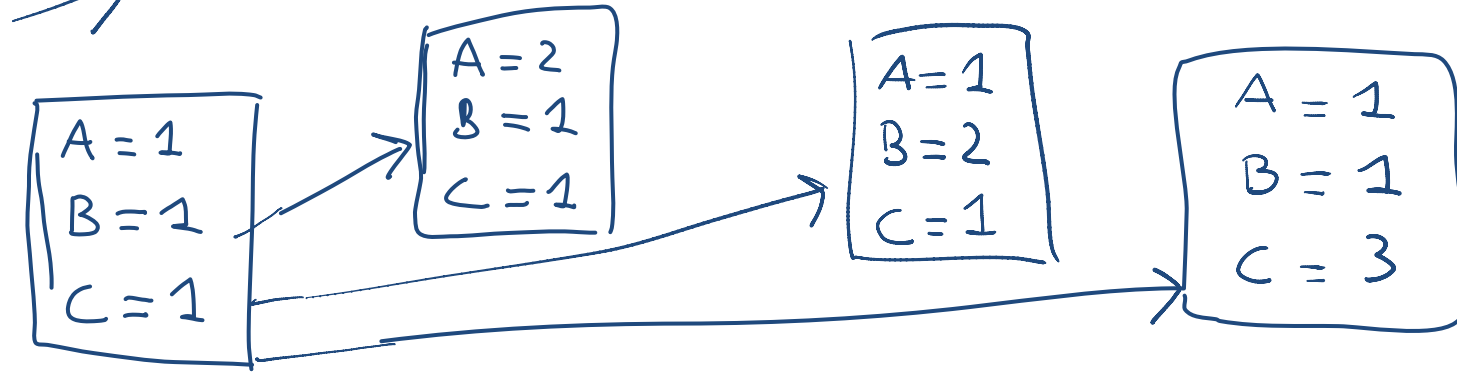
Minimal number of constraint violations

Iterative Best Improvement

- How to determine the neighbor node to be selected?
- **Iterative Best Improvement:**
 - select the neighbor that optimizes some evaluation function
- Which strategy would make sense? Select
Minimal number of constraint violations
- **Evaluation function:**
 $h(n)$: number of constraint violations in state n
- **Greedy descent:** evaluate $h(n)$ for each neighbour, pick the neighbour n with minimal $h(n)$
- **Hill climbing:** equivalent algorithm for maximization problems
 - Here: maximize the number of constraints satisfied

Selecting the best neighbor

- Example: A,B,C same domain {1,2,3} , (A=B, A>1, C≠3)



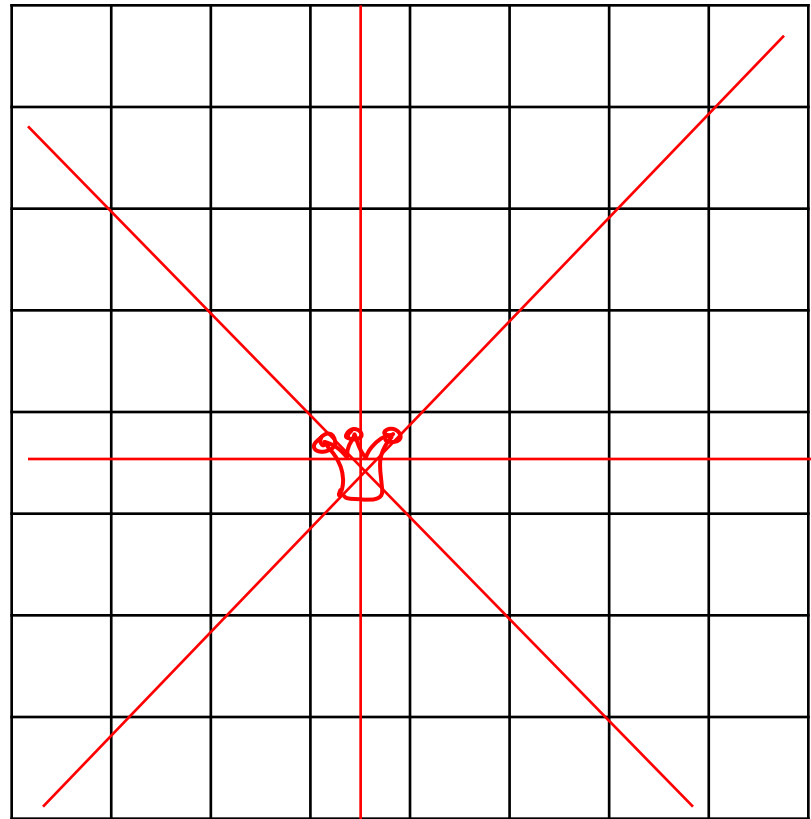
A common component of the scoring function (heuristic) => select the neighbor that results in the

- the **min conflicts** heuristics

Example: N-Queens

- Put n queens on an $n \times n$ board with **no two queens** on the same row, column, or diagonal (i.e attacking each other)

- Positions a queen can attack



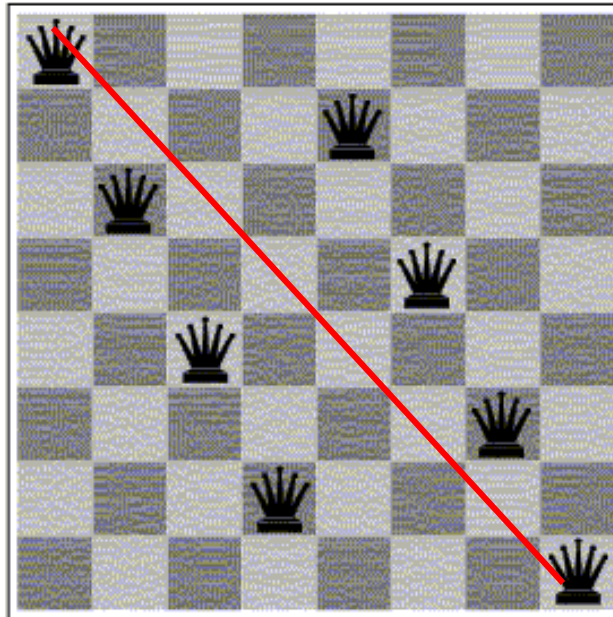
Example: N-queen as a local search problem

CSP: N-queen CSP

- One variable per column; domains $\{1, \dots, N\} \Rightarrow$ row where the queen in the i^{th} column seats;
- Constraints: no two queens in the same row, column or diagonal

Neighbour relation: value of a single column differs

Scoring function: number of constraint violations (i.e., number of



Example: n -queens

Put n queens on an $n \times n$ board with **no two queens** on the same row, column, or diagonal (i.e attacking each other)

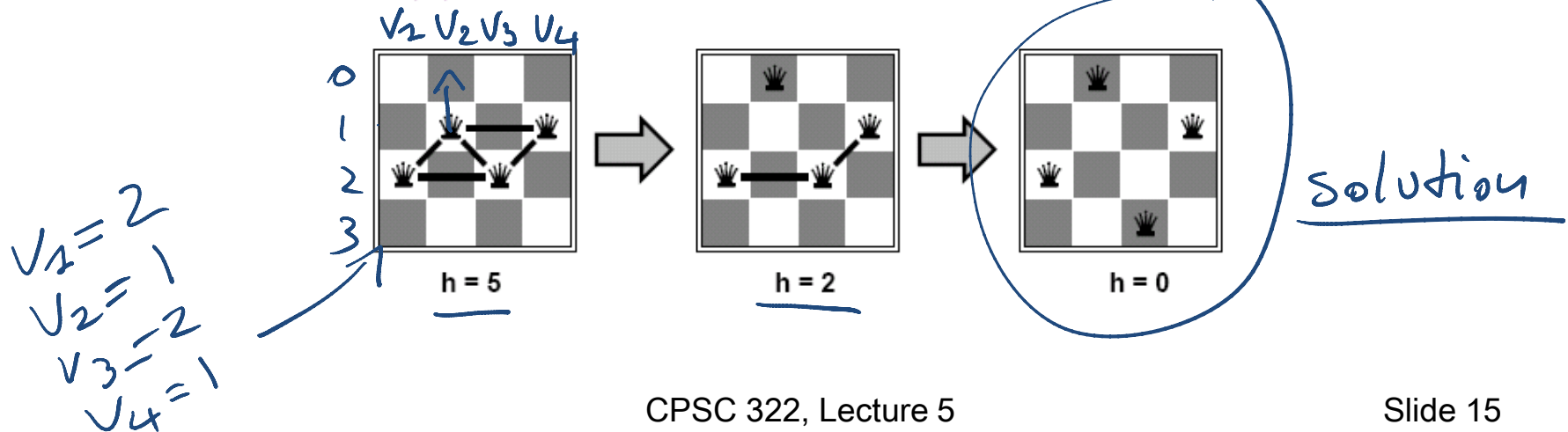
Example: 4-Queens

~~States~~: 4 queens in 4 columns ($4^4 = 256$ states)

Operators: move queen in column (to generate neighbors)

Goal test: no attacks

Evaluation: $h(n)$ = number of attacks



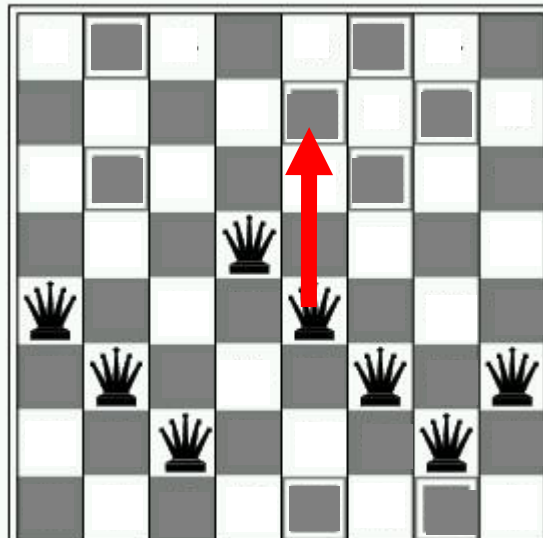
Example: Greedy descent for N-Queen

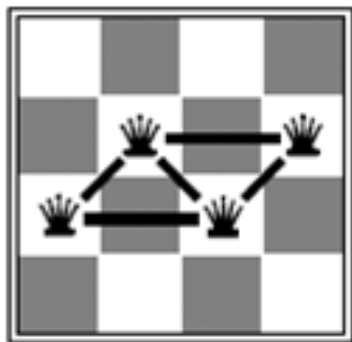
For each column, assign randomly each queen to a row
(a number between 1 and N)

Repeat

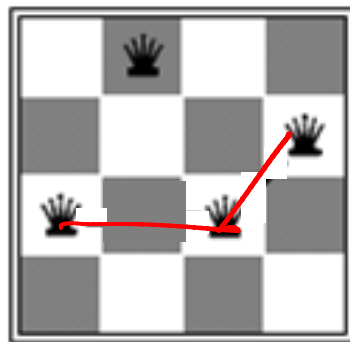
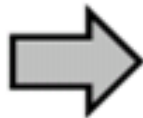
- For each column & each number: Evaluate how many constraint violations changing the assignment would yield
- Choose the column and number that leads to the fewest violated constraints; **change it**

Until solved

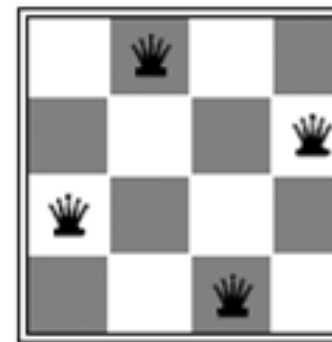




$h = 5$



$h = ?$



$h = ?$



n -queens, Why?



Why this problem?

Lots of research in the 90' on local search for CSP was generated by the observation that the run-time of local search on n -queens problems is **independent of problem size!**

Given random initial state, can solve n -queens in almost constant time for arbitrary n with high probability (e.g., $n = 10,000,000$)

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- Recap
- Local search
- **Constrained Optimization**
- Greedy Descent / Hill Climbing:
Problems

Constrained Optimization Problems

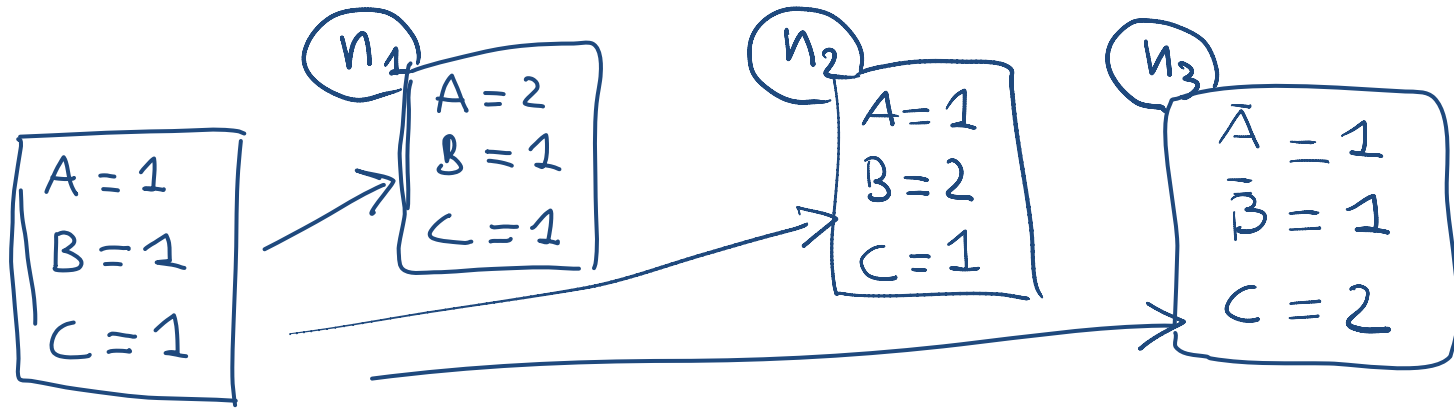
So far we have assumed that we just want to find a possible world that satisfies all the constraints.

But sometimes solutions may have different **values / costs**

- We want to find the **optimal solution** that
 - **maximizes the value** or
 - **minimizes the cost**

Constrained Optimization Example

- Example: A,B,C same domain {1,2,3} , (A=B, A>1, C≠3)
- Value = (C+A) so we want a solution that maximize that



The scoring function we'd like to maximize might be:

$$f(n) = (C + A) + \# \text{-of-satisfied-const}$$

$\begin{matrix} n_1 & n_2 & n_3 \\ (1+2)+2 & (1+1)+1 & (2+1)+2 \end{matrix}$

Hill Climbing means selecting the neighbor which best improves a (value-based) scoring function.

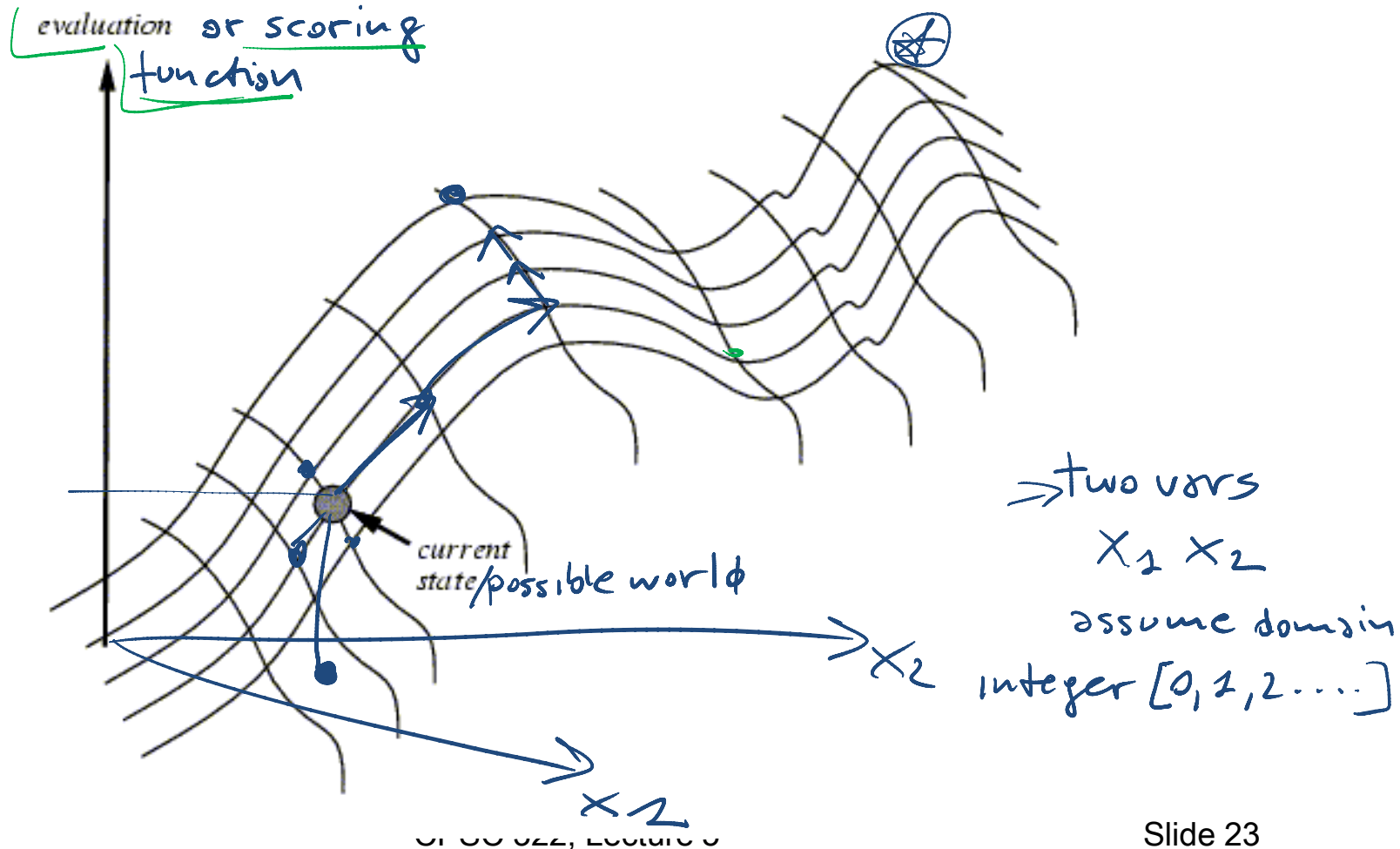
Greedy Descent means selecting the neighbor which minimizes a (cost-based) scoring function. *cost + # of conflicts*

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Hill Climbing

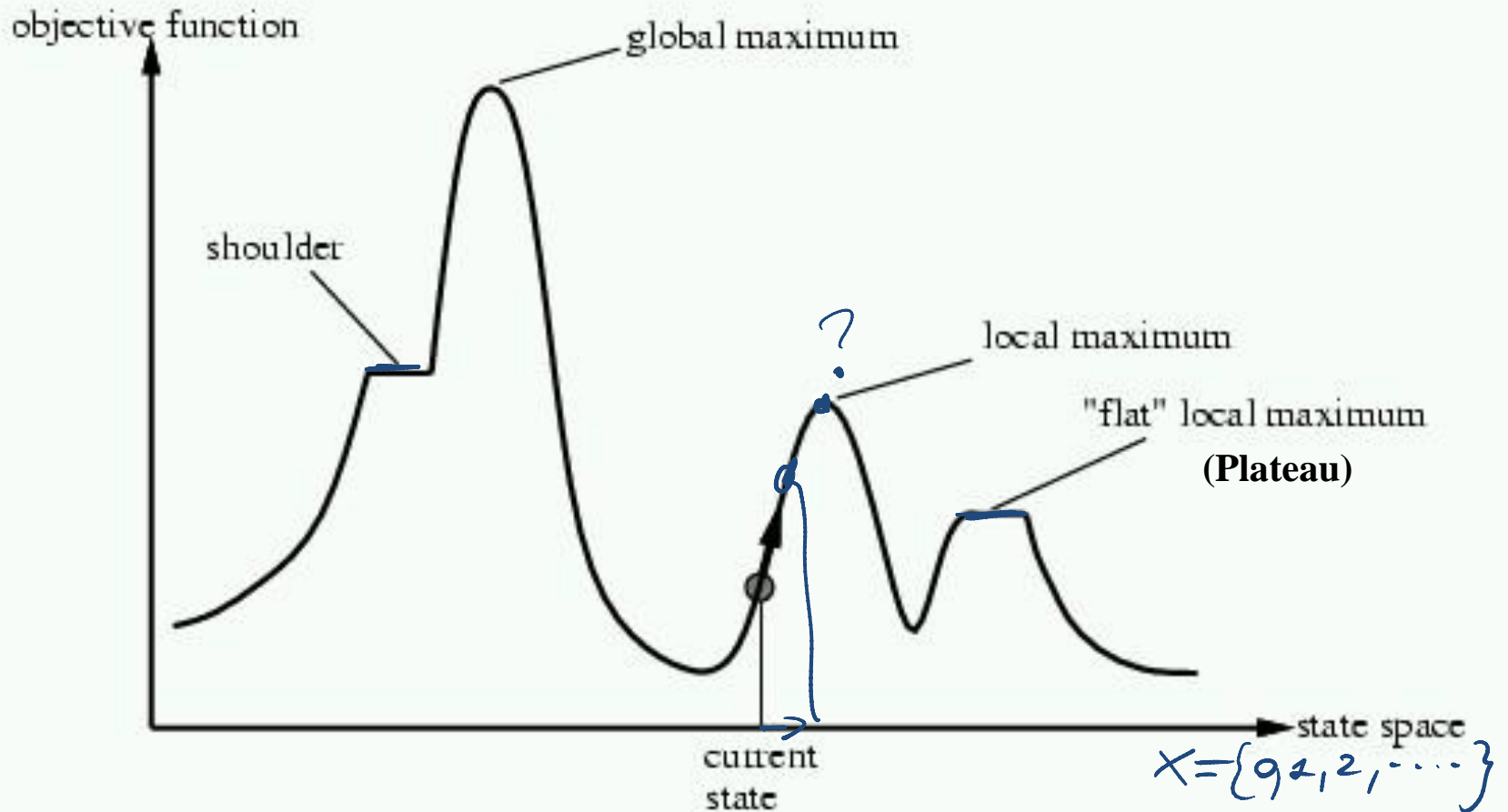
NOTE: Everything that will be said for Hill Climbing is also true for Greedy Descent



Problems with Hill Climbing

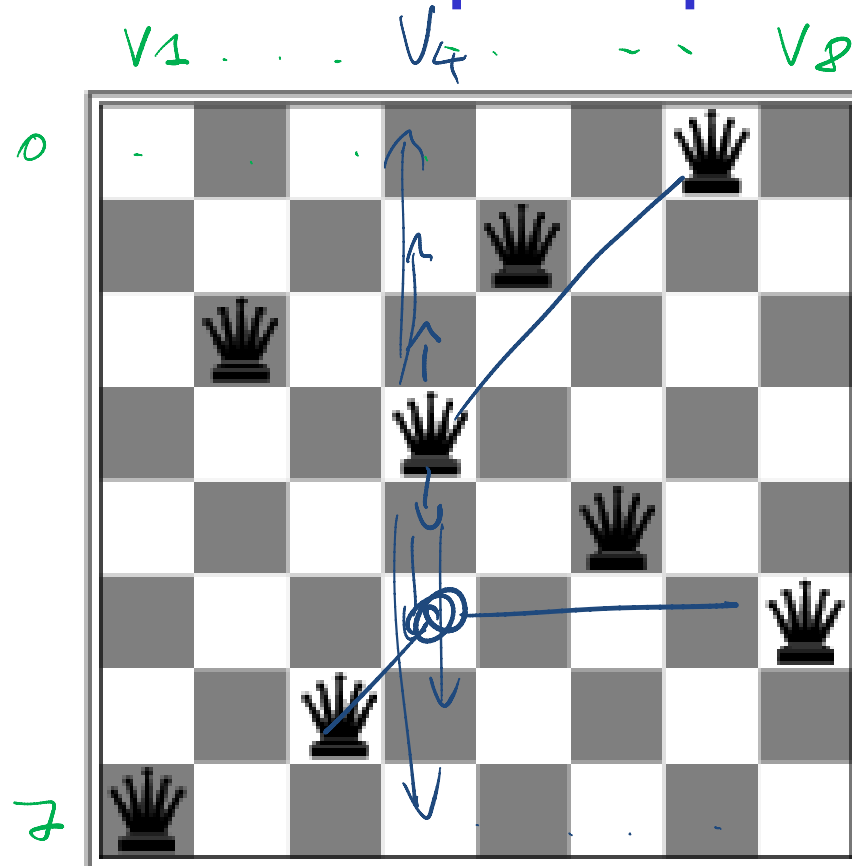
Local Maxima.

Plateau - Shoulders



Corresponding problem for GreedyDescent

Local minimum example: 8-queens problem



for all the
moves
(neighbors)
 $h > 1$

A local minimum with $h = 1$

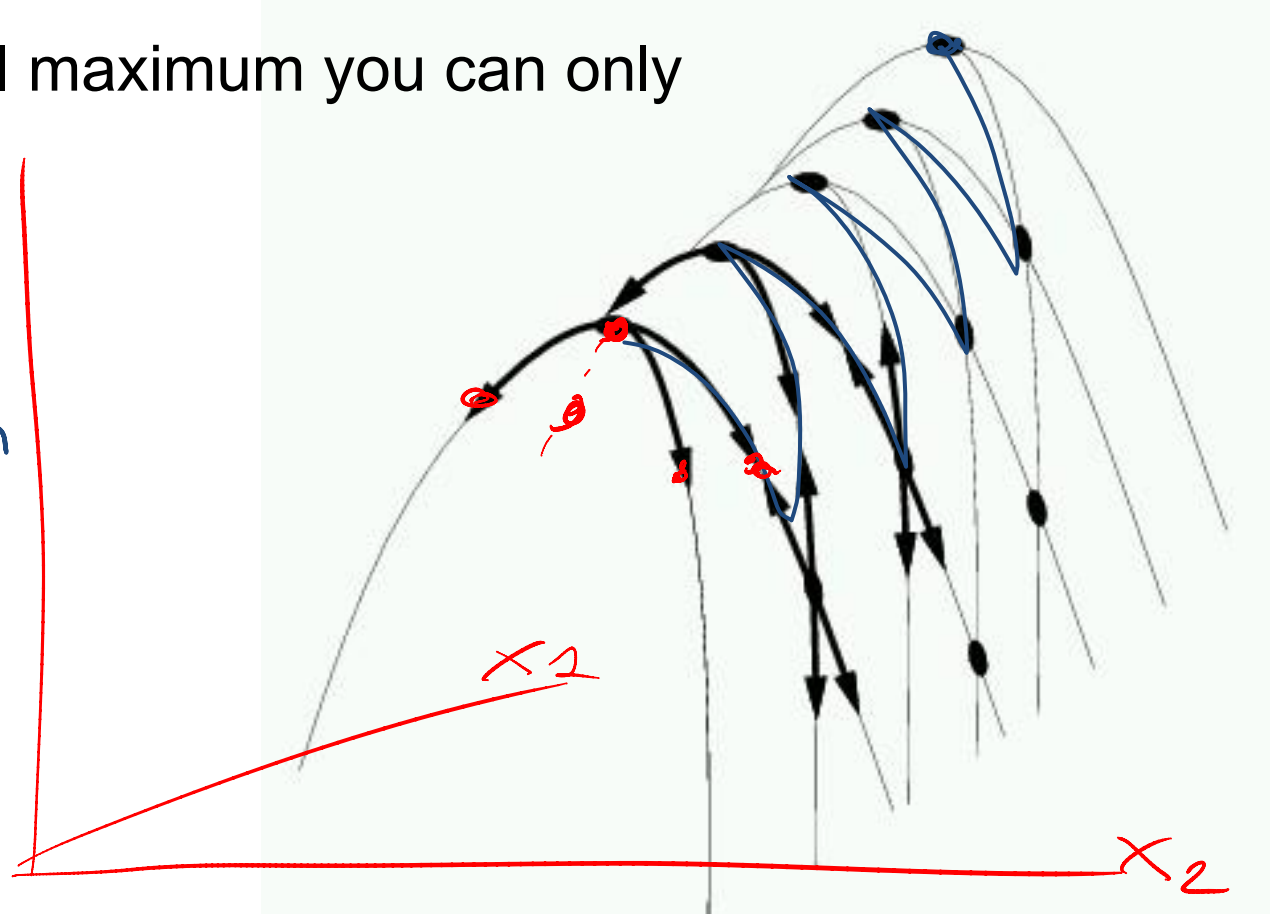
$h = 0$
for solution

Even more Problems in higher dimensions

E.g., Ridges – sequence of local maxima not directly connected to each other

From each local maximum you can only go downhill

scoring
function



Local Search: Summary

- A useful method for large CSPs
 - Start from a **possible world** (randomly chosen)
 - Generate some **neighbors** (“similar” possible worlds)
e.g. differ from current poss. world only by one variable's value
 - Move from current node to a neighbor, selected to minimize/maximize a scoring function which combines:
 - ✓ Info about how many constraints are violated
 - ✓ Information about the cost/quality of the solution (you want the best solution, not just a solution)

Learning Goals for today's class

You can:

- Implement **local search** for a CSP.
- Implement different ways to **generate neighbors**
- Implement **scoring functions** to solve a CSP by local search through either **greedy descent** or **hill-climbing**.

Next Class

- How to address problems with Greedy Descent / Hill Climbing?

Stochastic Local Search (SLS)

322 Feedback 😊 or 😞

- Lectures
- Slides
- Practice Exercises
- Assignments
- Alspace
-
- Textbook
- Course Topics / Objectives
- TAs
- Learning Goals
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