

Finish Search

Computer Science cpssc322, Lecture 10

(Textbook Chpt 3.6)

January, 25, 2010

Lecture Overview

- **Optimal Efficiency**
- Pruning Cycles and Repeated states Examples
- Dynamic Programming
- 8-puzzle Applet
- Search Recap

Optimal efficiency of A^*

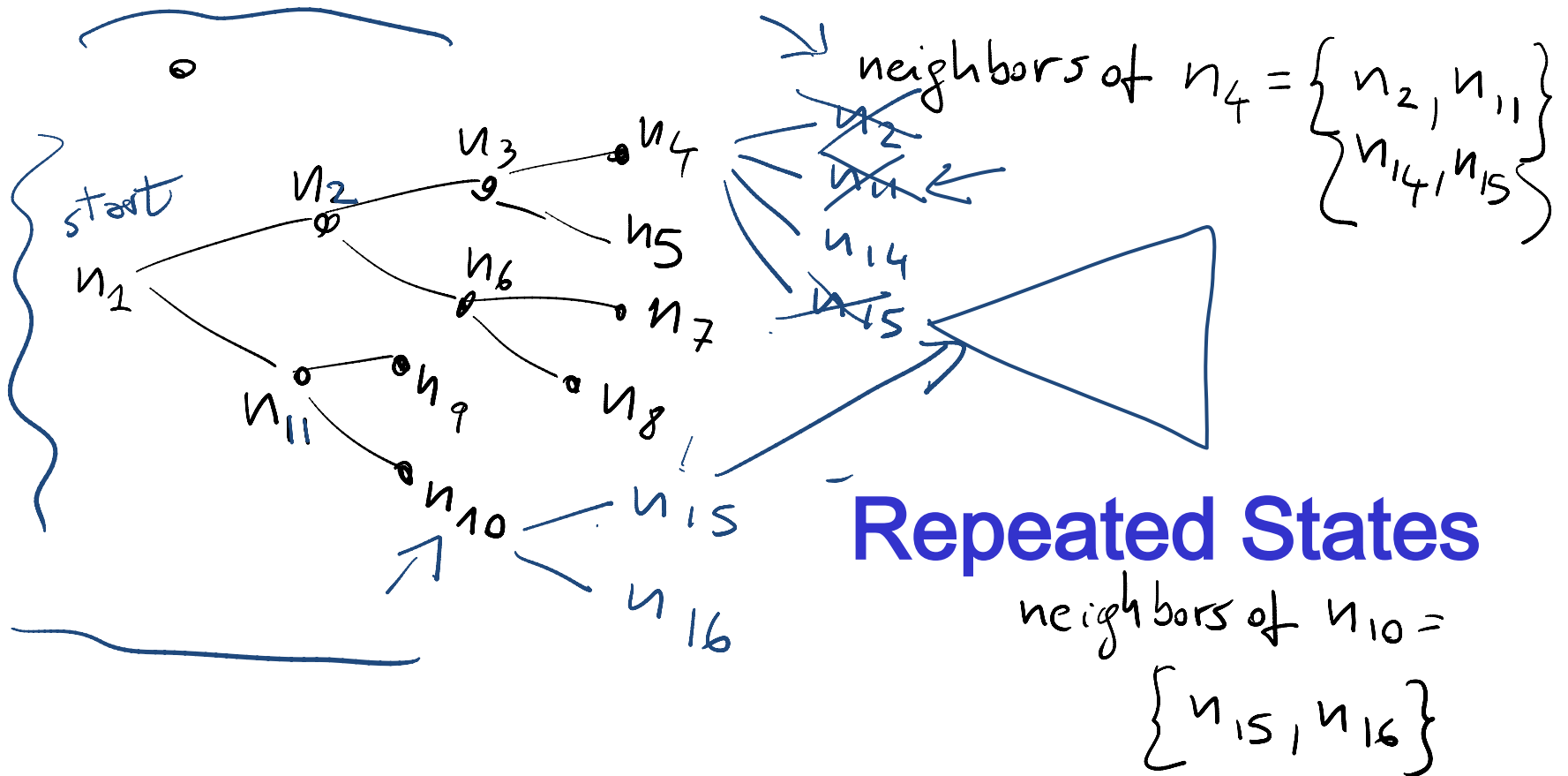
- In fact, we can prove something even stronger about A^* : in a sense (given the particular heuristic that is available) no search algorithm could do better!
- **Optimal Efficiency:** Among all optimal algorithms that start from the same start node and use the same heuristic h , A^* expands the minimal number of paths.

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Example

Pruning Cycles



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Dynamic Programming

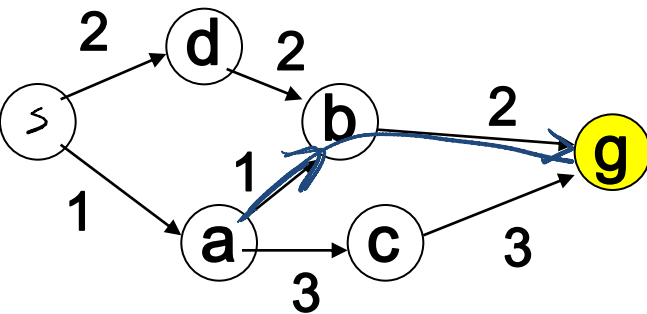
Idea: for statically stored graphs, build a table of $dist(n)$ the actual distance of the shortest path from node n to a goal. \leftarrow

This is the perfect... *heuristics*

This can be built **backwards** from the goal:

$$\underline{dist(n)} = \begin{cases} \underline{0} & \text{if } \underline{is_goal(n)}, \\ \min_{\langle n, m \rangle \in A} [\underline{cost(n, m)} + \underline{dist(m)}] & \text{otherwise} \end{cases}$$

all the neighbors m



$$dist(b) = \min[2 + 0] = 2$$

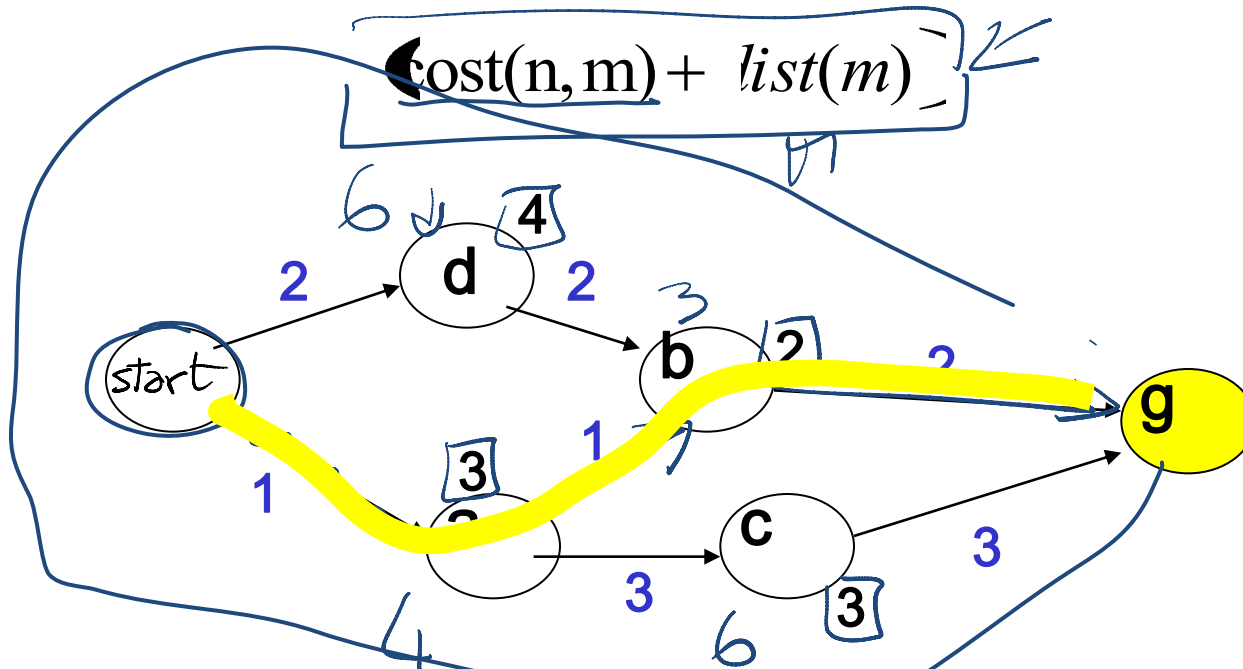
$$dist(c) = \min[(3 + 0)] = 3$$

$$dist(a) = \min[(3 + 3), (1 + 2)] = 2$$

Dynamic Programming

This can be used locally to determine what to do.

From each node n go to its neighbor which minimizes



But there are at least two main problems:

- You need enough space to store the graph.
- The *dist* function needs to be recomputed for each goal

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DFS, BFS, A* Animation Example

- The AI-Search animation system

<http://www.cs.rmit.edu.au/AI-Search/Product/>

- *To examine Search strategies when they are applied to the 8puzzle*
- Compare **only** DFS, BFS and A* (with only the two heuristics we saw in class)

AI-Search

Algorithm Problem Mode Settings Help

START

NEXT

BACK

PAUSE

RESET

slow

fast

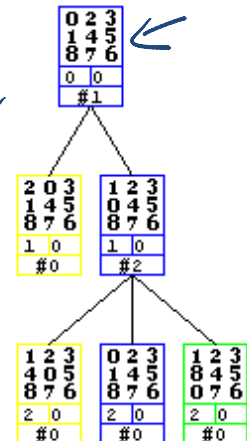
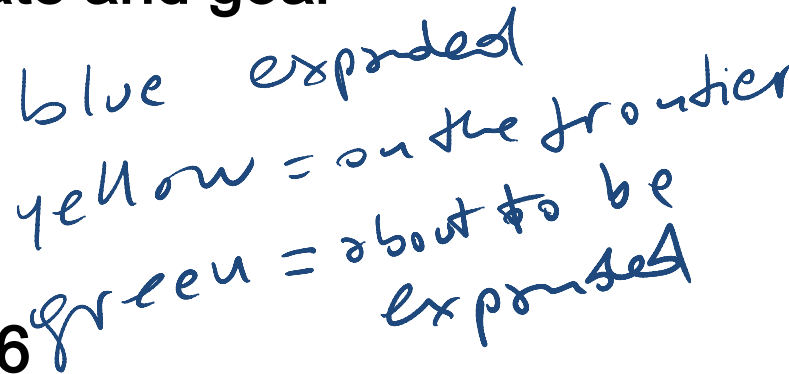
Open List
Closed List

- With default start state and goal

- DFS will find Solution at depth 32

- BFS will find
- Optimal solution depth 6

- **A* will also find opt. sol. expanding much less nodes**



n Puzzles are not always solvable

Half of the starting positions for the n -puzzle are impossible to resolve (for more info on 8puzzle) 

<http://www.isle.org/~sbay/ics171/project/unsolvable>

- So experiment with the AI-Search animation system with the default configurations.
- If you want to try new ones keep in mind that you may pick unsolvable problems

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Recap Search

| | Selection | Complete | Optimal | Time | Space |
|----------------|-----------------|----------|---------|----------|----------|
| DFS | LIFO | N | N | $O(b^m)$ | $O(mb)$ |
| BFS | FIFO | Y | Y | $O(b^m)$ | $O(b^m)$ |
| IDS(C) | LIFO | Y | Y | $O(b^m)$ | $O(mb)$ |
| LCFS | min cost | Y | Y | $O(b^m)$ | $O(b^m)$ |
| BFS | min h | N | N | $O(b^m)$ | $O(b^m)$ |
| A* | min $f = g + h$ | Y | Y | $O(b^m)$ | $O(b^m)$ |
| <u>B&B</u> | LIFO + pruning | N | Y | $O(b^m)$ | $O(mb)$ |
| <u>IDA*</u> | LIFO | Y | Y | $O(b^m)$ | $O(mb)$ |
| MBA* | min f | N | Y | $O(b^m)$ | $O(b^m)$ |

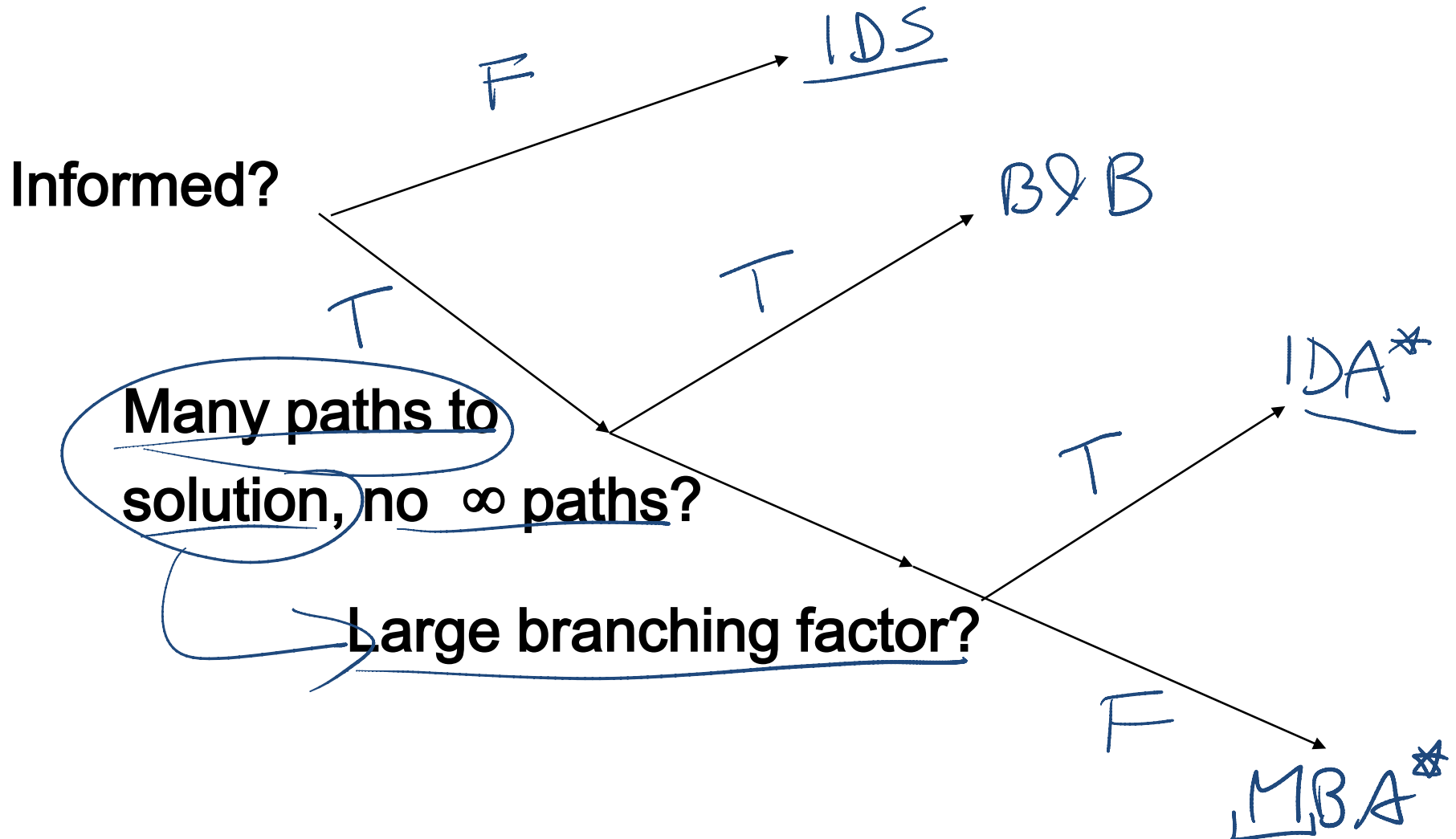
Recap Search (some qualifications)

| | Complete | Optimal | Time | Space |
|--------|----------|-------------------|----------|----------|
| DFS | N | N | $O(b^m)$ | $O(mb)$ |
| BFS | Y | Y | $O(b^m)$ | $O(b^m)$ |
| IDS(C) | Y | Y | $O(b^m)$ | $O(mb)$ |
| LCFS | Y | Y ? $C \geq 0$ | $O(b^m)$ | $O(b^m)$ |
| BFS | N | N | $O(b^m)$ | $O(b^m)$ |
| A* | Y | Y ? admissible | $O(b^m)$ | $O(b^m)$ |
| B&B | N | Y ? | $O(b^m)$ | $O(mb)$ |
| IDA* | Y | Y | $O(b^m)$ | $O(mb)$ |
| MBA* | N | Y | $O(b^m)$ | $O(b^m)$ |

Search in Practice

| | Complete | Optimal | Time | Space |
|----------------|----------|---------|--------------|---------------------------|
| DFS | N | N | $O(b^m)$ | $O(mb)$ |
| BFS | Y | Y | $O(b^m)$ | $O(b^m)$ |
| <u>IDS(C)</u> | → Y | → Y | $O(b^m)$ | <u>$O(mb)$</u> |
| LCFS | Y | Y | $O(b^m)$ | $O(b^m)$ |
| BFS | N | N | $O(b^m)$ | $O(b^m)$ |
| A* | Y | Y | $O(b^m)$ | $O(b^m)$ |
| <u>B&B</u> | N | Y | $O(b^m)$ | $O(mb)$ |
| <u>IDA*</u> | Y | Y | $O(b^m)$ | $O(mb)$ |
| <u>MBA*</u> | N | Y | $O(b^m)$ | $O(b^m)$ |
| BDS | Y | Y | $O(b^{m/2})$ | $O(b^{m/2})$ |

Search in Practice (cont')



➤ (Adversarial) Search: Chess

Deep Blue's Results in the second tournament:

- second tournament: won 3 games, lost 2, tied 1
- 30 CPUs + 480 chess processors
- Searched 126.000.000 nodes per sec
- Generated 30 billion positions per move reaching depth 14 routinely

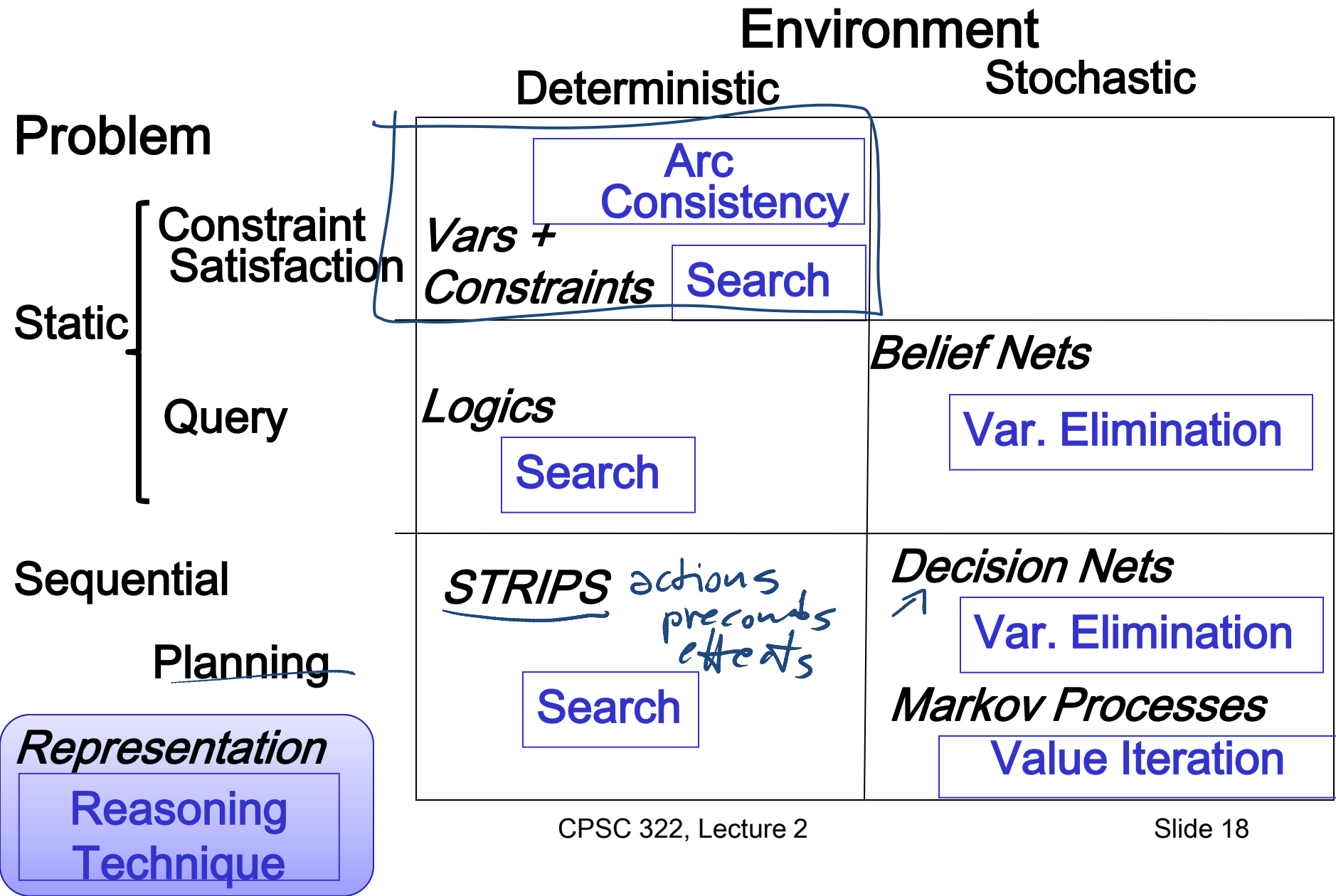
May 11th, 1997
Computer won world champion of chess
(Deep Blue) (Garry Kasparov)



(Reuters = Kyodo News)

- Iterative Deepening with evaluation function (similar to a heuristic) based on 8000 features (e.g., sum of worth of pieces: pawn 1, rook 5, queen 10)

Modules we'll cover in this course: R&Rsys



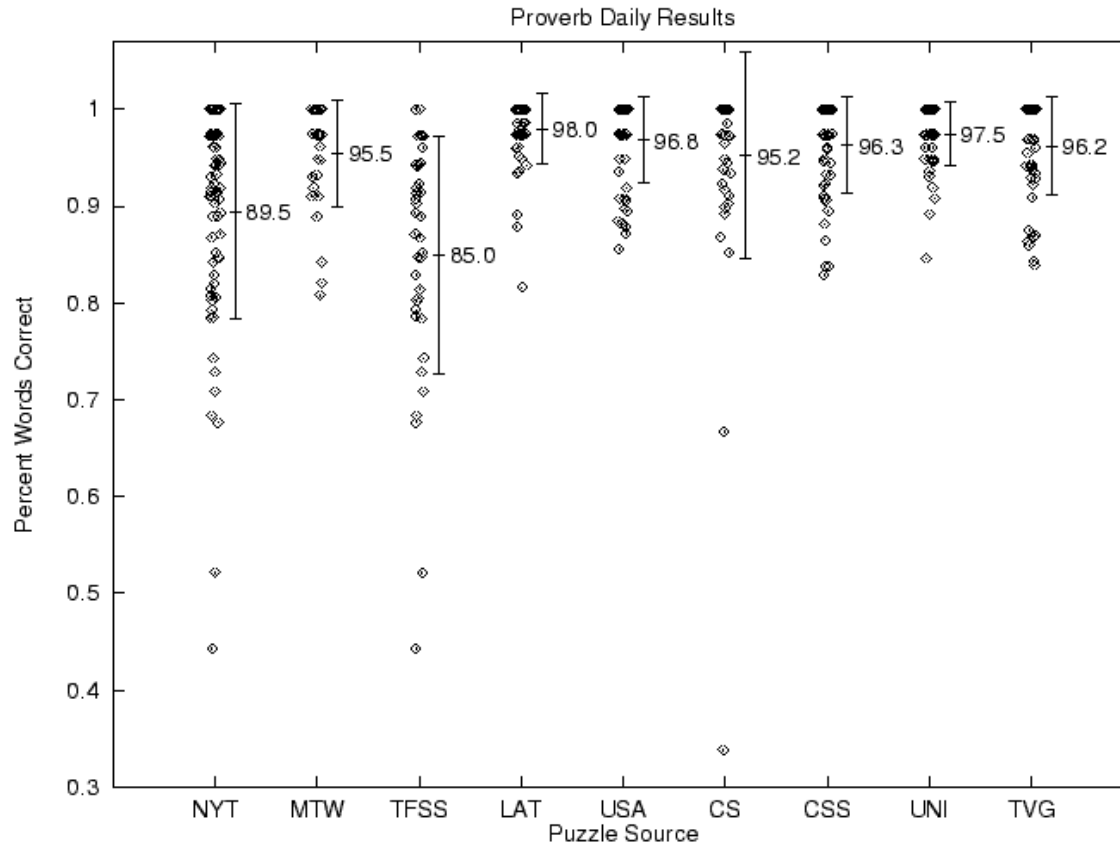
CSPs: Crossword Puzzles

Daily Puzzles

370 puzzles from 7 sources.

Summary statistics:

- ♦ 95.3% words correct (miss three or four words per puzzle)
- ♦ 98.1% letters correct
- ♦ 46.2% puzzles completely correct



| | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| P | O | L | O | N | E | | P | A | L | O | M | I | N | O |
| A | S | I | M | O | V | | I | S | O | L | A | T | E | D |
| S | L | E | E | V | E | | T | H | W | A | R | T | E | D |
| T | I | G | G | E | R | | C | O | R | N | Y | | | |
| A | N | E | A | L | E | | A | R | I | D | | J | A | M |
| | | | | | | E | S | P | I | E | S | | L | O |
| S | E | A | O | T | T | E | R | | E | E | N | O | N | |
| A | B | B | O | T | | A | N | A | | U | S | A | G | E |
| B | O | O | Z | E | S | | S | N | A | P | S | H | O | T |
| E | N | V | Y | | | P | L | I | N | T | H | | | |
| R | Y | E | | | | H | I | E | S | | T | E | A | S |
| | | | | | | K | A | R | E | L | | I | M | P |
| M | A | R | I | N | A | R | A | | | | M | I | A | S |
| A | B | E | R | D | E | E | N | | | | E | S | C | H |
| B | H | N | K | Y | A | R | D | | | | S | M | E | A |

Source: *Michael Littman*

CSPs: Radio link frequency assignment

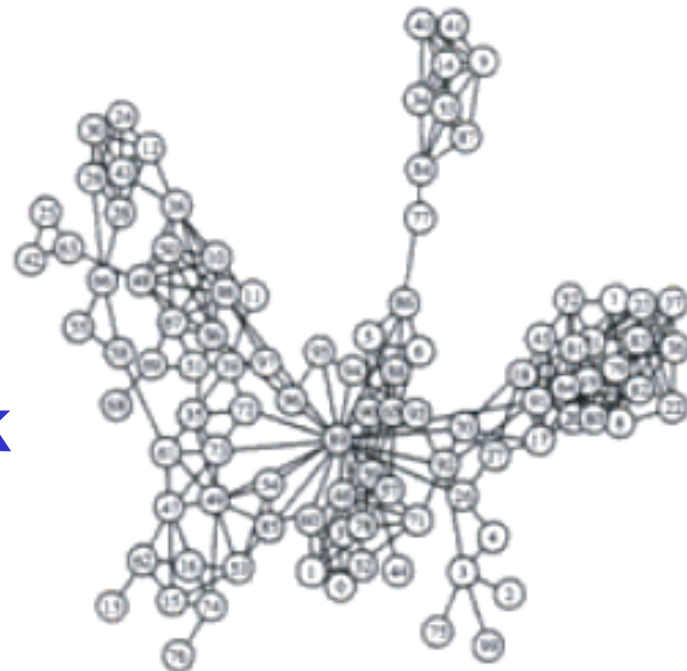
Assigning **frequencies** to a set of radio **links** defined between pairs of sites in order to avoid interferences.

Constraints on frequency depend on position of the links and on physical environment .

Source: *INRIA*

Sample Constraint network

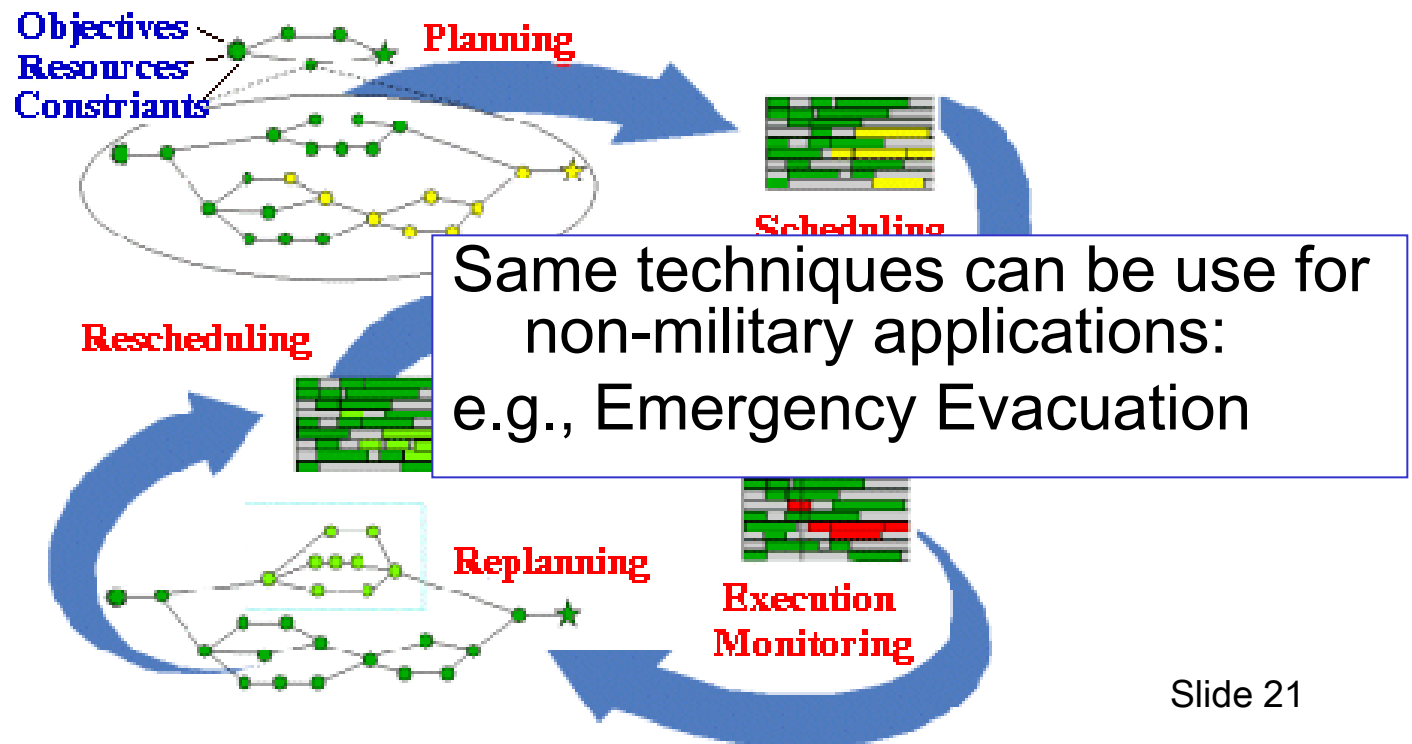
CPS



Planning & Scheduling: Logistics

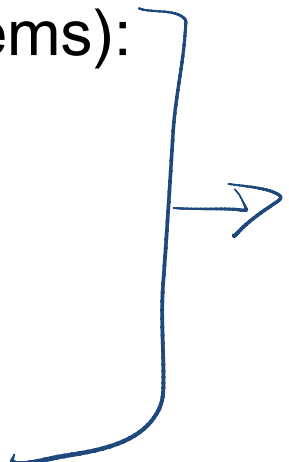
Dynamic Analysis and Replanning Tool (Cross & Walker)

- logistics planning and scheduling for military transport
- used for the first time in the 1991 Gulf War by the US
- problems had 50,000 entities (e.g., vehicles); different starting points and destinations



Standard Search vs. Specific R&R systems

Constraint Satisfaction (Problems):

- State
 - Successor function
 - Goal test
 - Solution
 - Heuristic function
- 

Planning :

- State
- Successor function
- Goal test
- Solution
- Heuristic function

Inference

- State
- Successor function
- Goal test
- Solution
- Heuristic function

Next class

Start Constraint Satisfaction Problems (CSPs)

Textbook 4.1-4.3