416 Distributed Systems

Networks review
Sep 11, 2018
Distributed Systems vs. Networks

• Low level (c/go)
• Run forever
• Support others
• Adversarial environment
• Distributed & concurrent
• Resources matter

• And have it implemented/run by vast numbers of different people with different goals/skills
Keep an eye out for…

- **Modularity, Layering, and Decomposition:**
  - Techniques for dividing the work of building systems
  - Hiding the complexity of components from each other
  - Hiding implementation details to deal with heterogeneity

- **Naming/lookup/routing**

- **Resource sharing and isolation**

- **Models and assumptions about the environment and components**
Today’s Lecture (317 review-ish)

- Network links and LANs
- Layering and protocols
- Internet design
Basic Building Block: Links

- Electrical questions
  - Voltage, frequency, …
  - Wired or wireless?

- Link-layer issues: How to send data?
  - When to talk – can either side talk at once?
  - What/how to say – low-level format?
Model of a communication channel

- Latency - how long does it take for the first bit to reach destination
- Jitter - how much variation in latency?
- Capacity - how many bits/sec can we push through? (often termed “bandwidth”)
- Loss / Reliability - can the channel drop packets?
- Reordering
Basic Building Block: Links

• But what if we want more hosts?

• Scalability?!

One wire

Wires for everybody!
Multiplexing

- Need to share network resources

- How? Switched network
  - Party “A” gets resources sometimes
  - Party “B” gets them sometimes

- Interior nodes act as “Switches”

- What mechanisms to share resources?
In the Old Days… Circuit Switching
Packet Switching

- Source sends information as self-contained packets that have an address.
  - Source may have to break up single message in multiple packets.

- Each packet travels independently to the destination host.
  - Switches use the address in the packet to determine how to forward the packets.
  - Store and forward of packets.

- Analogy: a letter in surface mail (snail mail).
Packet Switching – Statistical Multiplexing

- Switches arbitrate between inputs
- Can send from *any* input that’s ready
  - Links never idle when traffic to send
  - (Efficiency!)
What if Network is Overloaded?

Problem: Network Overload

- Short bursts: buffer
- What if buffer overflows?
  - Packets dropped
  - Sender adjusts rate until load = resources → “congestion control”

Solution: Buffering and Congestion Control
Example: Ethernet Packet

- Sending adapter **encapsulates** IP datagram (or other network layer protocol packet) in Ethernet frame
Each protocol layer needs to provide some hooks to upper layer protocols

- Demultiplexing: identify which upper layer protocol packet belongs to
- E.g., port numbers allow TCP/UDP to identify target application
- Ethernet uses Type field

**Type:** 2 bytes

- Indicates the higher layer protocol, mostly IP but others may be supported such as Novell IPX and AppleTalk
Ethernet Frame Structure (cont.)

- **Addresses**: 6 bytes
  - Each adapter is given a globally unique address at manufacturing time
    - Address space is allocated to manufacturers
      - 24 bits identify manufacturer
      - E.g., 0:0:15:* → 3com adapter
    - Frame is received by all adapters on a LAN and dropped if address does not match
  - Special addresses
    - Broadcast – FF:FF:FF:FF:FF:FF is “everybody”
    - Range of addresses allocated to multicast
      - Adapter maintains list of multicast groups node is interested in
Packet Switching

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Frame Forwarding

- A machine with MAC Address lies in the direction of number port of the bridge
- For every packet, the bridge “looks up” the entry for the packet’s destination MAC address and forwards the packet on that port.
  - Other packets are broadcast – why?
- Timer is used to flush old entries
Learning Bridges

- Manually filling in bridge tables?
  - Time consuming, error-prone
- Keep track of source address of packets arriving on every link, showing what segment hosts are on
  - Fill in the forwarding table based on this information
Today’s Lecture

• Network links and LANs
• Layering and protocols
• Internet design
Internet

• An inter-net: a network of networks.
  • Networks are connected using routers that support communication in a hierarchical fashion
  • Often need other special devices at the boundaries for security, accounting, ..

• The Internet: the interconnected set of networks of the Internet Service Providers (ISPs)
  • About 17,000 different networks make up the Internet
Challenges of an internet

- Heterogeneity
- Address formats
- Performance – bandwidth/latency
- Packet size
- Loss rate/pattern/handling
- Routing
- Diverse network technologies → satellite links, cellular links, carrier pigeons
- In-order delivery
How To Find Nodes?

Computer 1

Internet

Need naming and routing

Computer 2
Naming

What’s the IP address for www.cmu.edu?
It is 128.2.11.43

Computer 1

Local DNS Server

Translates human readable names to logical endpoints
Routing

Routers send packet towards destination

H: Hosts
R: Routers
What is the base (IP) service model?

- Ethernet/Internet: *best-effort* – packets can get lost, etc.

What if you want more?

- Performance guarantees (QoS)
- Reliability
  - Corruption
  - Lost packets
- Flow and congestion control
- Fragmentation
- In-order delivery
- Etc…
Aside: failure models

- Fail-stop:
  - When something goes wrong, the process stops / crashes / etc.

- Fail-slow or fail-stutter:
  - Performance may vary on failures as well

- Byzantine:
  - Anything that can go wrong, will.
  - Including malicious entities taking over your computers and making them do whatever they want.

- These models are useful for proving things;
- The real world typically has a bit of everything.

- Deciding which model to use is important!
Fancier Network Service Models

• What if network had reliable, in-order, mostly no-corruption, stream-oriented communication (i.e. TCP)

• Programmers don’t have to implement these features in every application

• But note limitations: this can’t turn a byzantine failure model into a fail-stop model...
What if the Data gets Corrupted?

Problem: Data Corruption

GET index.html → Internet → GET inrex.html

Solution: Add a checksum

0,9 9 6,7,8 21 4,5 7 1,2,3 6

X

Internet
What if the Data gets Lost?

Problem: Lost Data

Internet

GET index.html

Solution: Timeout and Retransmit

GET index.html

Internet

GET index.html
What if the Data is Out of Order?

Problem: Out of Order

ml → inde → x.ht → GET

Solution: Add Sequence Numbers

ml → 4 → inde → 2 → x.ht → 3 → GET → 1

GET index.html
Networks [including end points] Implement Many Functions

- Link
- Multiplexing
- Routing
- Addressing/naming (locating peers)
- Reliability
- Flow control
- Fragmentation
- Etc….
What is Layering?

- Modular approach to network functionality
- Example:

```
Application
Application-to-application channels
Host-to-host connectivity
Link hardware
```
What is Layering?

Modular approach to network functionality
Layering Characteristics

• Each layer relies on services from layer below and exports services to layer above
• Interface defines interaction with peer on other hosts
• Hides implementation - layers can change without disturbing other layers (black box)
What are Protocols?

- An agreement between parties on how communication should take place

- Module in layered structure

- Protocols define:
  - Interface to higher layers (API)
  - Interface to peer (syntax & semantics)
    - Actions taken on receipt of a messages
    - Format and order of messages
    - Error handling, termination, ordering of requests, etc.

- Example: Buying airline ticket

Friendly greeting

Muttered reply

Destination?

Honolulu

Thank you
IP Layering (control flow)

- Relatively simple

Diagram showing the layering of IP and its components (Application, Transport, Network, Link, Physical) and the roles of different devices (Host, Bridge/Switch, Router/Gateway).
The Internet Protocol Suite

The "thin" waist facilitates interoperability
Layer Encapsulation (data flow)
Multiplexing and Demultiplexing

- There may be multiple implementations of each layer.
  - How does the receiver know what version of a layer to use?
- Each header includes a demultiplexing field that is used to identify the next layer.
  - Filled in by the sender
  - Used by the receiver
- Multiplexing occurs at multiple layers. E.g., IP, TCP, …
Multiplexing and Demultiplexing

There may be multiple implementations of each layer. How does the receiver know what version of a layer to use? Each header includes a demultiplexing field that is used to identify the next layer. Filled in by the sender, used by the receiver.

Multiplexing occurs at multiple layers. E.g., IP, TCP, …

List of IP protocol numbers

From Wikipedia, the free encyclopedia

<table>
<thead>
<tr>
<th>Decimal</th>
<th>Hex</th>
<th>Keyword</th>
</tr>
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<tbody>
<tr>
<td>0</td>
<td>0x00</td>
<td>HOPOPT</td>
</tr>
<tr>
<td>1</td>
<td>0x01</td>
<td>ICMP</td>
</tr>
<tr>
<td>2</td>
<td>0x02</td>
<td>IGMP</td>
</tr>
<tr>
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<td>GGP</td>
</tr>
<tr>
<td>4</td>
<td>0x04</td>
<td>IP-in-IP</td>
</tr>
<tr>
<td>5</td>
<td>0x05</td>
<td>ST</td>
</tr>
<tr>
<td>6</td>
<td>0x06</td>
<td>TCP</td>
</tr>
<tr>
<td>7</td>
<td>0x07</td>
<td>CBT</td>
</tr>
<tr>
<td>8</td>
<td>0x08</td>
<td>EGP</td>
</tr>
<tr>
<td>9</td>
<td>0x09</td>
<td>IGP</td>
</tr>
<tr>
<td>10</td>
<td>0x0A</td>
<td>BBN-RCC-MON</td>
</tr>
<tr>
<td>11</td>
<td>0x0B</td>
<td>NVP-II</td>
</tr>
<tr>
<td>12</td>
<td>0x0C</td>
<td>PUP</td>
</tr>
<tr>
<td>13</td>
<td>0x0D</td>
<td>ARGUS</td>
</tr>
<tr>
<td>14</td>
<td>0x0E</td>
<td>EMSS</td>
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V/HL | TOS | Length | Flags/Offset | H. Checksum | Source IP address | Destination IP address | Options..
Protocol Demultiplexing

- Multiple choices at each layer
Today’s Lecture

- Network links and LANs
- Layering and protocols
- Internet design
0. Connect existing networks
   initially ARPANET and ARPA packet radio network

1. Survivability
   ensure communication service even in the presence of
   network and router failures

2. Support multiple types of services

3. Must accommodate a variety of networks

4. Allow distributed management

5. Allow host attachment with a low level of effort

6. Be cost effective

7. Allow resource accountability
Goal 1: Survivability

- If network is disrupted and reconfigured…
  - Communicating entities should not care!
  - No higher-level state reconfiguration

- How to achieve such reliability?
  - Where can communication state be stored?

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<td>Simple</td>
</tr>
<tr>
<td>Routing state</td>
<td>Maintain state</td>
<td>Stateless</td>
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Fate Sharing

• Definition: *lose state information for an entity if and only if the entity itself is lost.*

• Examples:
  - OK to lose TCP state if one endpoint crashes
    - NOT okay to lose if an intermediate router reboots
  - Is this still true in today’s network?
    - NATs and firewalls

• Tradeoffs
  - Less information available to the network
  - Must trust endpoints more
Networks [including end points] Implement Many Functions

- Link
- Multiplexing
- Routing
- Addressing/naming (locating peers)
- **Reliability**
- Flow control
- Fragmentation
- Etc….
Design Question

- If you want reliability, where should you implement it?

Option 1: Hop-by-hop (at switches)

Option 2: end-to-end (at end-hosts)
Options

- **Hop-by-hop**: Have each switch/router along the path ensure that the packet gets to the next hop
- **End-to-end**: Have just the end-hosts ensure that the packet made it through
- What do we have to think about to make this decision??
A question

• Is hop-by-hop enough?

• [hint: What happens if a switch crashes? What if it’s buggy and goofs up a packet?]
End-to-End Argument

• Deals with where to place functionality
  • Inside the network (in switching elements)
  • At the edges

• Guideline not a law

• Argument
  • If you have to implement a function end-to-end anyway (e.g., because it requires the knowledge and help of the end-point host or application), don’t implement it inside the communication system
  • Unless there’s a compelling performance enhancement

Further Reading: “End-to-End Arguments in System Design.” Saltzer, Reed, and Clark.
Questions to ponder

• If you have a whole file to transmit, how do you send it over the Internet?
  • You break it into packets (packet-switched medium)
  • TCP, roughly speaking, has the sender tell the receiver “got it!” every time it gets a packet. The sender uses this to make sure that the data’s getting through.
  • But by e2e, if you have to acknowledge the correct receipt of the entire file... why bother acknowledging the receipt of the individual packets???
Questions to ponder

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  • But by e2e, if you have to acknowledge the correct receipt of the entire file... why bother acknowledging the receipt of the individual packets???

• The answer: if you want performance, then you better do it this way (a mixture of e2e and in-network); imagine the waste if you had to retransmit the entire file because one packet was lost!
Internet Design: Types of Service

- **Principle**: network layer provides one simple service: best effort datagram (packet) delivery
  - All packets are treated the same

- Relatively simple core network elements
- Building block from which other services (such as reliable data stream) can be built
- Contributes to scalability of network

- No QoS support assumed from below
  - In fact, some underlying nets (e.g., link/physical layer) only provide reliable delivery (not best effort)
    - This made Internet datagram service less useful!
  - Hard to implement QoS without network support
  - QoS is an ongoing debate…
User Datagram Protocol (UDP): An Analogy

<table>
<thead>
<tr>
<th>UDP</th>
<th>Postal Mail</th>
</tr>
</thead>
<tbody>
<tr>
<td>Single socket to receive messages</td>
<td>Single mailbox to receive letters</td>
</tr>
<tr>
<td>No guarantee of delivery</td>
<td>Unreliable ☹</td>
</tr>
<tr>
<td>Not necessarily in-order delivery</td>
<td>Not necessarily in-order delivery</td>
</tr>
<tr>
<td>Datagram – independent packets</td>
<td>Letters sent independently</td>
</tr>
<tr>
<td>Must address each packet</td>
<td>Must address each letter</td>
</tr>
</tbody>
</table>

Example UDP applications
Multimedia, voice over IP
Transmission Control Protocol (TCP): An Analogy

TCP
- Reliable – guarantee delivery
- Byte stream – in-order delivery
- Connection-oriented – single socket per connection
- Setup connection followed by data transfer

Telephone Call
- Guaranteed delivery
- In-order delivery
- Connection-oriented
- Setup connection followed by conversation

Example TCP applications
Web, Email, Telnet
Why not always use TCP?

• TCP provides “more” than UDP
• Why not use it for everything??

• A: Nothing comes for free...
  • Connection setup (take on faith) -- TCP requires one round-trip time to setup the connection state before it can chat...
  • How long does it take, using TCP, to fix a lost packet?
    • At minimum, one “round-trip time” (2x the latency of the network)
    • That could be 100+ milliseconds!
  • If I guarantee in-order delivery, what happens if I lose one packet in a stream of packets?
  • Has semantics that may be too strong for the app (e.g., Netflix streaming)
Design trade-off

• If you’re building an app...
  • Do you need everything TCP provides?
  • If not:
    • Can you deal with its drawbacks to take advantage of the subset of its features you need?
    OR
    • You’re going to have to implement the ones you need on top of UDP
      • Caveat: There are some libraries, protocols, etc., that can help provide a middle ground.
      • Takes some looking around
Socket API Operation Overview

Client

- `socket`
- `connect` (Client)
- `write`
- `read` (Client)
- `close`

Server

- `socket`
- `bind`
- `listen`
- `accept`
- `read`
- `write`
- `close`

Connection request: `open_listenfd` from Server to Client

Client / Server Session

- `open_clientfd`
- `read` (Client)
- `write` (Client)
- `read` (Server)
- `write` (Server)
- `close`
- `EOF`
Blocking sockets

- What happens if an application write()s to a socket waaaaay faster than the network can send the data?
- TCP figures out how fast to send the data...
- And it builds up in the kernel socket buffers at the sender... and builds...
- until they fill. The next write() call blocks (by default).
- What’s blocking? It suspends execution of the blocked thread until enough space frees up...
In contrast to UDP

- UDP doesn’t figure out how fast to send data, or make it reliable, etc.
- So if you write() like mad to a UDP socket...
- It often silently disappears. *Maybe* if you’re lucky the write() call will return an error. But no promises.
Summary: Internet Architecture

- Packet-switched datagram network
- IP is the “compatibility layer”
  - Hourglass architecture
  - All hosts and routers run IP
- Stateless architecture
  - *no per flow state inside network*
Summary: Minimalist Approach

• Dumb network
  • IP provide minimal functionalities to support connectivity
    • Addressing, forwarding, routing

• Smart end system
  • Transport layer or application performs more sophisticated functionalities
    • Flow control, error control, congestion control

• Advantages
  • Accommodate heterogeneous technologies (Ethernet, modem, satellite, wireless)
  • Support diverse applications (telnet, ftp, Web, X windows)
  • Decentralized network administration