

Single agent or multiple agents

- Many domains are characterized by multiple agents rather than a single agent.
- **Game theory** studies what agents should do in a multi-agent setting.
- Agents can be cooperative, competitive or somewhere in between.
- Agents that are strategic can't be modeled as nature.

Multi-agent framework

- Each agent can have its own values.
- Agents select actions autonomously.
- Agents can have different information.
- The outcome can depend on the actions of all of the agents.
- Each agent's value depends on the outcome.

Fully Observable + Multiple Agents

- If agents act sequentially and can observe the state before acting: **Perfect Information Games.**
- Can do dynamic programming or search:
Each agent maximizes for itself.
- Multi-agent MDPs: value function for each agent.
each agent maximizes its own value function.
- Multi-agent reinforcement learning: each agent has its own Q function.
- Two person, competitive (zero sum) \implies minimax.

Normal Form of a Game

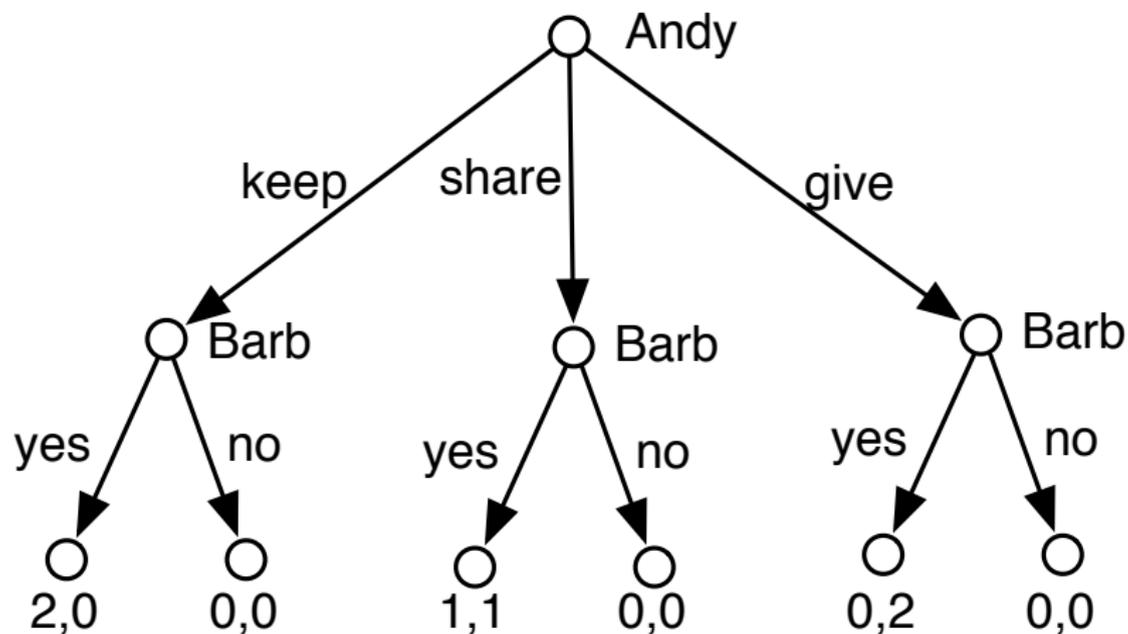
The **strategic form of a game** or **normal-form game**:

- a finite set I of agents, $\{1, \dots, n\}$.
- a set of actions A_i for each agent $i \in I$.
An **action profile** σ is a tuple $\langle a_1, \dots, a_n \rangle$, means agent i carries out a_i .
- a utility function $utility(\sigma, i)$ for action profile σ and agent $i \in I$, gives the expected utility for agent i when all agents follow action profile σ .

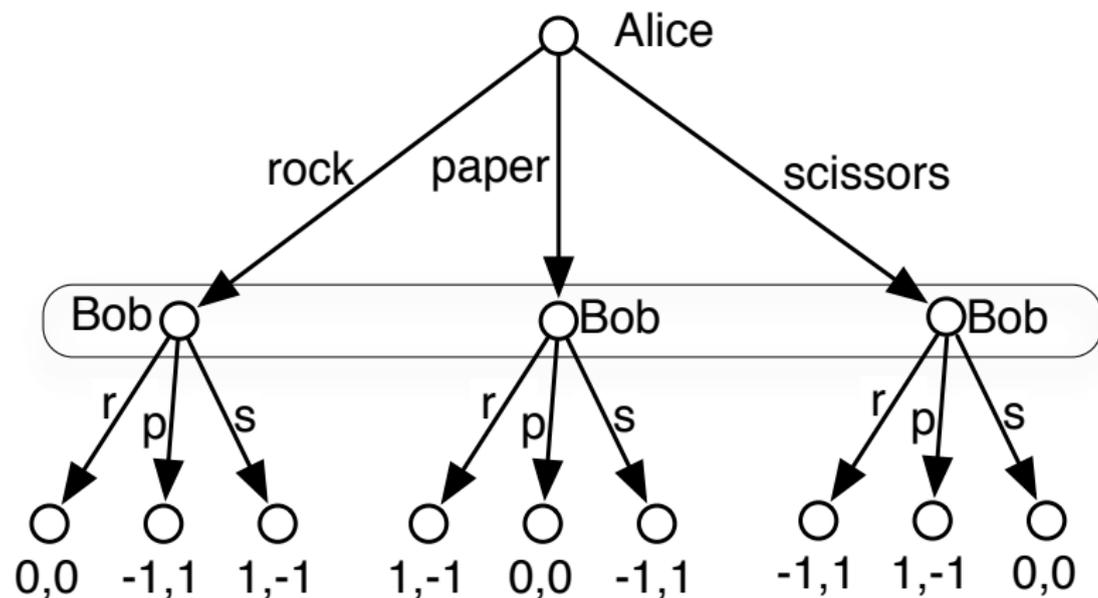
Rock-Paper-Scissors

		Bob		
		<i>rock</i>	<i>paper</i>	<i>scissors</i>
Alice	<i>rock</i>	0, 0	-1, 1	1, -1
	<i>paper</i>	1, -1	0, 0	-1, 1
	<i>scissors</i>	-1, 1	1, -1	0, 0

Extensive Form of a Game

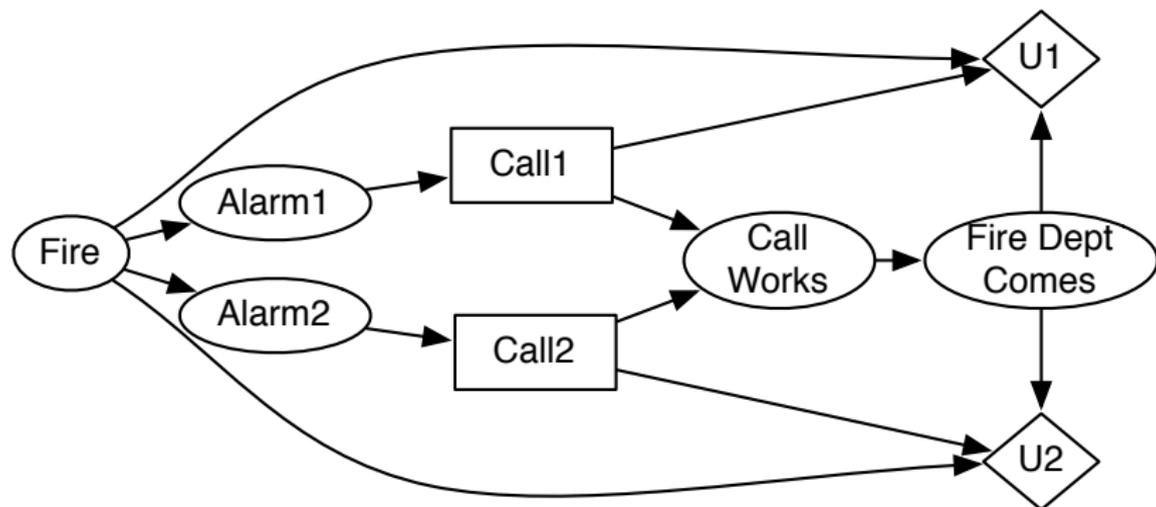


Extensive Form of an imperfect-information Game



Bob cannot distinguish the nodes in an **information set**.

Multiagent Decision Networks

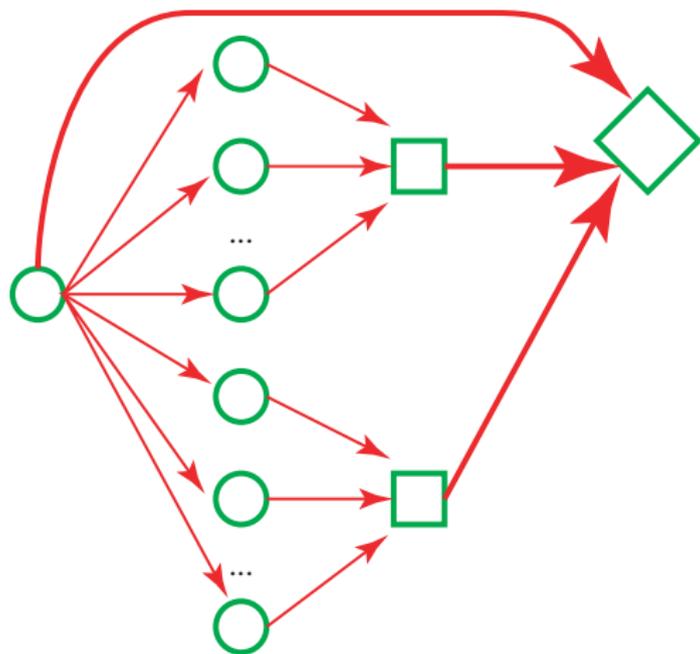


Value node for each agent.

Each decision node is owned by an agent.

Utility for each agent.

Multiple Agents, shared value



Complexity of Multi-agent decision theory

- It can be exponentially harder to find optimal multi-agent policy even with a shared values.
- **Why?** Because dynamic programming doesn't work:
 - ▶ If a decision node has n binary parents, dynamic programming lets us solve 2^n decision problems.
 - ▶ This is much better than d^{2^n} policies (where d is the number of decision alternatives).
- Multiple agents with shared values is equivalent to having a single forgetful agent.

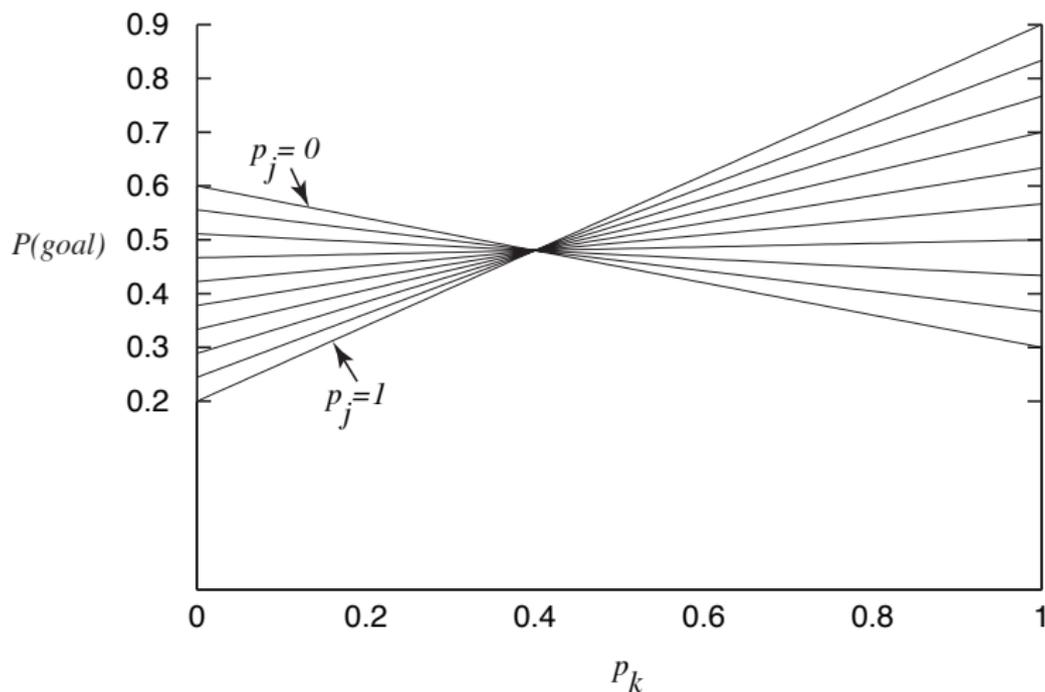
Partial Observability and Competition



		goalie	
		left	right
kicker	left	0.6	0.2
	right	0.3	0.9

Probability of a goal.

Stochastic Policies



Strategy Profiles

- Assume a general n -player game,
- A **strategy** for an agent is a probability distribution over the actions for this agent.
- A **strategy profile** is an assignment of a strategy to each agent.
- A strategy profile σ has a utility for each agent. Let $utility(\sigma, i)$ be the utility of strategy profile σ for agent i .
- If σ is a strategy profile:
 σ_i is the strategy of agent i in σ ,
 σ_{-i} is the set of strategies of the other agents.
Thus σ is $\sigma_i\sigma_{-i}$

Nash Equilibria

- σ_i is a **best response** to σ_{-i} if for all other strategies σ'_i for agent i ,

$$utility(\sigma_i \sigma_{-i}, i) \geq utility(\sigma'_i \sigma_{-i}, i).$$

- A strategy profile σ is a **Nash equilibrium** if for each agent i , strategy σ_i is a best response to σ_{-i} . That is, a Nash equilibrium is a strategy profile such that no agent can be better by unilaterally deviating from that profile.
- Theorem [Nash, 1950] Every finite game has at least one Nash equilibrium.

Multiple Equilibria

Hawk-Dove Game:
Agent 2

		Agent 2	
		dove	hawk
Agent 1	dove	$R/2, R/2$	$0, R$
	hawk	$R, 0$	$-D, -D$

D and R are both positive with $D \gg R$.

Just because you know the Nash equilibria doesn't mean you know what to do:

		Agent 2	
		shopping	football
Agent 1	shopping	2,1	0,0
	football	0,0	1,2

Prisoner's Dilemma

Two strangers are in a game show. They each have the choice:

- Take \$100 for yourself
- Give \$1000 to the other player

This can be depicted as the payoff matrix:

		Player 2	
		take	give
Player 1	take	100,100	1100,0
	give	0,1100	1000,1000

Tragedy of the Commons

Example:

- There are 100 agents.
- There is an common environment that is shared amongst all agents. Each agent has $1/100$ of the shared environment.
- Each agent can choose to do an action that has a payoff of +10 but has a -100 payoff on the environment or do nothing with a zero payoff

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- If every agent does the action the total payoff is

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- If every agent does the action the total payoff is $1000 - 10000 = -9000$

Computing Nash Equilibria

To compute a Nash equilibria for a game in strategic form:

- Eliminate dominated strategies
- Determine which actions will have non-zero probabilities. This is the **support set**.
- Determine the probability for the actions in the support set

Eliminating Dominated Strategies

		Agent 2		
		d_2	e_2	f_2
Agent 1	a_1	3,5	5,1	1,2
	b_1	1,1	2,9	6,4
	c_1	2,6	4,7	0,8

Computing probabilities in randomized strategies

Given a support set:

- Why would an agent will randomize between actions $a_1 \dots a_k$?

Computing probabilities in randomized strategies

Given a support set:

- Why would an agent will randomize between actions $a_1 \dots a_k$? Actions $a_1 \dots a_k$ have the same value for that agent given the strategies for the other agents.
- This forms a set of simultaneous equations where variables are probabilities of the actions
- If there is a solution with all the probabilities in range $(0,1)$ this is a Nash equilibrium.

Search over support sets to find a Nash equilibrium

Learning to Coordinate

- Each agent maintains $P[A]$ a probability distribution over actions.
- Each agent maintains $Q[A]$ an estimate of value of doing A given policy of other agents.
- Repeat:
 - ▶ select action a using distribution P ,
 - ▶ do a and observe payoff
 - ▶ update Q :

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- Repeat:
 - ▶ select action a using distribution P ,
 - ▶ do a and observe payoff
 - ▶ update Q : $Q[a] \leftarrow Q[a] + \alpha(\text{payoff} - Q[a])$
 - ▶ incremented probability of best action by δ .
 - ▶ decremented probability of other actions