Minmax and Maxmin

ISCI 330 Lecture 8

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Minmax and Maxmin

ISCI 330 Lecture 8, Slide 1

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Lecture Overview





ISCI 330 Lecture 8, Slide 2

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Minmax and Maxmin

Computing Mixed Nash Equilibria

- Guess the support
- If a player has a support of size 2 or more, he must be indifferent between these actions
- Set up an equation that expresses these constraints:

• e.g.,
$$u_1(B, (p, 1-p)) = u_1(F, (p, 1-p))$$

• Solve the equation to find *p*.

Lecture Overview







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Minmax and Maxmin

Max-Min Strategies

- Player *i*'s maxmin strategy is a strategy that maximizes *i*'s worst-case payoff, in the situation where all the other players (whom we denote -i) happen to play the strategies which cause the greatest harm to *i*.
- The maxmin value (or safety level) of the game for player *i* is that minimum amount of payoff guaranteed by a maxmin strategy.
- Why would *i* want to play a maxmin strategy?

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- The maxmin value (or safety level) of the game for player *i* is that minimum amount of payoff guaranteed by a maxmin strategy.
- Why would *i* want to play a maxmin strategy?
 - a conservative agent maximizing worst-case payoff
 - a paranoid agent who believes everyone is out to get him

Definition

The maxmin strategy for player *i* is $\arg \max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$, and the maxmin value for player *i* is $\max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$.

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Min-Max Strategies

- Player *i*'s minmax strategy in a 2-player game is a strategy that minimizes the other player -i's best-case payoff.
- The minmax value of the 2-player game for player *i* is that maximum amount of payoff that -i could achieve under *i*'s minmax strategy.
- Why would *i* want to play a minmax strategy?

Min-Max Strategies

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- The minmax value of the 2-player game for player *i* is that maximum amount of payoff that -i could achieve under *i*'s minmax strategy.
- Why would *i* want to play a minmax strategy?
 - to punish the other agent as much as possible

Definition

The maxmin strategy for player i is $\arg \max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$, and the maxmin value for player i is $\max_{s_i} \min_{s_{-i}} u_i(s_1, s_2)$.

Definition

In a two-player game, the minmax strategy for player i is $\arg \min_{s_i} \max_{s_{-i}} u_{-i}(s_1, s_2)$, and the minmax value for player i is $\min_{s_i} \max_{s_{-i}} u_{-i}(s_1, s_2)$.

Minmax and Maxmin

Minmax Theorem

Theorem (Minmax theorem (von Neumann, 1928))

In any finite, two-player, zero-sum game, in any Nash equilibrium each player receives a payoff that is equal to both his maxmin value and his minmax value.

- The maxmin value for one player is equal to the minmax value for the other player. By convention, the maxmin value for player 1 is called the value of the game.
- For both players, the set of maxmin strategies coincides with the set of minmax strategies.
- Any maxmin strategy profile (or, equivalently, minmax strategy profile) is a Nash equilibrium. Furthermore, these are all the Nash equilibria. Consequently, all Nash equilibria have the same payoff vector (namely, those in which player 1 gets the value of the game).

Geometric Representation: Saddle Point



• Can you see why this picture illustrates the maxmin and minmax values?

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How to find maxmin and minmax strategies

Consider maxmin strategies for player i in a 2-player game.

- Notice that i's maxmin strategy depends only on i's utilities
 - thus changes to -i's utilities do not change i's maxmin strategy
- Consider the game where player i has the same utilities as before, but player -i's utilities are replaced with the negatives of i's utilities
 - this is now a zero-sum game
- Because of the minmax theorem, we know that any Nash equilibrium strategy in this game is also a maxmin strategy
 - Thus, find player *i*'s equilibrium strategy in the new game and we have *i*'s maxmin strategy in the original game
- We can use a similar approach for minmax.

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