

Uninformed Search

CPSC 322 – Search 3

Textbook §3.5

Lecture Overview

- 1 Recap
- 2 Depth-First Search
- 3 Breadth-First Search

Graph Search Algorithm

Input: a graph,
a set of start nodes,
Boolean procedure $goal(n)$ that tests if n is a goal node.
 $frontier := \{ \langle s \rangle : s \text{ is a start node} \};$
while $frontier$ is not empty:
 select and remove path $\langle n_0, \dots, n_k \rangle$ from $frontier$;
 if $goal(n_k)$
 return $\langle n_0, \dots, n_k \rangle$;
 for every neighbor n of n_k
 add $\langle n_0, \dots, n_k, n \rangle$ to $frontier$;
end while

- After the algorithm returns, it can be asked for more answers and the procedure continues.
- Which value is selected from the frontier defines the search strategy.
- The *neighbor* relationship defines the graph.
- The *goal* function defines what is a solution.

Comparing Algorithms

Definition (complete)

A search algorithm is **complete** if, whenever at least one solution exists, the algorithm is guaranteed to find a solution within a finite amount of time

Definition (time complexity)

The **time complexity** of a search algorithm is an expression for the worst-case amount of time it will take to run, expressed in terms of the maximum path length m and the maximum branching factor b .

Definition (space complexity)

The **space complexity** of a search algorithm is an expression for the worst-case amount of memory that the algorithm will use, expressed in terms of m and b .

Lecture Overview

- 1 Recap
- 2 Depth-First Search
- 3 Breadth-First Search

Depth-first Search

- **Depth-first search** treats the frontier as a stack
- It always selects one of the last elements added to the frontier.
- **Example:**
 - the frontier is $[p_1, p_2, \dots, p_r]$
 - neighbours of p_1 are $\{n_1, \dots, n_k\}$
- What happens?
 - p_1 is selected, and tested for being a goal.
 - Neighbours of p_1 replace p_1 at the beginning of the frontier.
 - Thus, the frontier is now $[(p_1, n_1), \dots, (p_1, n_k), p_2, \dots, p_r]$.
 - p_2 is only selected when all paths extending p_1 have been explored.

DFS Example

- `http://aispace.org/search/`
- “simple tree graph”

Analysis of Depth-first Search

- Is DFS **complete**?

Analysis of Depth-first Search

- Is DFS **complete**?
 - Depth-first search isn't guaranteed to halt on infinite graphs or on graphs with cycles.
 - However, DFS *is* complete for finite trees.

Analysis of Depth-first Search

- Is DFS **complete**?
 - Depth-first search isn't guaranteed to halt on infinite graphs or on graphs with cycles.
 - However, DFS *is* complete for finite trees.
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?

Analysis of Depth-first Search

- Is DFS **complete**?
 - Depth-first search isn't guaranteed to halt on infinite graphs or on graphs with cycles.
 - However, DFS *is* complete for finite trees.
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?
 - The time complexity is $O(b^m)$: must examine every node in the tree.
 - Search is unconstrained by the goal until it happens to stumble on the goal.

Analysis of Depth-first Search

- Is DFS **complete**?
 - Depth-first search isn't guaranteed to halt on infinite graphs or on graphs with cycles.
 - However, DFS *is* complete for finite trees.
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?
 - The time complexity is $O(b^m)$: must examine every node in the tree.
 - Search is unconstrained by the goal until it happens to stumble on the goal.
- What is the **space complexity**?

Analysis of Depth-first Search

- Is DFS **complete**?
 - Depth-first search isn't guaranteed to halt on infinite graphs or on graphs with cycles.
 - However, DFS *is* complete for finite trees.
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?
 - The time complexity is $O(b^m)$: must examine every node in the tree.
 - Search is unconstrained by the goal until it happens to stumble on the goal.
- What is the **space complexity**?
 - Space complexity is $O(bm)$: the longest possible path is m , and for every node in that path must maintain a fringe of size b .

Using Depth-First Search

- When is DFS **appropriate**?

Using Depth-First Search

- When is DFS **appropriate**?
 - space is restricted
 - solutions tend to occur at the same depth in the tree
 - you know how to order nodes in the list of neighbours so that solutions will be found relatively quickly

Using Depth-First Search

- When is DFS **appropriate**?
 - space is restricted
 - solutions tend to occur at the same depth in the tree
 - you know how to order nodes in the list of neighbours so that solutions will be found relatively quickly

- When is DFS **inappropriate**?

Using Depth-First Search

- When is DFS **appropriate**?
 - space is restricted
 - solutions tend to occur at the same depth in the tree
 - you know how to order nodes in the list of neighbours so that solutions will be found relatively quickly

- When is DFS **inappropriate**?
 - some paths have infinite length
 - the graph contains cycles
 - some solutions are very deep, while others are very shallow

Lecture Overview

- 1 Recap
- 2 Depth-First Search
- 3 Breadth-First Search**

Breadth-first Search

- Breadth-first search treats the frontier as a **queue**
 - it always selects one of the earliest elements added to the frontier.
- **Example:**
 - the frontier is $[p_1, p_2, \dots, p_r]$
 - neighbours of p_1 are $\{n_1, \dots, n_k\}$
- What happens?
 - p_1 is selected, and tested for being a goal.
 - Neighbours of p_1 follow p_r at the end of the frontier.
 - Thus, the frontier is now $[p_2, \dots, p_r, (p_1, n_1), \dots, (p_1, n_k)]$.
 - p_2 is selected next.

BFS Example

- `http://aispace.org/search/`
- “simple tree graph”

Analysis of Breadth-First Search

- Is BFS **complete**?

Analysis of Breadth-First Search

- Is BFS **complete**?
 - Yes (but it wouldn't be if the branching factor for any node was infinite)
 - In fact, BFS is guaranteed to find the path that involves the fewest arcs (why?)

Analysis of Breadth-First Search

- Is BFS **complete**?
 - Yes (but it wouldn't be if the branching factor for any node was infinite)
 - In fact, BFS is guaranteed to find the path that involves the fewest arcs (why?)
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?

Analysis of Breadth-First Search

- Is BFS **complete**?
 - Yes (but it wouldn't be if the branching factor for any node was infinite)
 - In fact, BFS is guaranteed to find the path that involves the fewest arcs (why?)
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?
 - The time complexity is $O(b^m)$: must examine every node in the tree.
 - The order in which we examine nodes (BFS or DFS) makes no difference to the worst case: search is unconstrained by the goal.

Analysis of Breadth-First Search

- Is BFS **complete**?
 - Yes (but it wouldn't be if the branching factor for any node was infinite)
 - In fact, BFS is guaranteed to find the path that involves the fewest arcs (why?)
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?
 - The time complexity is $O(b^m)$: must examine every node in the tree.
 - The order in which we examine nodes (BFS or DFS) makes no difference to the worst case: search is unconstrained by the goal.
- What is the **space complexity**?

Analysis of Breadth-First Search

- Is BFS **complete**?
 - Yes (but it wouldn't be if the branching factor for any node was infinite)
 - In fact, BFS is guaranteed to find the path that involves the fewest arcs (why?)
- What is the **time complexity**, if the maximum path length is m and the maximum branching factor is b ?
 - The time complexity is $O(b^m)$: must examine every node in the tree.
 - The order in which we examine nodes (BFS or DFS) makes no difference to the worst case: search is unconstrained by the goal.
- What is the **space complexity**?
 - Space complexity is $O(b^m)$: we must store the whole frontier in memory

Using Breadth-First Search

- When is BFS **appropriate**?

Using Breadth-First Search

- When is BFS **appropriate**?
 - space is not a problem
 - it's necessary to find the solution with the fewest arcs
 - although all solutions may not be shallow, at least some are
 - there may be infinite paths

Using Breadth-First Search

- When is BFS **appropriate**?
 - space is not a problem
 - it's necessary to find the solution with the fewest arcs
 - although all solutions may not be shallow, at least some are
 - there may be infinite paths

- When is BFS **inappropriate**?

Using Breadth-First Search

- When is BFS **appropriate**?
 - space is not a problem
 - it's necessary to find the solution with the fewest arcs
 - although all solutions may not be shallow, at least some are
 - there may be infinite paths

- When is BFS **inappropriate**?
 - space is limited
 - all solutions tend to be located deep in the tree
 - the branching factor is very large