

# Propositional Logic: Semantics and Bottom-Up Proofs

CPSC 322 Lecture 19

February 28, 2007

Textbook §4.2

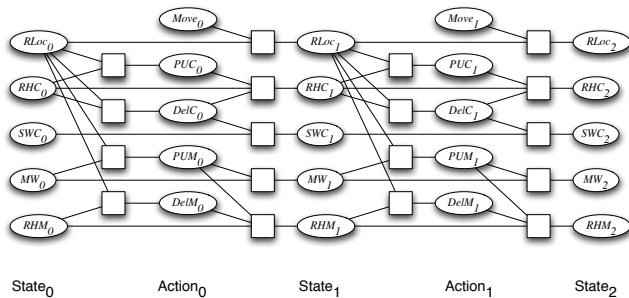
# Lecture Overview

- 1 Recap
- 2 Semantics
- 3 Using Logic to Model the World

# Planning as a CSP

- We can go forwards and backwards at the same time, if we set up a planning problem as a CSP
- To do this, we need to “unroll” the plan for a fixed number of steps
  - this is called the **horizon**
- To do this with a horizon of  $k$ :
  - construct a **variable for each feature at each time step** from 0 to  $k$
  - construct a boolean **variable for each action at each time step** from 0 to  $k - 1$ .

# CSP Planning: Robot Example



Do you see why CSP planning is both forwards and backwards?

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## Definition (knowledge base)

A **knowledge base** is a set of definite clauses



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Do any of these statements *mean* anything? Syntax doesn't answer this question.

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# Propositional Definite Clauses: Semantics

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An **interpretation**  $I$  assigns a truth value to each atom.

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We can use the interpretation to determine the truth value of clauses and knowledge bases:

## Definition (truth values of statements)

- A **body**  $b_1 \wedge b_2$  is true in  $I$  if and only if  $b_1$  is true in  $I$  and  $b_2$  is true in  $I$ .
- A **rule**  $h \leftarrow b$  is false in  $I$  if and only if  $b$  is true in  $I$  and  $h$  is false in  $I$ .
- A **knowledge base**  $KB$  is true in  $I$  if and only if every clause in  $KB$  is true in  $I$ .

# Models and Logical Consequence

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## Definition (logical consequence)

If  $KB$  is a set of clauses and  $g$  is a conjunction of atoms,  $g$  is a **logical consequence** of  $KB$ , written  $KB \models g$ , if  $g$  is *true* in every model of  $KB$ .

- we also say that  $g$  **logically follows** from  $KB$ , or that  $KB$  **entails**  $g$ .
- In other words,  $KB \models g$  if there is no interpretation in which  $KB$  is *true* and  $g$  is *false*.

# Example: Models

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	<i>p</i>	<i>q</i>	<i>r</i>	<i>s</i>
<i>I</i> <sub>1</sub>	<i>true</i>	<i>true</i>	<i>true</i>	<i>true</i>

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Which of the following is true?

- $KB \models q$ ,  $KB \models p$ ,  $KB \models s$ ,  $KB \models r$

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# User's view of Semantics

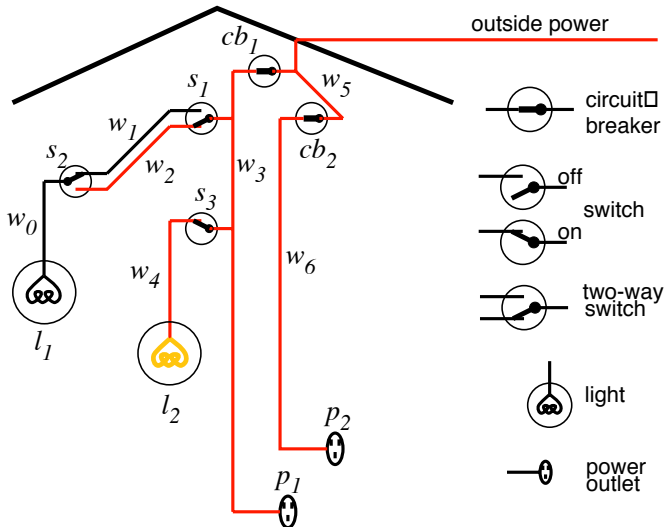
- 1 Choose a task domain: **intended interpretation**.
- 2 Associate an atom with each proposition you want to represent.
- 3 Tell the system clauses that are true in the intended interpretation: **axiomatizing the domain**.
- 4 Ask questions about the intended interpretation.
- 5 If  $KB \models g$ , then  $g$  must be true in the intended interpretation.
- 6 The user can interpret the answer using their intended interpretation of the symbols.



# Computer's view of semantics

- The computer doesn't have access to the intended interpretation.
  - All it knows is the knowledge base.
- The computer can determine if a formula is a logical consequence of KB.
  - If  $KB \models g$  then  $g$  must be true in the intended interpretation.
  - If  $KB \not\models g$  then there is a model of  $KB$  in which  $g$  is false. This could be the intended interpretation.

# Electrical Environment



# Representing the Electrical Environment

$light\_l_1.$

$light\_l_2.$

$down\_s_1.$

$up\_s_2.$

$up\_s_3.$

$ok\_l_1.$

$ok\_l_2.$

$ok\_cb_1.$

$ok\_cb_2.$

$live\_outside.$

$live\_l_1 \leftarrow live\_w_0$

$live\_w_0 \leftarrow live\_w_1 \wedge up\_s_2.$

$live\_w_0 \leftarrow live\_w_2 \wedge down\_s_2.$

$live\_w_1, \leftarrow live\_w_3 \wedge up\_s_1.$

$live\_w_2 \leftarrow live\_w_3 \wedge down\_s_1.$

$live\_l_2 \leftarrow live\_w_4.$

$live\_w_4 \leftarrow live\_w_3 \wedge up\_s_3.$

$live\_p_1 \leftarrow live\_w_3.$

$live\_w_3 \leftarrow live\_w_5 \wedge ok\_cb_1.$

$live\_p_2 \leftarrow live\_w_6.$

$live\_w_6 \leftarrow live\_w_5 \wedge ok\_cb_2.$

$live\_w_5 \leftarrow live\_outside.$

# Role of semantics

## In user's mind:

- $l2\_broken$ : light  $l2$  is broken
- $sw3\_up$ : switch is up
- $power$ : there is power in the building
- $unlit\_l2$ : light  $l2$  isn't lit
- $lit\_l1$ : light  $l1$  is lit

## In Computer:

$$l2\_broken \leftarrow sw3\_up$$
$$\wedge power \wedge unlit\_l2.$$
$$sw3\_up.$$
$$power \leftarrow lit\_l1.$$
$$unlit\_l2.$$
$$lit\_l1.$$

---

## Conclusion: $l2\_broken$

- The computer doesn't know the meaning of the symbols
- The user can interpret the symbols using their meaning