

Game Theory: Analyzing Games

CPSC 322 Lecture 35

April 5, 2006

Reading: excerpt from "Multiagent Systems", chapter 3.

Lecture Overview

Recap

Pareto Optimality

Best Response and Nash Equilibrium

Mixed Strategies

Non-Cooperative Game Theory

- ▶ What is it?
 - ▶ mathematical study of interaction between **rational**, **self-interested** agents
- ▶ Why is it called non-cooperative?
 - ▶ while it's most interested in situations where agents' interests conflict, it's not restricted to these settings
 - ▶ the key is that the individual is the basic modeling unit, and that individuals pursue their own interests
 - ▶ cooperative/coalitional game theory has teams as the central unit, rather than agents
- ▶ You can think of a non-cooperative game as a decision diagram where different agents control different decision nodes, and where each agent has his own utility node.

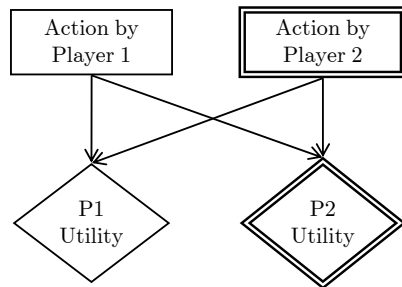
Defining Games

- ▶ Finite, n -person game: $\langle N, A, u \rangle$:
 - ▶ N is a finite set of n **players**, indexed by i
 - ▶ $A = A_1, \dots, A_n$ is a set of **actions** for each player i
 - ▶ $a \in A$ is an **action profile**
 - ▶ $u = \{u_1, \dots, u_n\}$, a **utility function** for each player, where $u_i : A \mapsto \mathbb{R}$
- ▶ Writing a 2-player game as a **matrix**:
 - ▶ row player is player 1, column player is player 2
 - ▶ rows are actions $a \in A_1$, columns are $a' \in A_2$
 - ▶ cells are outcomes, written as a tuple of utility values for each player

Games in Matrix Form

Here's the **TCP Backoff Game** written as a matrix ("normal form") and as a decision network.

	<i>C</i>	<i>D</i>
<i>C</i>	-1, -1	-4, 0
<i>D</i>	0, -4	-3, -3



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Analyzing Games

- ▶ We've defined some canonical games, and thought about how to play them. Now let's examine the games from the **outside**
- ▶ From the point of view of an outside observer, can some outcomes of a game be said to be **better** than others?

Analyzing Games

- ▶ We've defined some canonical games, and thought about how to play them. Now let's examine the games from the **outside**
- ▶ From the point of view of an outside observer, can some outcomes of a game be said to be **better** than others?
 - ▶ we have no way of saying that one agent's interests are more important than another's
 - ▶ intuition: imagine trying to find the revenue-maximizing outcome when you don't know what currency has been used to express each agent's payoff
- ▶ Are there situations where we can still prefer one outcome to another?

Pareto Optimality

- ▶ **Idea:** sometimes, one outcome o is at least as good for every agent as another outcome o' , and there is no agent who strictly prefers o' to o
 - ▶ in this case, it seems reasonable to say that o is better than o'
 - ▶ we say that o **Pareto-dominates** o' .

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 - ▶ can a game have more than one Pareto-optimal outcome?

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- ▶ An outcome o^* is **Pareto-optimal** if there is no other outcome which Pareto-dominates it.
 - ▶ can a game have more than one Pareto-optimal outcome?
 - ▶ does every game have at least one Pareto-optimal outcome?

Pareto Optimal Outcomes in Example Games

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B	2, 1	0, 0
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	Heads	Tails
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Best Response

- ▶ If you knew what everyone else was going to do, it would be easy to pick your own action
 - ▶ phrased as a decision diagram: observing the other players' decision nodes as evidence
- ▶ Let $a_{-i} = \langle a_1, \dots, a_{i-1}, a_{i+1}, \dots, a_n \rangle$.
 - ▶ now $a = (a_{-i}, a_i)$
- ▶ **Best response:** $a_i^* \in BR(a_{-i})$ iff
$$\forall a_i \in A_i, u_i(a_i^*, a_{-i}) \geq u_i(a_i, a_{-i})$$

Nash Equilibrium

- ▶ Now let's return to the setting where no agent knows anything about what the others will do
- ▶ What can we say about which actions will occur?

Nash Equilibrium

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- ▶ What can we say about which actions will occur?

- ▶ Idea: look for **stable** action profiles.
- ▶ $a = \langle a_1, \dots, a_n \rangle$ is a **Nash equilibrium** iff $\forall i, a_i \in BR(a_{-i})$.

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The paradox of *Prisoner's dilemma*: the Nash equilibrium is the only non-Pareto-optimal outcome!

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Mixed Strategies

- ▶ It would be a pretty bad idea to play any deterministic strategy in matching pennies
- ▶ Idea: confuse the opponent by playing **randomly**
- ▶ Define a **strategy** s_i for agent i as any probability distribution over the actions A_i .
 - ▶ **pure strategy**: only one action is played with positive probability
 - ▶ **mixed strategy**: more than one action is played with positive probability
- ▶ Let the set of **all strategies** for i be S_i
- ▶ Let the set of **all strategy profiles** be $S = S_1 \times \dots \times S_n$.

Utility under Mixed Strategies

- ▶ What is your **payoff** if all the players follow mixed strategy profile $s \in S$?
 - ▶ We can't just read this number from the game matrix anymore: we won't always end up in the same cell

Utility under Mixed Strategies

- ▶ What is your **payoff** if all the players follow mixed strategy profile $s \in S$?
 - ▶ We can't just read this number from the game matrix anymore: we won't always end up in the same cell
- ▶ Instead, use the idea of **expected utility** from decision theory:

$$u_i(s) = \sum_{a \in A} u_i(a) Pr(a|s)$$

$$Pr(a|s) = \prod_{j \in N} s_j(a_j)$$

Best Response and Nash Equilibrium

Our definitions of best response and Nash equilibrium generalize from actions to strategies.

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Our definitions of best response and Nash equilibrium generalize from actions to strategies.

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▶ $s = \langle s_1, \dots, s_n \rangle$ is a Nash equilibrium iff $\forall i, s_i \in BR(s_{-i})$

▶ **Every finite game has a Nash equilibrium!** [Nash, 1950]

▶ e.g., matching pennies: both players play heads/tails 50%/50%