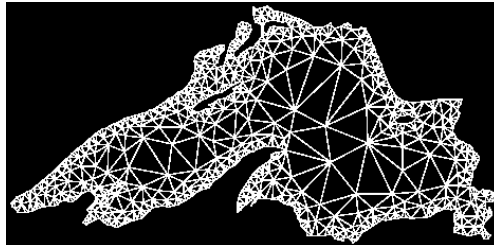
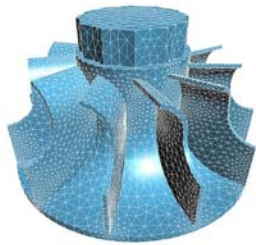
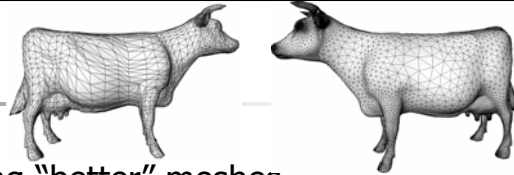


## Mesh Generation

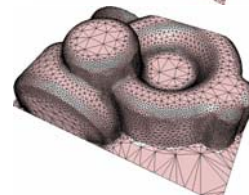
### Introduction & Meshing in 2D



## Motivation



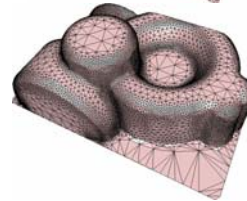
- Remeshing = creating “better” meshes
  - Element sizing
  - Element shape
  - Approximation quality
- Applications
  - Computer Graphics
    - Applications: rendering, editing
    - Compact but accurate description
  - Engineering
    - CAD – model description
    - **Analysis** – finite element meshing





## Mesh Properties

- Depend on application
- Universal
  - Approximation of data
  - Number of elements (less is better)
- Application dependant
  - Element size
    - Graphics – does not matter
    - FEM – defines solution accuracy (depends on equation)
  - Element shape (geometry)
    - Equilateral elements are optimal
    - How crucial depends on application

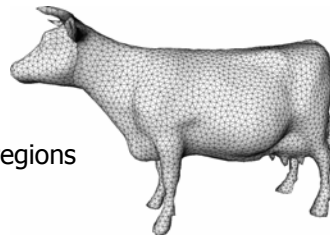


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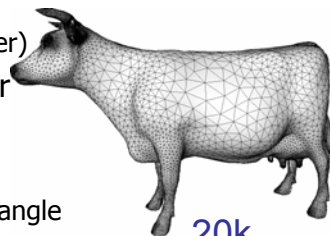


## Mesh Properties

- Approximation of data
  - 2D – accurate approximation of boundary
    - shorter edges in more curved regions
  - 3D – accurate approximation of surface
    - Smaller elements in more curved regions
    - Edges close to surface (more later)
- Sizing function – provided by user
  - In 2D –edge length at any  $(u,v)$
  - 3D –value per vertex
    - barycentric coordinates inside triangle



5k



20k

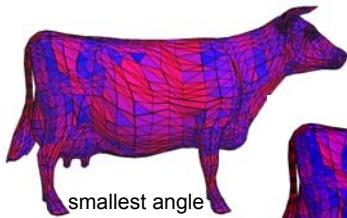


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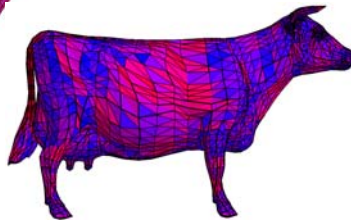


## Quality

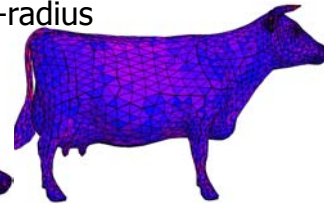
- Measure “closeness” to equilateral triangle
- Triangle quality measures
  - Ratio of in-radius to circum-radius
  - Smallest angle
  - Ratio of shortest edge to circum-radius



smallest angle



shortest edge to circum-radius



smallest angle

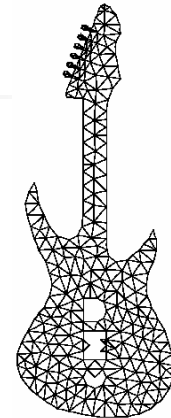


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## Meshing in 2D

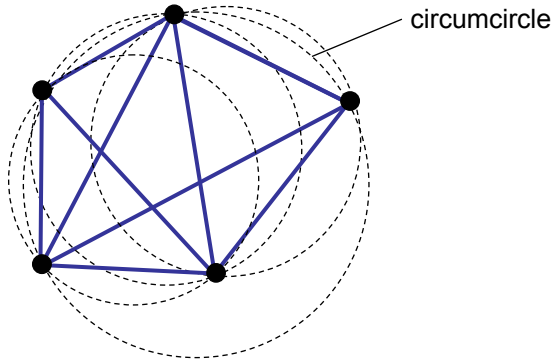
- Input:
  - Planar domain (polygon)
  - Optimal sizing
    - Gradation + budget (e.g. #vertices)
    - Sizing function at each  $(u,v)$
- Output: triangular mesh
- Motivation
  - 2D problems
  - 3D problems reduced to 2D (parameterization)
- Problem components
  - Vertex placement
  - Connectivity construction: Delaunay criterion



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# Constructing Connectivity Delaunay Criterion



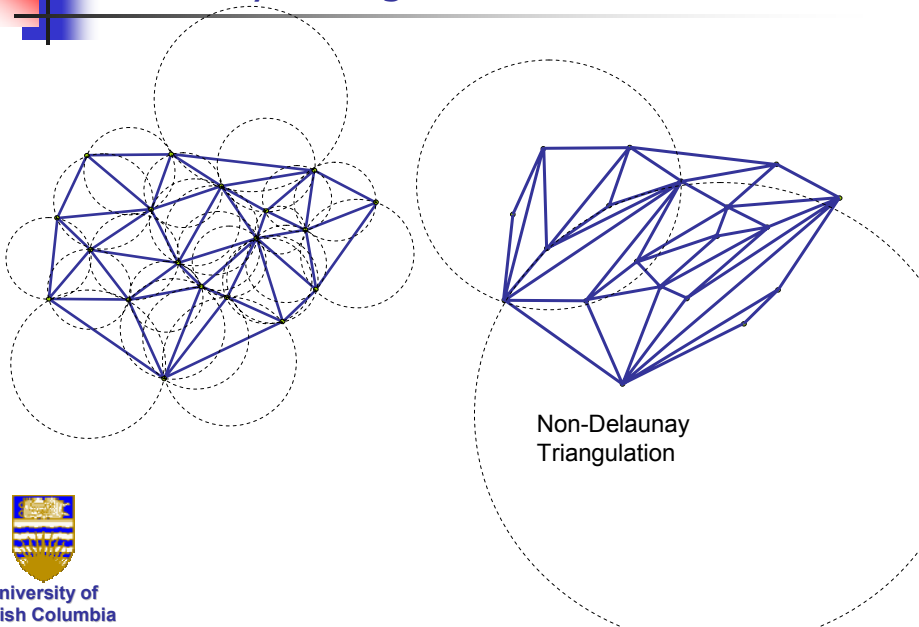
*Empty Circle Property:*  
No other vertex is contained within the  
circumcircle of any triangle



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# Delaunay Triangulation



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## Delaunay Triangulation

- Obeys empty-circle property
- Exists for any set of vertices
- Is **unique** (up to degenerate cases)
- Proven to provide best triangles in terms of quality for given vertex positions
- To test – enough to check pairs of triangles sharing common edge



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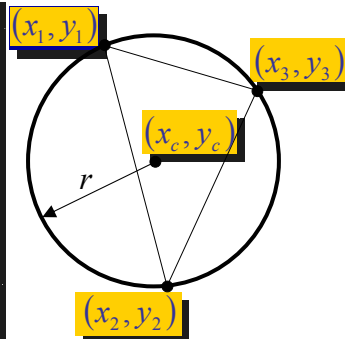
## Delaunay - Practicalities

- Computing circumcircle center & radius

$$Ax + By + C = x^2 + y^2$$

$$\begin{bmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{bmatrix} \begin{Bmatrix} A \\ B \\ C \end{Bmatrix} = \begin{bmatrix} x_1^2 + y_1^2 \\ x_2^2 + y_2^2 \\ x_3^2 + y_3^2 \end{bmatrix}$$

$$x_c = \frac{A}{2}, y_c = \frac{B}{2}, r = \sqrt{C + x_c^2 + y_c^2}$$



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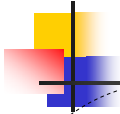


## Triangulation Methods

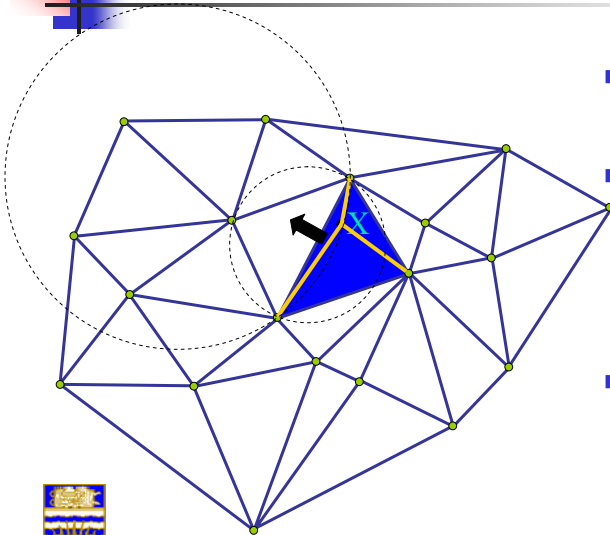
- Edge flip algorithm
  - Start with any triangulation of the vertices
  - Test all edges if satisfy Delaunay criterion
    - test triangles on both sides of edge
  - If edge does not satisfy it, flip edge
  - Repeat until all edges satisfy criterion
- Proven to terminate & give Delaunay mesh
- Slow  $O(n^2)$
- Alternative – additive construction
  - Keep Delaunay mesh
  - Add one vertex at a time



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## Vertex Insertion



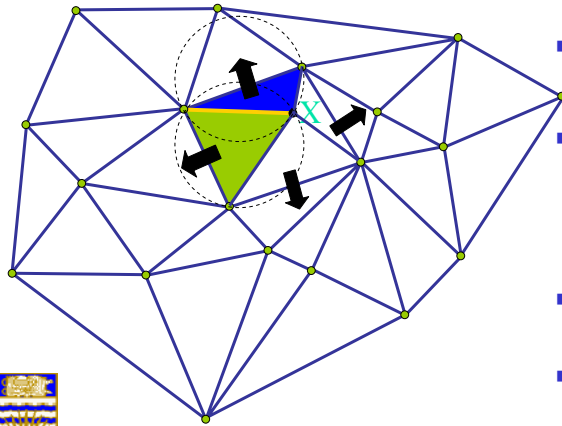
- Locate triangle containing X
- Subdivide triangle
- Recursively check adjoining triangles to ensure empty-circle property



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## Vertex Insertion



- Locate triangle containing X
- Subdivide triangle
- Recursively check adjoining triangles to ensure empty-circle property
- Swap diagonal if needed
- Expected time  $O(n \log(n))$

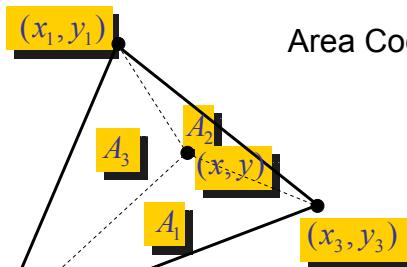


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## Practicalities- Find Triangle

- Search through mesh
- Test inside:



Area Coordinates  $(A_1, A_2, A_3)$

$$A_i = y_{i+2}(x_{i+1} - x) + y_{i+1}(x - x_{i+2}) + y(x_{i+2} - x_{i+1})$$

$(x, y)$  is inside triangle if  $A_i > 0$



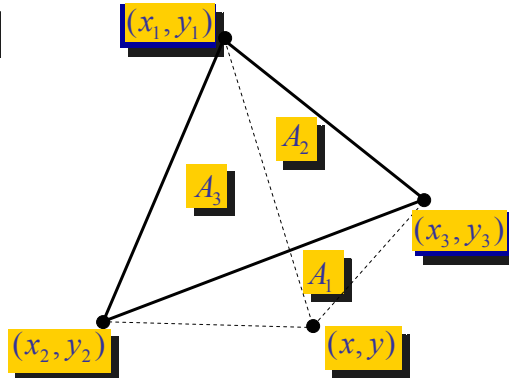
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## Find Triangle

If  $A_i < 0$   
then move to adjacent  
triangle at edge

$(x_{i+1}, y_{i+1}), (x_{i+2}, y_{i+2})$

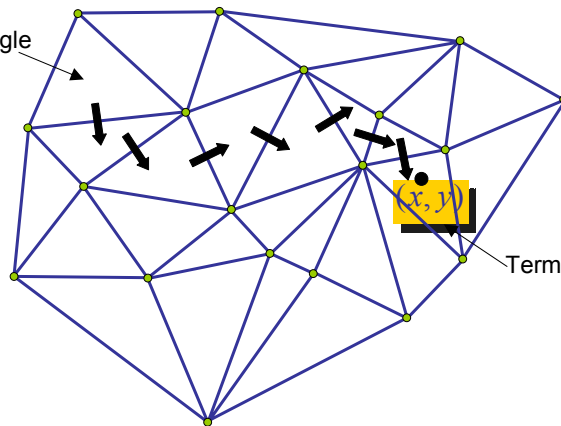


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## Find Triangle

Starting triangle



Terminal triangle

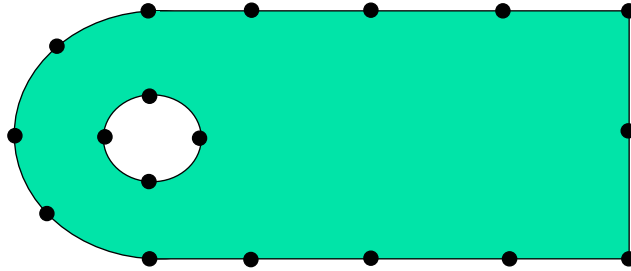


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## Example: Boundary Insertion



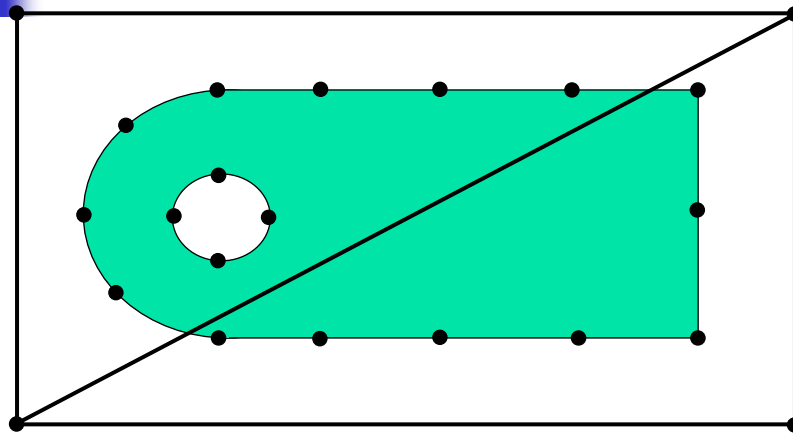
- Place vertices on boundary at cord-length intervals based on sizing



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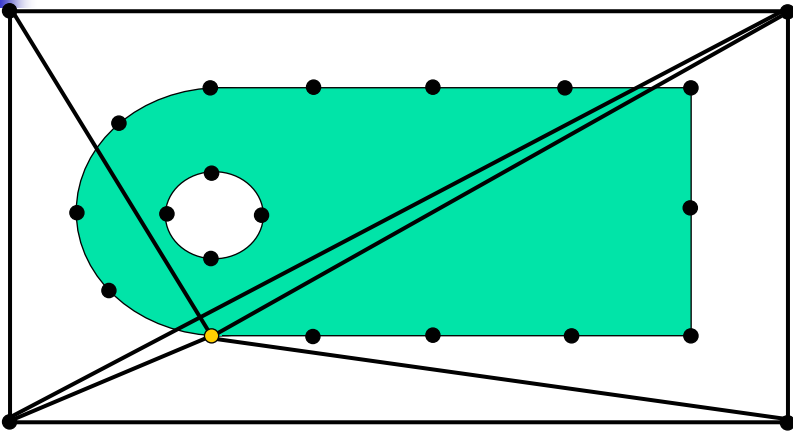
## Boundary Insertion



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- Create bounding triangles

## Boundary Insertion



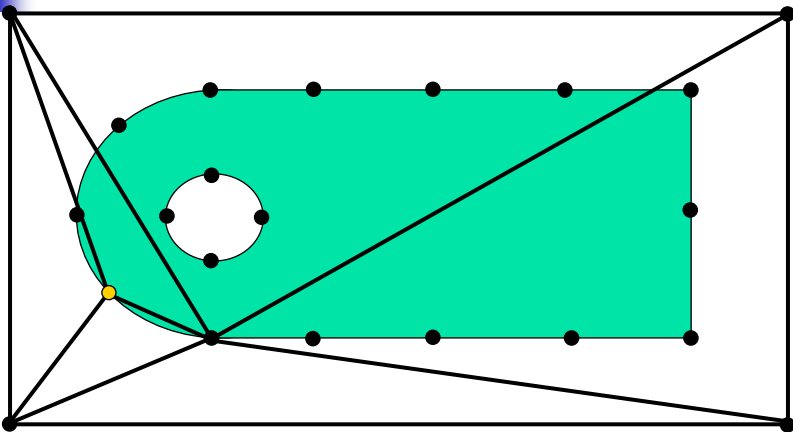
The diagram shows a rectangular domain with a hole on the left side. A yellow dot is positioned on the bottom boundary of the domain. The domain is divided into several triangles by lines connecting the corners and the hole's boundary. The yellow dot is located within one of these triangles.

□ Insert vertices using Delaunay method




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## Boundary Insertion



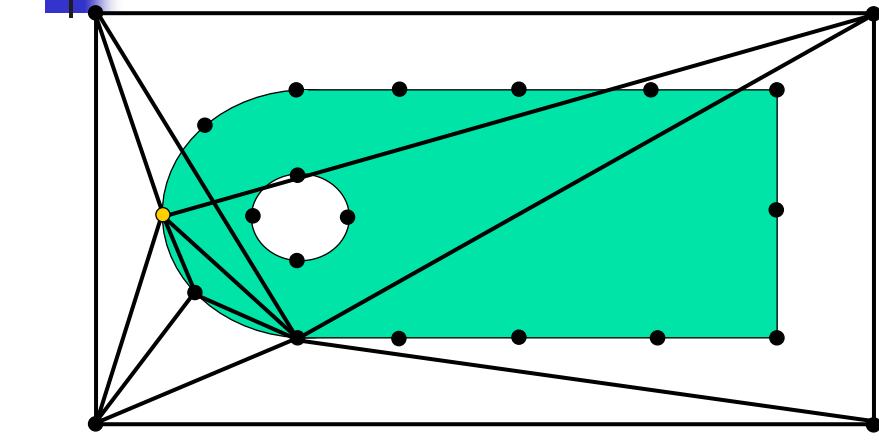
The diagram shows the same rectangular domain with a hole and a yellow dot on the bottom boundary. In this step, the yellow dot has moved to the left boundary of the domain, specifically at the point where the bottom boundary meets the hole's boundary.

□ Insert vertices using Delaunay method




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# Boundary Insertion

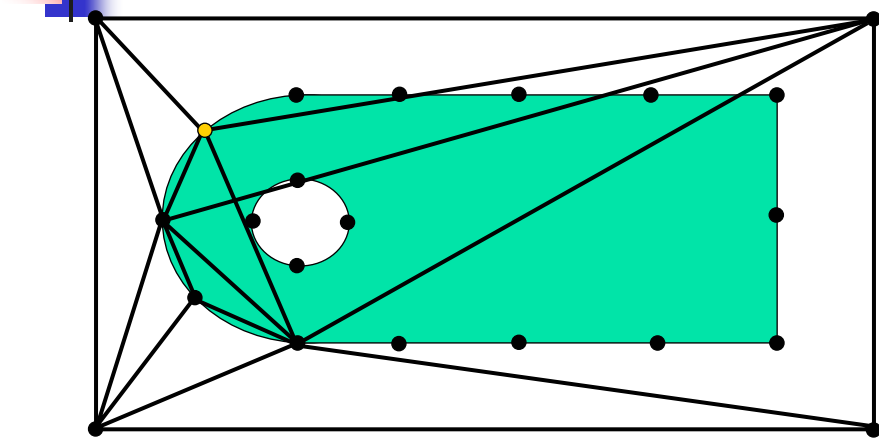


The diagram shows a rectangular domain with a hole on the left side. A set of black dots represents vertices. A yellow dot is positioned on the left boundary of the hole. Black lines connect the vertices to form a mesh. The region to the right of the hole is shaded green. The yellow dot is currently outside the green region.


  Insert vertices using Delaunay method

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# Boundary Insertion



The diagram is identical to the one above, but the yellow dot has moved to the left boundary of the hole, now positioned within the green shaded region.

  Insert vertices using Delaunay method

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# Boundary Insertion

Insert vertices using Delaunay method

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# Boundary Insertion

Delete outside triangles

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## Geometry: Placing Vertices

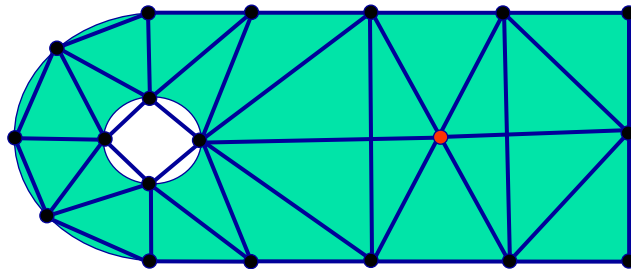
- Beforehand – domain sampling
- After:
  - Refinement
    - Split any edge longer than specified
      - (typically > 1.5 or 2 of required)
  - Coarsening
    - Collapse short edges
  - Smoothing
- Can combine both
- After approach is simpler to implement
- Without “smart” smoothing - poor distribution
- Each vertex added to mesh as described



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## Refinement - Edge Split



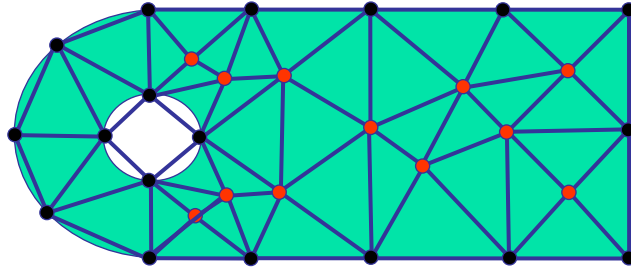
- Split long edges – one at a time (Why?)
  - Typically start from longest
- Add to mesh (one by one)



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## Coarsening - Edge Collapse



- Short edges – result of splitting nearby edges
- Collapse – check if any flips needed



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## Smoothing

- Perform a few iterations of smoothing (e.g. Laplacian) with boundary vertices fixed
  - Other weights better oriented towards good spacing exist
- Note: domain not convex – test for foldovers
- After smoothing perform edge flips on edges not satisfying Delaunay requirements



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## Sampling

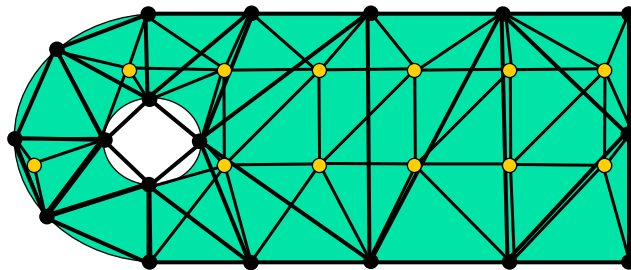
- Sampling provides initial placement for vertices before any triangulation exists
- Similar to quantization:
  - have "budget" (estimated number of vertices)
  - Need to optimally spread thru the domain
- Choices
  - Grid – place vertices on uniform 2D grid
  - Random / Smart random
  - Centroidal Voronoi diagram



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## Uniform Sampling



- Place vertices on regular grid
- Artifacts on boundary
- 90/45/45 triangles
  - Very simple
  - Non-Uniform distribution (why?)

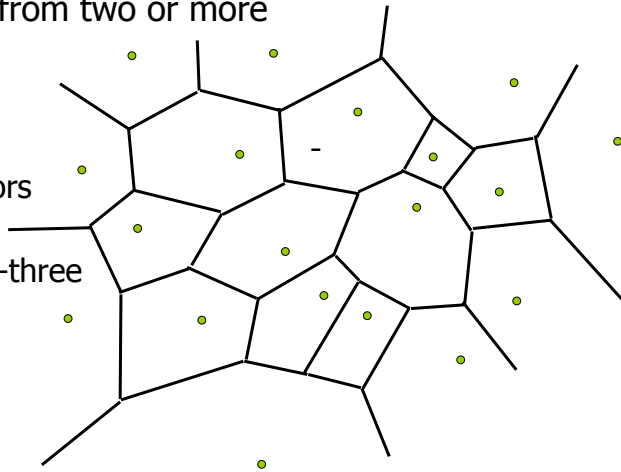


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## Voronoi Diagram

- Voronoi Diagram for given set of vertices: union of all locations at equal distance from two or more vertices
- Consists of
  - straight lines vertex bisectors
  - bisector intersections -three or more lines
  - (why?)

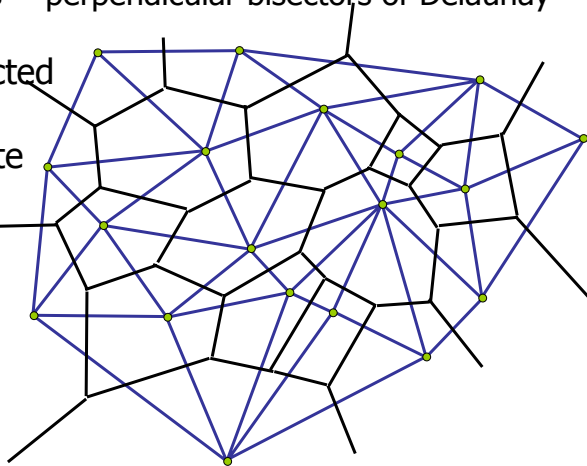


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## Voronoi Diagram

- Dual to Delaunay Triangulation
  - Vertices correspond to faces
  - Voronoi edges = perpendicular bisectors of Delaunay edges
- Can be constructed directly
- Easier – compute Delaunay & compute dual



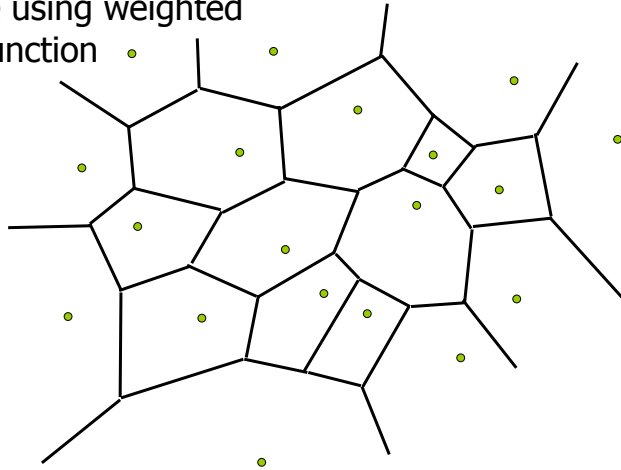
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## Voronoi Diagram

- Diagram partitions space into regions "closer" to one vertex than other
- Can define using weighted distance function

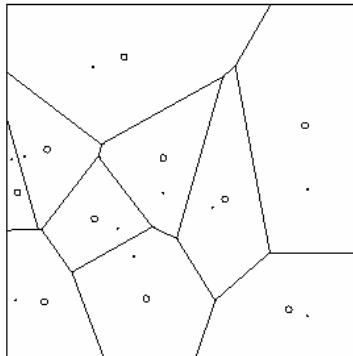


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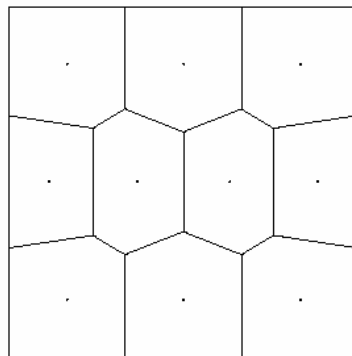


## Centroidal Voronoi diagram

Vertices (sites) **coincide** with centroids (center of mass)



Ordinary Voronoi diagram



Centroidal Voronoi diagram



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## Centroidal VD

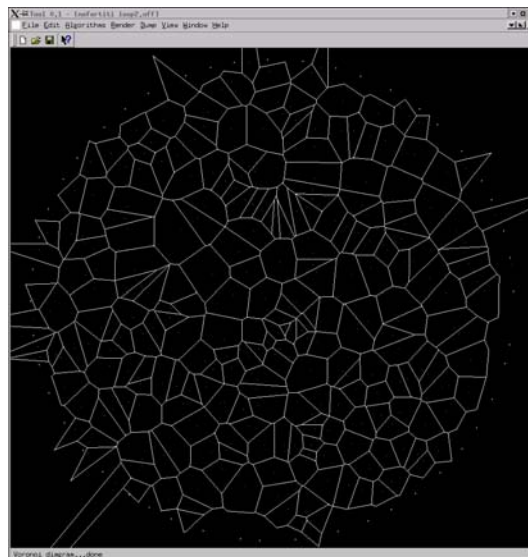
- To compute use Lloyd iterations:
  - Start with set of sites
  - Do
    - Compute VD
    - Compute centers of mass for each Voronoi cell
    - If sites = centers of mass
      - Stop
    - Set sites to centers of mass
    - Repeat
  - Guaranteed to converge
  - Provides optimal repartitioning of density among vertices



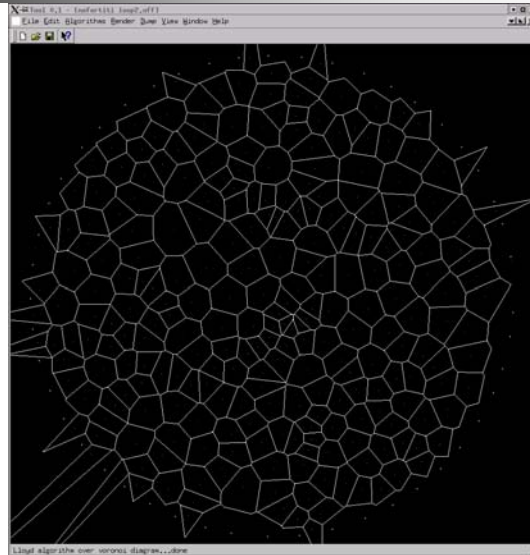
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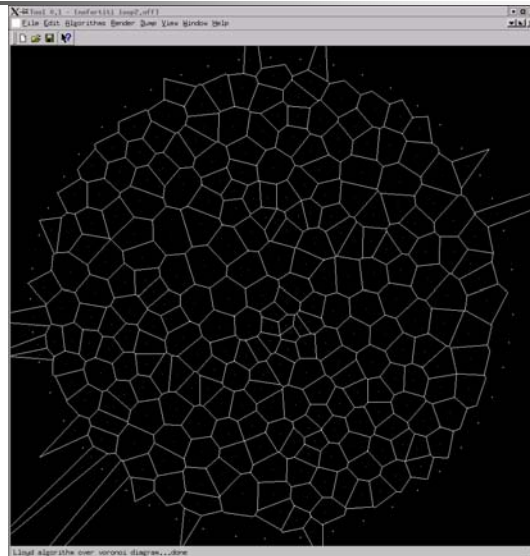
## Lloyd & uniform density



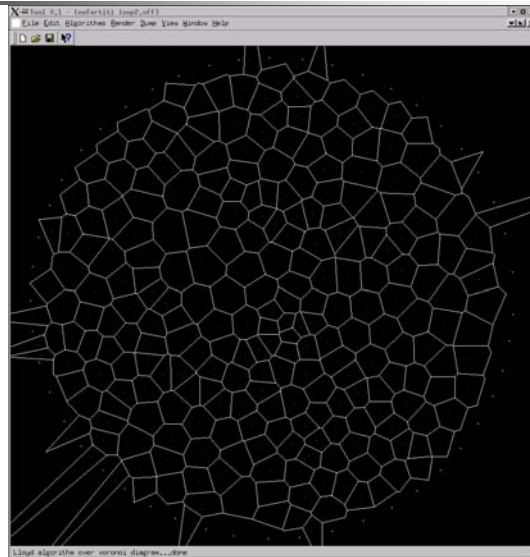
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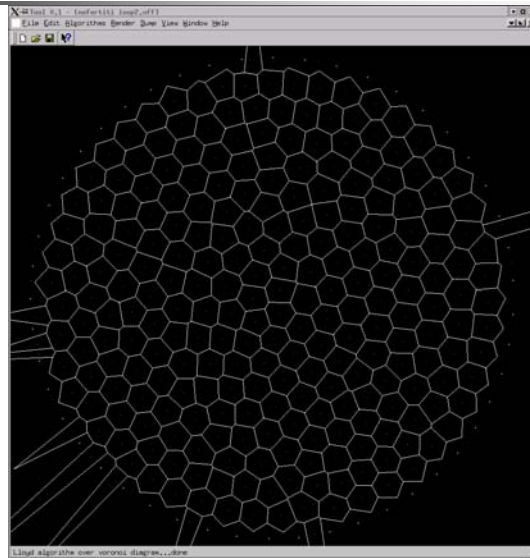
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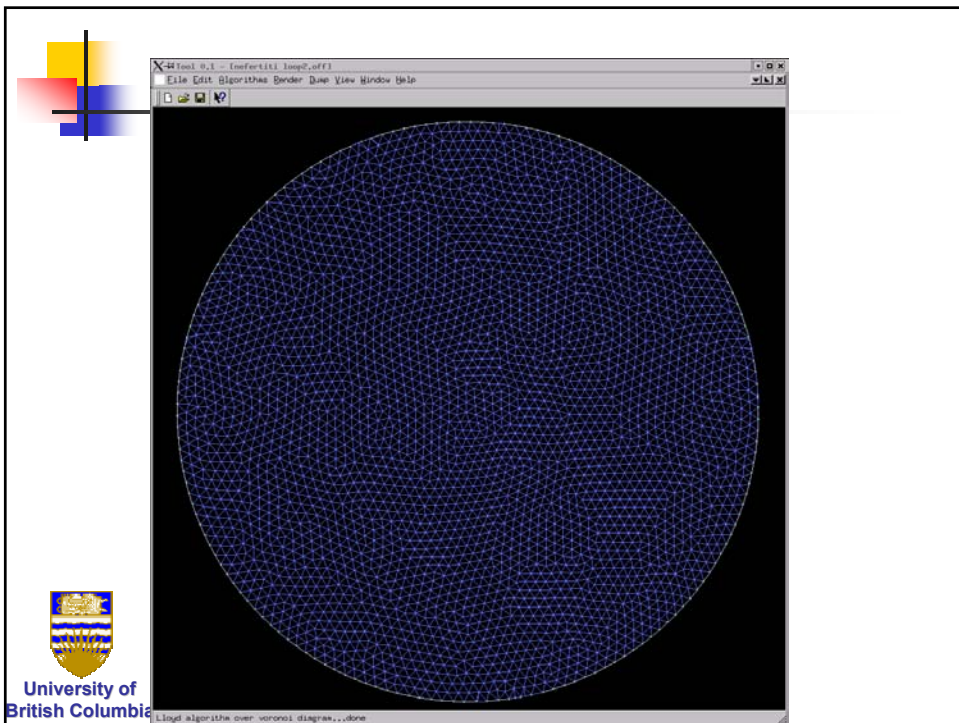
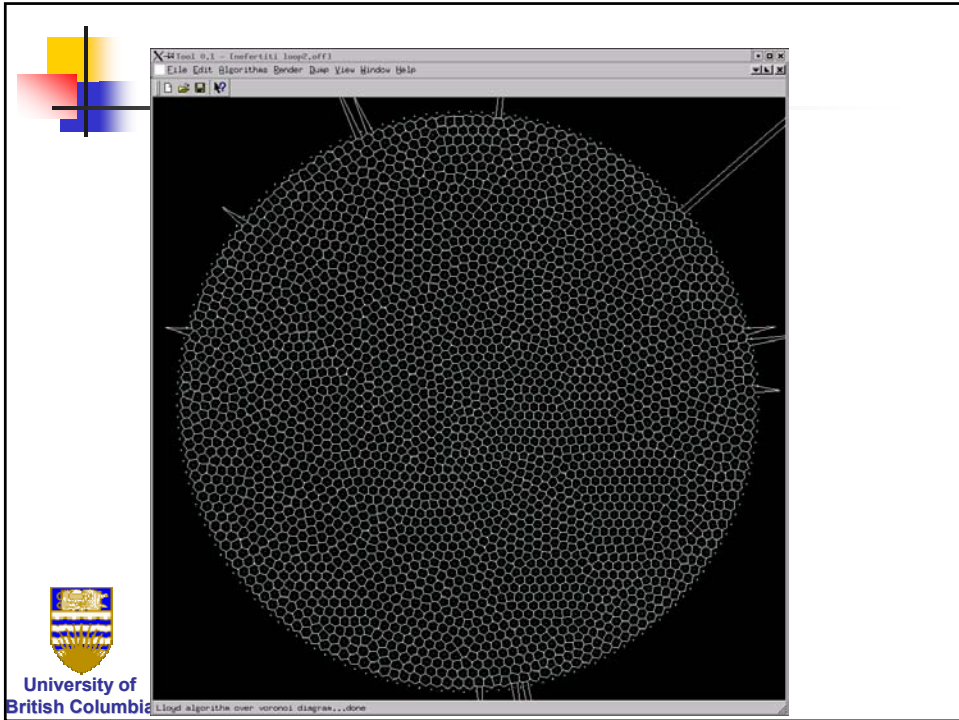
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## Improving Quality

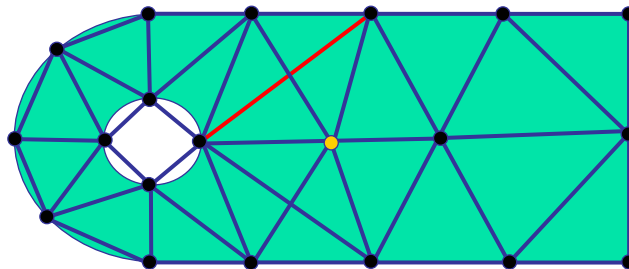
- “Good” sizing does not guarantee good element shape
- Solution: Insert additional vertices
- Insertion criterion - use measures above (min angle, edge ratio,..)
- Insertion strategies
  - Edge split
  - Circumcenters



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## Quality - Edge Split



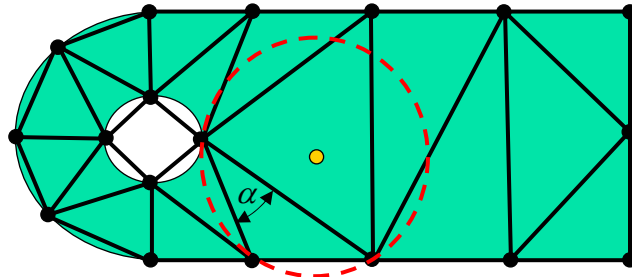
- Split edge if opposite very obtuse angle
- Start from longest
  - Start from worst triangle
  - If offender edge belongs to triangle with even longer edge, consider it (recursively) first
  - Why?



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## Quality: Vertex Insertion – Circumcenter



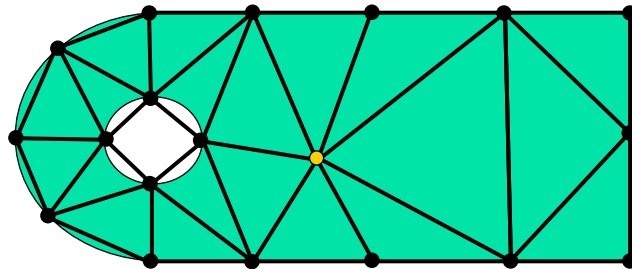
- Vertices introduced at triangle circumcenters
- “Guaranteed Quality”
- Continue till minimal angle > predefined minimum



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## Vertex Insertion - Circumcenter



- Add vertices on boundary if circumcenter is outside (if can)
- Performs better than edge split
- More complex to implement



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## Meshing 2D

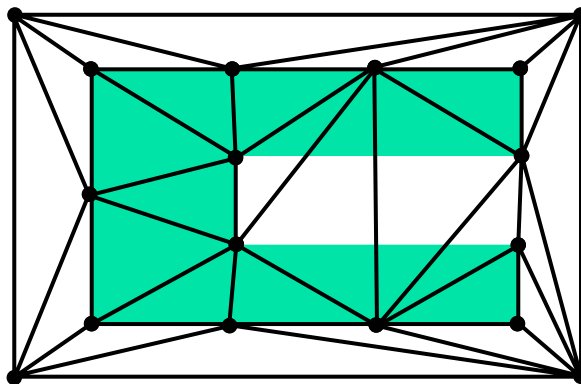
- Place vertices on boundary
- Use sampling for initial placement inside
- Construct Delaunay triangulation
- Iterate
  - Refinement
  - Coarsening
  - Smoothing
  - Each time perform necessary edge flips
- Last component: boundary recovery



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## Boundary Recovery



- Delaunay triangulation does not have to obey polygon boundary

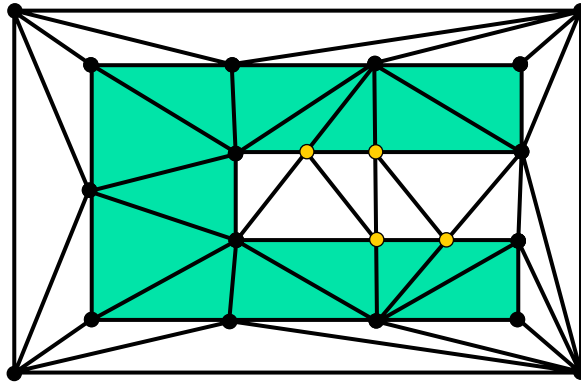


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## Boundary Recovery

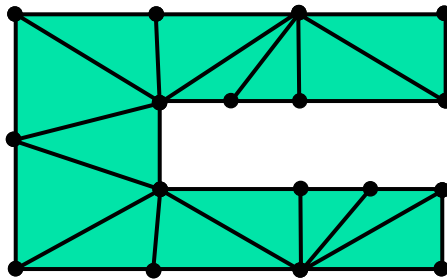


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- Boundary Conforming – add vertices at intersections
- Repeat if necessary



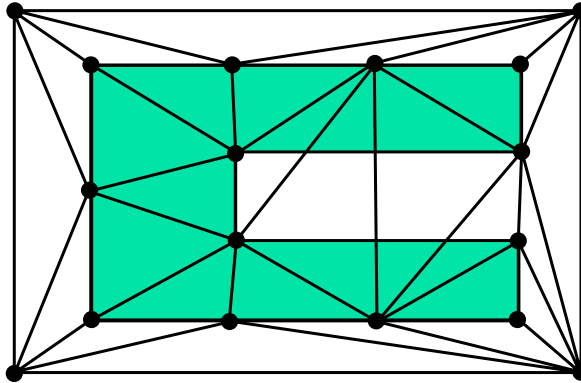
## Boundary Conforming Delaunay



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## Boundary Recovery - Constrained



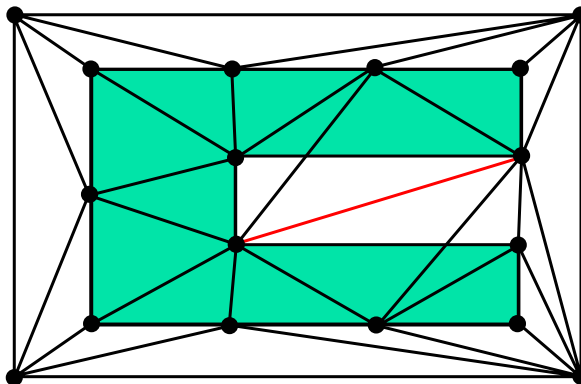
- Not always can add boundary vertices (shared edges)



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## Boundary Recovery - Constrained



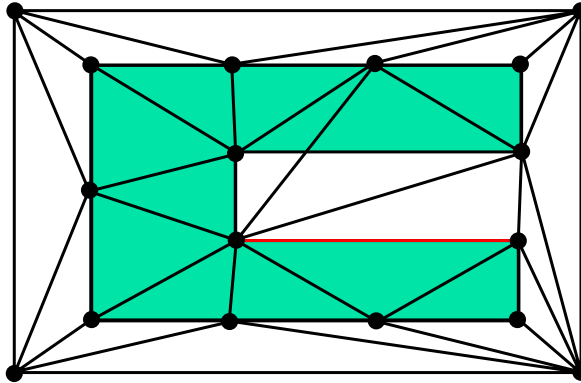
- Swap edges between adjacent pairs of triangles
- Repeat till recover the boundary



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## Boundary Recovery - Constrained



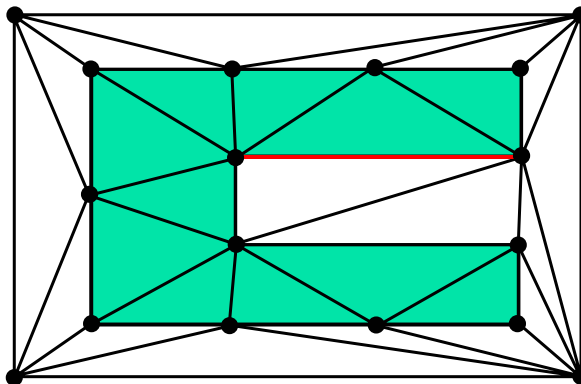
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Swap edges between adjacent pairs of triangles

Repeat till recover the boundary



## Boundary Recovery - Constrained

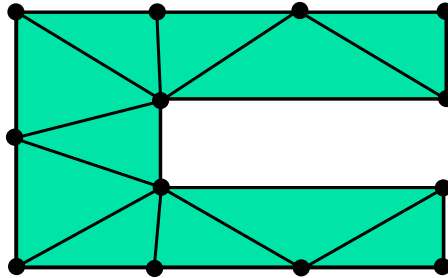


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- Swap edges between adjacent pairs of triangles
- Repeat till recover the boundary



## Boundary Recovery - Constrained



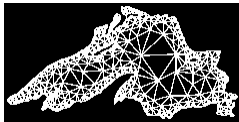
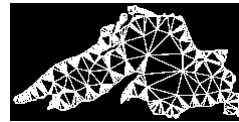
- Does not maintain Delaunay criterion !!!



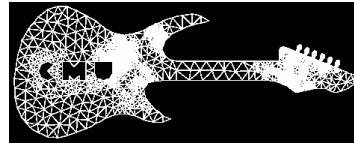
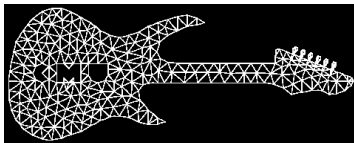
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## Examples



Boundary Conforming Meshes with different  
minimal angle constraints ( $0^\circ$ ,  $5^\circ$  &  $15^\circ$ )



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