

# CPSC 340: Machine Learning and Data Mining

Markov Chains

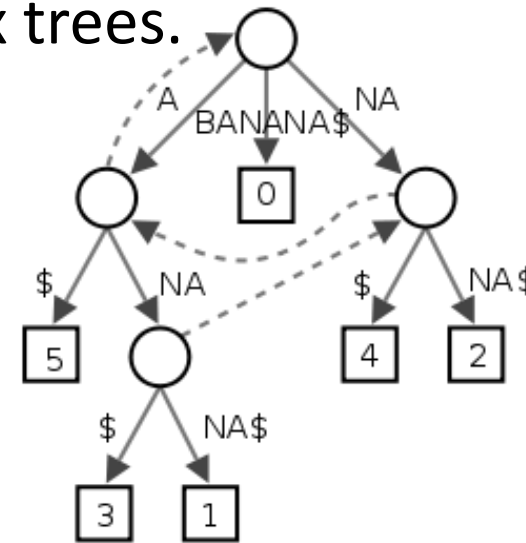
Fall 2015

# Admin

- Assignment 6 due Friday.
- Final exam details:
  - December 15: 8:30-11 (WESB 100).
  - 4 pages of cheat sheet allowed.
  - 9 questions.
  - Practice questions and list of topics posted.

# Last Time: Sequence Alignment

- We discussed **finding similar sequences** or **aligning sequence** parts.
  - Can find longest common substrings in linear time using suffix trees.
  - Using **dynamic programming**, we can compute ‘edit’ distance:
    - How many insertions/deletions/replacements to transform A into B?
  - **Local alignment**:
    - Find **local regions with small edit distance**.
  - BLAST:
    - Fast substring search to prune search for local alignments.
  - Multiple sequence alignment:
    - Hierarchical clustering helps align multiple sequences.



Smith-Waterman Scoring

	D	E	-	S				
D	D	0	5	4	3	4	4	3
E	E	0	4	10	9	8	7	6
A	A	0	3	9	9	8	7	6
S	S	0	2	8	14	13	12	11

Match = +5  
 Mismatch = -1  
 Gap = -1

Aligned:  
 1: DESIGN 1: DE-S  
 2: IDEAS 2: || |  
 2: DEAS

# Last Time: Dynamic Programming

- **Dynamic programming:**
  - Solves seemingly exponential-sized problems in polynomial-time.
- **3 ingredients:**
  1. Given results of **recursive calls**, can solve problem efficiently.
  2. **Limited number of possible arguments** recursive calls.
  3. **Memorize the results of recursive calls.**
- **Standard ways to implement dynamic programming:**
  1. Start from final result, use recursion but 'check' if result is in global table.
  2. Bottom-up: start filling out entries in the table in order.

# Example: Matrix Chain Multiplication

- Supposed we want to multiply matrices of different sizes:

- ABCDE.

- Cost depends on order of multiplication:

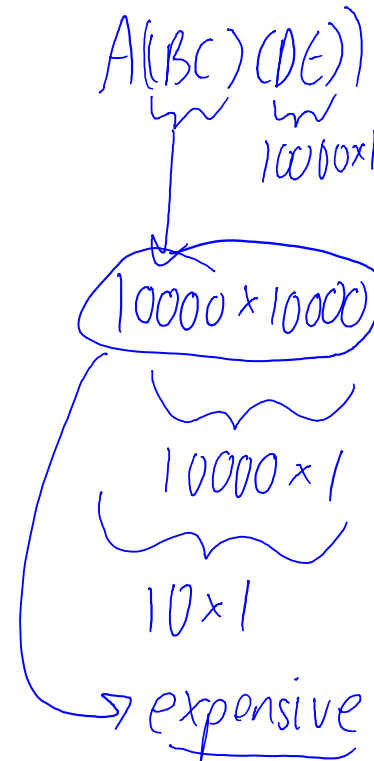
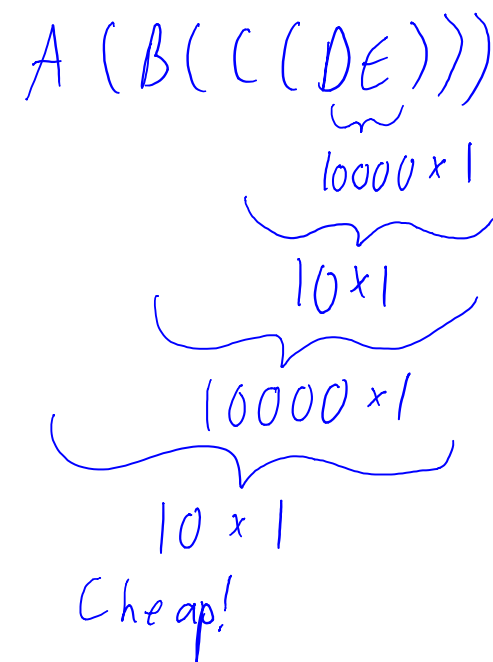
- $A(B(C(DE)))$  vs.  $A((BC)(DE))$ .

A:  $10 \times 10000$   
B:  $10000 \times 10$   
C:  $10 \times 10000$   
D:  $10000 \times 10$   
E:  $10 \times 1$

- What is the optimal order?

- There are an exponential number of possible orders.

- With 'n' matrices, we can solve this in  $O(n^3)$  using dynamic programming.



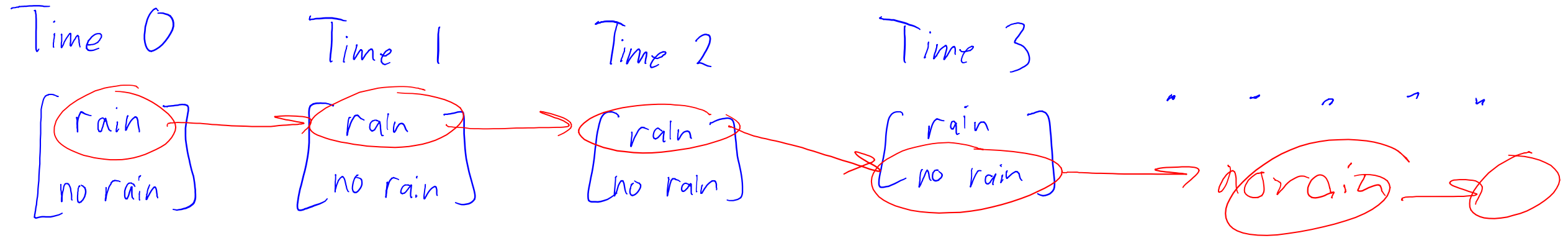
# Example: Matrix Chain Multiplication

- Define the solution recursively:
  - Cost(ABCDE) is minimum of:
    - Cost(ABCD) plus Cost(E) plus cost of combining result.
    - Cost(ABC) plus Cost(~~C~~DE) plus cost of combining result.
    - Cost(AB) plus Cost(CDE) plus cost of combining result.
    - Cost(A) plus Cost(BCDE) plus cost of combining result.
  - We could solve problem in  $O(n)$  given recursive calls.
- There are only  $O(n^2)$  possible recursive calls:
  - Ordering restricts possible recursions:
    - Cost(A), Cost(B), Cost(C), Cost(D), Cost(E),
    - Cost(AB), Cost(BC), Cost(CD), Cost(DE), Cost(ABC), Cost(BCD), Cost(CDE),
    - Cost(ABCD), Cost(BCDE).
- If you load results instead of re-computing, total cost is  $O(n^3)$ .

Cost( $X_1, X_2, X_3, \dots$ )  
{  
if (already computed)  
{  
load result  
return  
}  
else  
{  
do the work...  
store result

# Markov Chain Example

- Markov chains have a set of 'times' and possible 'states':



Our 'states' are 'rain' and 'not rain', at each time you have to

We define probabilities over initial state and transition between states. be in one of these states

$p(x_0 = \text{'rain'})$  is probability of it raining at time 0.

$p(x_t = \text{'rain'} | x_{t-1} = \text{'not rain'})$  is probability of it raining at time 't' if it was not raining at time t-1.

# Markov Chains

- Modeling the **probability of a sequence**  $x_0, x_1, x_2, x_3, \dots$ 
  - At ‘time’ 0, we have a **probability**  $p(x_0 = s)$  that ‘ $x_0$ ’ will be in each ‘**state**’ ‘ $s$ ’.
  - At ‘time’  $t$ , we have probability  $p_t(x_t = s_t \mid x_{t-1} = s_{t-1}, x_{t-2} = s_{t-2}, \dots, x_0 = s_0)$ :
    - Probability that ‘ $x_t$ ’ is in state ‘ $s_t$ ’, given what has happened so far.
  - Based on product rule:
$$p(x_t, x_{t-1}, \dots, x_0) = p(x_t \mid x_{t-1}, x_{t-2}, \dots, x_0) p(x_{t-1}, x_{t-2}, \dots, x_0)$$
$$= \prod_{i=0}^t [p(x_i \mid x_{i-1}, x_{i-2}, \dots, x_0)] p(x_0)$$
- **Markov chains** assume the **Markov property**:
  - Conditional independence assumption:  $x_t \perp x_{t-2}, x_{t-1}, \dots, x_0 \mid x_{t-1}$ :
    - $p_t(x_t = s_t \mid x_{t-1} = s_{t-1}, x_{t-2} = s_{t-2}, \dots, x_0 = s_0) = p_t(x_t = s_t \mid x_{t-1} = s_{t-1})$ .
  - ‘Memorylessness’:
$$\text{So } p(x_t, x_{t-1}, \dots, x_0) = p(x_0) \prod_{i=1}^t p(x_i \mid x_{i-1})$$
    - **Probability only depends on where you are now**, not where you were before.



# Markov Chain Examples

- We have already seen several examples:
  - PageRank, spectral clustering, graph-based SSL (A6Q2).

- PageRank example:

- Each 'state' is a webpage on the internet.
- We start on a random page:  $p(x_0 = s) = 1/|S|$ .

- E.g., [www.youtube.com](http://www.youtube.com) (not really random).

- At each time 't', we click on a random link:

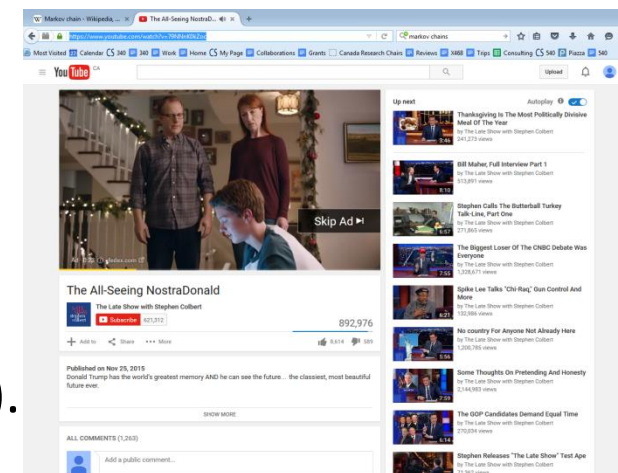
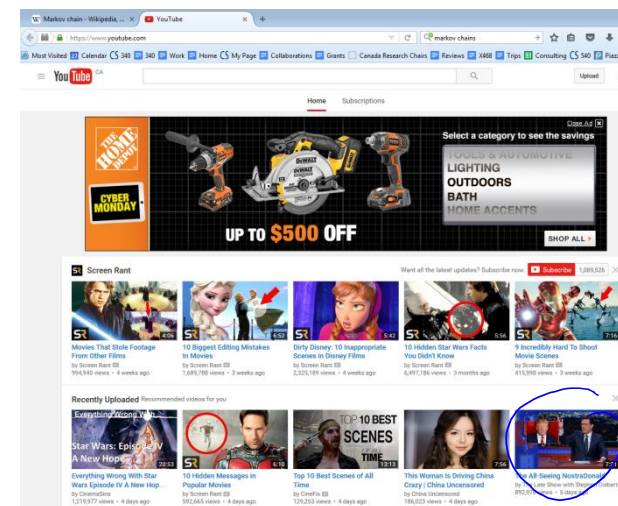
$$p(x_t = s_2 | x_{t-1} = s_1) = \frac{1}{d_{s_1}} \mathbb{I}[s_1 \text{ to } s_2 \text{ link exists}]$$

- If allow moving to new random page ('damping'):

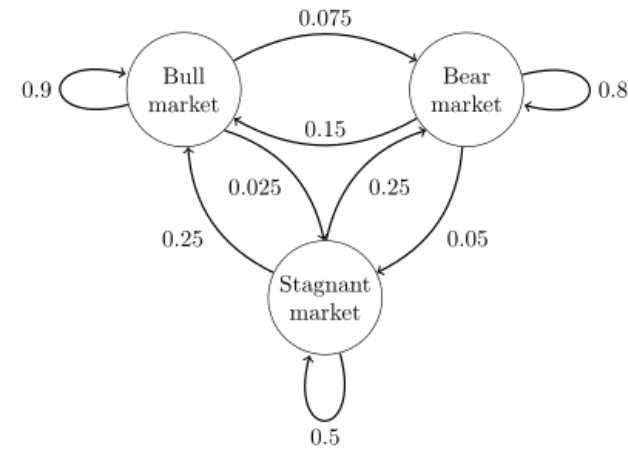
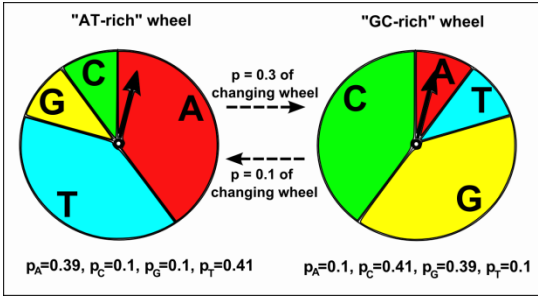
- Markov property is still satisfied.

- If time 't' depends on time 't-5':

- Markov property is not satisfied (can't look at browser history).



# Markov Chain Examples



- In our examples, probabilities were **homogeneous**:

$$p(x_t = s_2 | x_{t-1} = s) = p(x_i = s_2 | x_{i-1} = s_1)$$

for any  $i > 0$

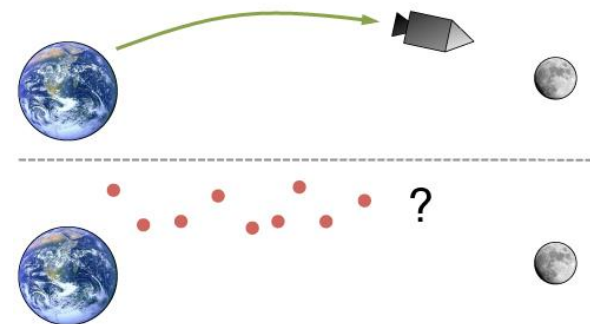
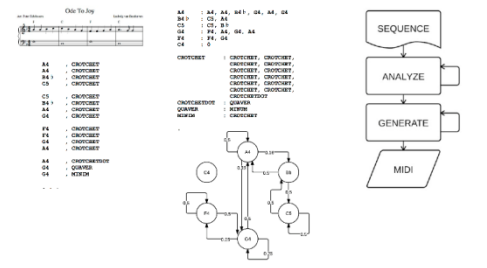
- Important special case.
- In practice, we **often allow time-dependent probabilities**.
- Incredible number of other applications:

– Bioinformatics, physics/chemistry, speech recognition, predator-prey models, language tagging/generation, computing integrals, economic models, tracking missiles/players, modeling music.

## Melody Generator

Generates a random melody using Markov Chains built from states and transitions extracted from an analysis of existing songs.

- Sequence
- Analysis
- Generate + Output



[https://en.wikipedia.org/wiki/Markov\\_chain](https://en.wikipedia.org/wiki/Markov_chain)  
<http://www.cs.uml.edu/ecg/index.php/AI/fall11/MarkovMelodyGenerator>  
<http://a-little-book-of-r-for-bioinformatics.readthedocs.org/en/latest/src/chapter10.html>  
<https://plus.maths.org/content/understanding-unseen>  
<http://www.cs.ubc.ca/~okumak/research.html>

# Markov Chain Tasks

- We are going to focus on **discrete time and states**.
- Common tasks we want to do with Markov chains:
  1. **Sampling**: given model, simulate from  $p(x_t, x_{t-1}, \dots, x_0)$ .
  2. **Learning**: estimating  $p(x_t = s_1 \mid x_t = s_2)$  to make model.
  3. **Inference**: given model, compute  $p_t(x_t = s)$ .
  4. **Stationary distribution**:  $p_\infty(x_\infty = s)$ .
  5. **Decoding**:  $\max_{y_1, y_2, \dots, y_t} p(x_t, x_{t-1}, x_{t-2}, \dots, x_0)$ .
  6. **Conditional inference**:  $p(x_t = s_1 \mid x_{t-1} = s_2, x_{t+10} = s_3)$ .

# Sampling from Markov Chains

- **Sampling** from a Markov chain:
  - Generate a sequence  $x_0, x_1, \dots, x_t$  following the joint distribution:
$$p(x_t, x_{t-1}, \dots, x_0) = p(x_0) \prod_{i=1}^t p(x_i | x_{i-1})$$
  - E.g., can we simulate a ‘random web surfer’?
- Easy for discrete time/states, a **random walk** model:
  - Generate  $x_0$  according to  $p(x_0)$ .
  - For  $i = 1, 2, \dots, t$ 
    - Given  $x_{i-1}$ , generate  $x_i$  according to  $p(x_i | x_{i-1})$ .

- **Why this works:** *Probability space*

$x_i = \text{'rain'}$	$x_i = \text{'not rain'}$	} $x_0 = \text{'not rain'}$
$x_i = \text{'rain'}$	$x_i = \text{'not rain'}$	

*Random walk gives each area the correct probability.*

# Learning Parameters of Markov Chain

- **Learning** in Markov chains:

- Given sample(s) of Markov chain, estimate what probabilities should be.

- **Maximum likelihood estimates:**

- $p(x_0 = s) = N(x_0 = s)/N$ .  
→ number of samples we have.  
→ number of times we start with  $x_0 = s$ .

- Inhomogeneous case:

- $p_t(x_t = s_2 | x_{t-1} = s_1) = N(x_{t-1} = s_1, x_t = s_2) / N(x_{t-1} = s_1)$ .  
→ number of times we were in  $s_1$  at time  $(t-1)$ .  
→ number of samples where we went from  $s_1$  to  $s_2$  at time  $(t-1)$ .

- Homogeneous case:

- $p(x_t = s_2 | x_{t-1} = s_1) = N(x_{i-1} = s_1, x_i = s_2) / N(x_{i-1} = s_1)$  for any 'i'.  
→ number of times we were in  $s_1$ .  
→ number of times we went from  $s_1$  to  $s_2$ .

- **Initial  $p(x_0 = s)$  and inhomogeneous case need multiple sequences.**

- Need a lot of data if number of states is very large.

# Inference in Markov Chains

- **Inference**: compute probability of being in state 's' at time 't'.
- We are given this for time 0, what about time 't'?
- Inference in length-2 chain:

$$p(x_1) = \sum_{x_0} p(x_1, x_0) \quad (\text{marginalization rule})$$

$$= \sum_{x_0} p(x_1 | x_0) p(x_0) \quad (\text{product rule})$$

We sum over values of  $x_0$ , summing probability of leading to  $x_1$ .

Matrix notation:

Let  $u_0$  be a row-vector with elements  $p(x_0)$ .

Let  $P$  be a matrix with elements  $P_{ij} = p(x_1 = j | x_0 = i)$ .

Let  $u_1$  be a row-vector with elements  $p(x_1)$ .

Then  $u_1 = u_0 P$ .



# Inference in Markov Chains

- Inference in length-3 chain:

$$p(x_1) = \sum_{x_0} \sum_{x_2} p(x_2, x_1, x_0) \quad (\text{marg. rule})$$

$$= \sum_{x_0} \sum_{x_2} p(x_2 | x_1, x_0) p(x_1 | x_0) p(x_0) \quad (\text{product rule})$$

$$= \sum_{x_0} \sum_{x_2} p(x_2 | x_1) p(x_1 | x_0) p(x_0) \quad (\text{Markov property})$$

$$= \sum_{x_0} p(x_1 | x_0) p(x_0) \left( \sum_{x_2} p(x_2 | x_1) \right) \quad (\text{distributive law: } \sum_i c a_i = c \sum_i a_i)$$

$$= \sum_{x_0} p(x_1 | x_0) p(x_0) \quad (\text{probabilities sum to 1})$$

Note: same as length-2 result.

- possible 'futures' do not change present.

- to get all probabilities,

$$u_1 = u_0 P$$

# Inference in Markov Chains

- Inference in length-3 chain:

$$\begin{aligned} p(x_2) &= \sum_{x_1} \sum_{x_0} p(x_2, x_1, x_0) \quad (\text{marg. rule}) \\ &= \sum_{x_1} \sum_{x_0} p(x_2 | x_1, x_0) p(x_1 | x_0) p(x_0) \quad (\text{prod. rule}) \\ &= \sum_{x_1} \sum_{x_0} p(x_2 | x_1) p(x_1 | x_0) p(x_0) \quad (\text{Markov}) \\ &= \sum_{x_1} p(x_2 | x_1) \sum_{x_0} p(x_1 | x_0) p(x_0) \quad (\text{distr. law}) \\ &= \sum_{x_1} p(x_2 | x_1) p(x_1) \quad (\text{definition of } x_1) \end{aligned}$$

Note: once you have  $p(x_1)$  for all values of  $x_1$ , it's easy to compute  $p(x_2)$ .

In matrix notation:

$$\begin{aligned} \mu_2 &= \mu_1 P \\ &= \mu_0 P^2 \\ &= \mu_0 P^2 \end{aligned} \quad \text{"Chapman-Kolmogorov equation"}$$

$$\text{Similarly, } p(x_3) = \sum_{x_2} p(x_3 | x_2) p(x_2)$$

$$\mu_3 = \mu_0 P^3$$

(If inhomogeneous,  $\mu_3 = \mu_0 P_1 P_2 P_3$ )



# Stationary Distribution of Markov Chains

- After 't' steps, the distribution of states is given by matrix power:

$$\mu_t = \mu_0 P^t$$

- After each step, we start 'forgetting' initial state.
- **Stationary distribution** is a steady-state:

$$\mu_\infty = \mu_\infty P$$

- A 'row' eigenvector of 'P' with an eigenvalue of 1.
  - Normalized to sum to 1.
- Often approximated using power of P (PageRank: 'power method').
- If  $P_{ij} > 0$ , guarantees **unique** stationary distribution.
  - Otherwise, could have multiple possibilities.
  - Many other conditions guarantee uniqueness.

# Decoding in Markov Chains

- A subtle difference:
  - Inference let's us compute  $p(x_t = s)$  for any 't' and 's'.
  - This lets us find the most **likely state at each time**.
  - But probability of states is dependent, **may not be mostly likely sequence**.
- Finding most likely sequence is called **decoding**:

$$\operatorname{argmax}_{x_0, x_1, \dots, x_t} \left\{ p(x_t, x_{t-1}, \dots, x_0) \right\}$$

- Seems harder than inference:
  - optimal sequence of length 3 may not contain optimal length 2 sequence.
  - but we can solve this problem with **dynamic programming**.

# Decoding in Markov Chains

- Let's consider the best length-2 sequence that ends in state 'x<sub>1</sub>':

$$V_1(x_1) = \max_{x_0} p(x_1, x_0) = \max_{x_0} p(x_1 | x_0) p(x_0)$$

*Handwritten notes:* "state" (with arrow pointing to x<sub>1</sub>), "value" (with arrow pointing to the max operator), "time" (with arrow pointing to the subscript 1).

"value" of best sequence up to time '1' that ends in state x<sub>1</sub>.

- Now let's consider a length-3 sequence that ends in state 'x<sub>2</sub>':

$$V_2(x_2) = \max_{x_1, x_0} p(x_2, x_1, x_0)$$

$$= \max_{x_1} \left\{ \max_{x_0} p(x_2, x_1, x_0) \right\}$$
 ("max of max" trick)

$$= \max_{x_1} \left\{ \max_{x_0} p(x_2 | x_1) p(x_1 | x_0) p(x_0) \right\}$$
 (prod. rule and Markov)

$$= \max_{x_1} \left\{ p(x_2 | x_1) \max_{x_0} p(x_1 | x_0) p(x_0) \right\} = \max_{x_1} \left\{ p(x_2 | x_1) V_1(x_1) \right\}$$

*Handwritten notes:* "Definition of V<sub>1</sub>(x<sub>1</sub>)" (with arrow pointing to the V<sub>1</sub>(x<sub>1</sub>) term in the final equation).

$$\max_i a b_i = a \max_i b_i$$
 (for a ≥ 0)

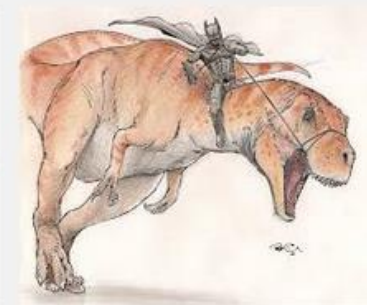
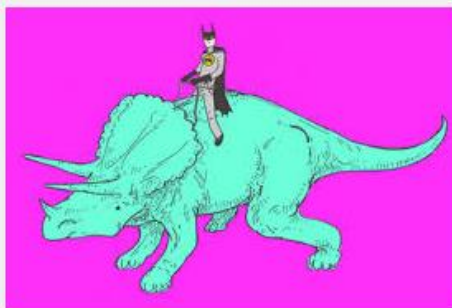
# Decoding Markov Chains

- General formula:  $V_0(x_0) = p(x_0)$   
 $V_t(x_t) = \max_{x_{t-1}} p(x_t | x_{t-1}) V_{t-1}(x_{t-1})$

- We have the ingredients for dynamic programming:
  - Efficiently compute solution given recursion result:

$$\max_{x_0, x_1, \dots, x_t} \{ p(x_t, x_{t-1}, \dots, x_0) \} = \max_{x_t} \{ V_t(x_t) \}$$

- Number number of possible recursions:  $V_{t'}(x_{t'})$  for each time  $t'$  and state  $x_{t'}$
  - If we memorize results of recursions, we can solve this efficiently.



# Summary

- **Markov chains** used for sequences/time-series/random-walks.
- **Sampling** is task of simulating sequence according to model.
  - Done by running a **random walk**.
- **Learning** is task of estimating parameters from sequences.
  - Done by **counting**.
- **Inference** is task of computing probabilities at particular times.
  - Done by **matrix multiplication**.
- **Stationary distribution** is steady-state after running for a long time.
  - Done by **normalizing eigenvector** with largest eigenvalue.
- **Decoding** is task of computing most likely sequence of states.
  - Done by **dynamic programming**.
- Next time: how Markov models are actually useful (HMMs and belief nets).