CPSC 504 – Background
(aka, all you need to know about databases for this course in two lectures)

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Overview of the next two classes
- Entity Relationship (ER) diagrams
- Relational databases
- Object Oriented Databases (OODBs)
- XML
- Other data types
- Database internals (Briefly)

Levels of Abstraction
- A major purpose of a DB system is to provide an abstract view of the data.
- Three abstraction levels:
  - Physical level: how data are actually stored
  - Conceptual (or Logical) level: how data is perceived by the users
  - External (or View) level: describes different part of the database to different users
    - Convenience, security, etc.
    - Compare views of student, registrar, & database admin.

Schema and Instances
- We create the schema – the logical structure of the database (e.g., students take courses)
  - Conceptual (or logical) schema: db design at the logical level
  - Physical schema: db design at the physical level: indexes, etc
- Later we’ll populate instances – the actual content of the database at a particular point in time
  - E.g., currently there are no grades for CPSC 504 Winter term 2
- Physical Data Independence – the ability to modify the physical schema without changing the logical schema
  - Applications depend on the conceptual schema
- Logical Data Independence – Provided by the views
  - Ability to change the conceptual scheme without changing the applications

Conceptual Database Design
- What are the entities and relationships in the enterprise?
  - Entitles are usually nouns, e.g., “course” “prof”
  - Relationships are statements about 2 or more objects. Often, verbs., e.g., “a prof teaches a course”
- What information about these entities and relationships should we store in the database?
- What integrity constraints or other rules hold?
- In relational databases, this data is generally encoded in an Entity-Relationship (ER) Diagram
Entity / Relationship Diagrams

Entities
- Product
Attributes
- address
Relationships between entities
- buys

Keys in E/R Diagrams
- Every entity set must have a key which is identified by an underline

Roles in Relationships
- What if we need an entity set twice in one relationship?

Attributes on Relationships

Subclasses in E/R Diagrams
- EducationalProduct
- SoftwareProduct
- Platforms
- AgeGroup
Summarizing ER diagrams
- Basics: entities, relationships, and attributes
- Also showed inheritance
- Has things other things like cardinality
- Used to design databases...

But how do you store data in them?

Overview of the next two classes
- Entity Relationship (ER) diagrams
- Relational databases
  - How did we get here?
  - What’s in a relational schema?
  - From ER to relational
  - Query Languages
- Object Oriented Databases (OODBs)
- XML
- Other data types
- Database internals (Briefly)

How did we get the relational model?
- Prior to the relational model, there were two main contenders
  - Network databases
  - Hierarchical databases
- Network databases had a complex data model
- Hierarchical databases integrated the application in the data model

Example Hierarchical Model

Example IMS (Hierarchical) query: Print the names of all the provinces admitted during a Liberal Government

Relational model to the rescue!
- Introduced by Edgar Codd (IBM) in 1970
- Most widely used model today.
  - Vendors: IBM, Informix, Microsoft, Oracle, Sybase, etc.
- Competitor: object-oriented model
  - ObjectStore, Versant, Ontos
- A synthesis emerging: object-relational model
  - Informix Universal Server, UniSQL, O2, Oracle, DB2
- Recent competitor: XML data model
Key points of the relational model

- Exceedingly simple to understand – main abstraction is represented as a table
- Simple query language separate from application language
- Lots of bells and whistles to do complicated things

Structure of Relational Databases

- **Relational database**: a set of relations
- **Relation**: made up of 2 parts:
  - **Schema**: specifies name of relation, plus name and domain (type) of each field (or column or attribute).
    - e.g., `Student` (id: string, name: string, address: string, phone: string, major: string).
  - **Instance**: a table with rows and columns. #Rows = cardinality, #fields = dimension / arity / degree
- **Relational Database Schema**: collection of schemas in the database
- **Database Instance**: a collection of instances of its relations (e.g., currently no grades in CPSC 504)

Example of a Relation Instance

<table>
<thead>
<tr>
<th>Product</th>
<th>Name</th>
<th>Price</th>
<th>Category</th>
<th>Manufacturer</th>
</tr>
</thead>
<tbody>
<tr>
<td>gizmo</td>
<td></td>
<td>$19.99</td>
<td>gadgets</td>
<td>GizmoWorks</td>
</tr>
<tr>
<td>Power gizmo</td>
<td>$29.99</td>
<td>gadgets</td>
<td>GizmoWorks</td>
<td></td>
</tr>
<tr>
<td>SingleTouch</td>
<td>$149.99</td>
<td>photography</td>
<td>Canon</td>
<td></td>
</tr>
<tr>
<td>MultiTouch</td>
<td>$203.99</td>
<td>household</td>
<td>Hitachi</td>
<td></td>
</tr>
</tbody>
</table>

Tuples or rows
Relation or table
Order of rows isn't important
Formal Definition:
Product(Name: string, Price: double, Category: string, Manufacturer: string)

Overview of the next two classes

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  - What's in a relational schema?
  - From ER to relational
  - Query Languages
- Object Oriented Databases (OODBs)
- XML
- Other data types
- Database internals (Briefly)

From E/R Diagrams to Relational Schema

- Entity set \(\rightarrow\) relation
- Relationship \(\rightarrow\) relation

Entity Set to Relation

<table>
<thead>
<tr>
<th>name</th>
<th>category</th>
<th>price</th>
</tr>
</thead>
<tbody>
<tr>
<td>gizmo</td>
<td>gadgets</td>
<td>$19.99</td>
</tr>
</tbody>
</table>
Overview of the next two classes
- Entity Relationship (ER) diagrams
- Relational databases
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  - What’s in a relational schema?
  - From ER to relational
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To Find products where the manufacturer is GizmoWorks:

<table>
<thead>
<tr>
<th>Product</th>
<th>Name</th>
<th>Price</th>
<th>Category</th>
<th>Manufacturer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Gizmo</td>
<td>$19.99</td>
<td>Gadgets</td>
<td>GizmoWorks</td>
<td></td>
</tr>
<tr>
<td>Powergizmo</td>
<td>$29.99</td>
<td>Gadgets</td>
<td>GizmoWorks</td>
<td></td>
</tr>
<tr>
<td>MultiTouch</td>
<td>$203.99</td>
<td>Household</td>
<td>Hitachi</td>
<td></td>
</tr>
</tbody>
</table>

Selection: $\sigma_{\text{Manufacturer} = \text{GizmoWorks}}$ Product
Find the Name, Price, and Manufacturers of products whose price is greater than 100

<table>
<thead>
<tr>
<th>Product</th>
<th>Price</th>
<th>Category</th>
<th>Manufacturer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Gizmo</td>
<td>$19.99</td>
<td>Gadget</td>
<td>Gizmo/Works</td>
</tr>
<tr>
<td>Powergizmo</td>
<td>$29.99</td>
<td>Gadget</td>
<td>Gizo/Works</td>
</tr>
<tr>
<td>SingleTouch</td>
<td>$149.99</td>
<td>Photography</td>
<td>Canon</td>
</tr>
<tr>
<td>MultiTouch</td>
<td>$293.99</td>
<td>Household</td>
<td>Hitachi</td>
</tr>
</tbody>
</table>

Selection + Projection:
π_{Name, Price, Manufacturer}(π_{Price > 100}Product)

<table>
<thead>
<tr>
<th>Name</th>
<th>Price</th>
<th>Manufacturer</th>
</tr>
</thead>
<tbody>
<tr>
<td>SingleTouch</td>
<td>$149.99</td>
<td>Canon</td>
</tr>
<tr>
<td>MultiTouch</td>
<td>$293.99</td>
<td>Hitachi</td>
</tr>
</tbody>
</table>

Find the product names and price of products that cost less than $200 and have manufacturers where there is a Company that has a CName that matches the manufacturer, and its country is Japan

<table>
<thead>
<tr>
<th>Product</th>
<th>Price</th>
<th>Company</th>
<th>CName</th>
<th>Country</th>
</tr>
</thead>
<tbody>
<tr>
<td>Gizmo</td>
<td>$19.99</td>
<td>Gizo</td>
<td>Gizo</td>
<td>Gizo/Works</td>
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<tr>
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<tr>
<td>SingleTouch</td>
<td>$149.99</td>
<td>Canon</td>
<td>Canon</td>
<td>Japan</td>
</tr>
<tr>
<td>MultiTouch</td>
<td>$293.99</td>
<td>Hitachi</td>
<td>Hitachi</td>
<td>Japan</td>
</tr>
</tbody>
</table>

π_{Name, Price}(σ_{Price < 200}Product)\left(\pi_{CName = \text{Manufacturer}}(\pi_{Country = \text{Japan}}\text{Company})\right)

When are two relations related?
- You guess they are
- I tell you so
- Constraints say so
- A key is a set of attributes whose values are unique; we underline a key Product{Name, Price, Category, Manufacturer}
- Foreign keys are a method for schema designers to tell you so
  - A foreign key states that an attribute is a reference to the key of another relation ex: Product.Manufacturer is foreign key of Company
- Gives information and enforces constraint

The SQL Query Language
- Developed by IBM (System R) in the 1970s
- Standards:
  - SQL-86
  - SQL-89 (minor revision)
  - SQL-92 (major revision, current standard)
  - SQL-99 (major extensions)

SQL
- Data Manipulation Language (DML)
  - Query one or more tables
  - Insert/delete/modify tuples in tables
- Data Definition Language (DDL)
  - Create/alter/delete tables and their attributes
- Transact-SQL
  - Idea: package a sequence of SQL statements ➔ server

Querying – SQL
Standard language for querying and manipulating data

Structured Query Language

Many standards out there:
- ANSI SQL
- SQL92 (a.k.a. SQL2)
- SQL99 (a.k.a. SQL3)
- Vendors support various subsets of these
- What we discuss is common to all of them
### SQL basics

- **Basic form:** (many many more bells and whistles in addition)

  Select attributes  
  From relations (possibly multiple, joined)  
  Where conditions (selections)

### SQL – Selections

**SELECT** *  
**FROM** Company  
**WHERE** country="Canada" AND stockPrice > 50

Some things allowed in the **WHERE** clause:  
- attribute names of the relation(s) used in the **FROM**.  
- comparison operators: =, <>, <, >, <=, >=  
- apply arithmetic operations: stockPrice*2  
- operations on strings (e.g., "||" for concatenation).  
- Lexicographic order on strings.  
- Pattern matching: s LIKE p  
- Special stuff for comparing dates and times.

### SQL – Projections

Select only a subset of the attributes  

**SELECT** name, stock price  
**FROM** Company  
**WHERE** country="Canada" AND stockPrice > 50

Rename the attributes in the resulting table  

**SELECT** name AS company, stockPrice AS price  
**FROM** Company  
**WHERE** country="Canada" AND stockPrice > 50

### SQL – Joins

**SELECT** name, store  
**FROM** Person, Purchase  
**WHERE** name=buyer AND city="Vancouver" AND product="gizmo"

Product (name, price, category, maker)  
Purchase (buyer, seller, store, product)  
Company (name, stock price, country)  
Person(name, phone number, city)

### Selection:

**σ**<sub>Manufacturer = GizmoWorks</sub> (Product)

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What’s the SQL?

### Selection + Projection:

**π**<sub>Name, Price, Manufacturer (**σ**Price > 100 Product)

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What’s the SQL?
What’s the SQL?

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π Name, Price (σ Price <= 200 Product) ≠ Cname (σ Country = 'Japan' Company)

Product

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Company

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What’s the Datalog?

π Name, Price (σ Price <= 200 Product) ≠ Cname (σ Country = 'Japan' Company)

π Name, Price (σ Price <= 200 Product)

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π Name, Price, Manufacturer (σ Price > 100 Product)

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<tr>
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</tr>
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</table>

Querying – Datalog (Our final query language)

- Enables expressing recursive queries
- More convenient for analysis
- Some people find it easier to understand
- Without recursion but with negation it is equivalent in power to relational algebra and SQL
- Limited version of Prolog (no functions)

Conjunctive Queries

- A subset of Datalog
- Only relations appear in the right hand side of rules
- No negation
- Functionally equivalent to Select, Project, Join queries
- Very popular in modeling relationships between databases

Dataylog Rules and Queries

A datalog rule has the following form:

head :- atom1, atom2, ..., atom,...
then :- if ...

Arithmetic comparison or interpreted predicate

ExpensiveProduct(N) :- Product(N,M,P) & P > $100
CanadianProduct(N) :- Product(N,M,P) & Company(M, “Canada”, SP)
IntlProd(N) :- Product(N,M,P) & NOT Company(M, “Canada”, SP)

Negated subgoal - also denoted by ~

Selection:

σ Manufacturer = GizmoWorks (Product)

Selection + Projection:

π Name, Price, Manufacturer (σ Price > 100 Product)
π_{Name, Price} ((σ_{Price <= 200} Product) \times (σ_{Country = 'Japan'} Company))

**Bonus Relational Goodness: Views**

Views are stored queries treated as relations, Virtual views are not physically stored. Materialized views are stored.

They are used (1) to define conceptually different views of the database and (2) to write complex queries simply.

View: purchases of telephony products:

```
CREATE VIEW telephony-purchases AS
SELECT product, buyer, seller, store
FROM Purchase, Product
WHERE Purchase.product = Product.name
AND Product.category = 'telephony'
```

---

**Summarizing Relational DBs**

- Relational perspective: Data is stored in relations. Relations have attributes. Data instances are tuples.
- SQL perspective: Data is stored in tables. Tables have columns. Data instances are rows.
- Query languages:
  - Relational algebra – mathematical base for understanding query languages
  - SQL – very widely used
  - Datalog – based on Prolog, very popular with theoreticians
- Views allow complex queries to be written simply

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**Object-Oriented DBMS’s**

- Started late 80’s
- Main idea:
  - Toss the relational model!
  - Use the OO model – e.g. C++ classes
- Standards group: ODMG = Object Data Management Group.
- OQL = Object Query Language, tries to imitate SQL in an OO framework.

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**Outline**

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- XML
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**The OO Plan**

ODMG imagines OO-DBMS vendors implementing an OO language like C++ with extensions (OQL) that allow the programmer to transfer data between the database and “host language” seamlessly.

A brief diversion: the impedance mismatch
OO Implementation Options

- Build a new database from scratch (O2)
- Elegant extension of SQL
- Later adopted by ODMG in the OQL language
- Used to help build XML query languages
- Make a programming language persistent (ObjectStore)
- No query language
- Niche market
- ObjectStore is still around, renamed to Exelon, stores XML objects now

ODL

- ODL is used to define persistent classes, those whose objects may be stored permanently in the database.
- ODL classes look like Entity sets with binary relationships, plus methods.
- ODL class definitions are part of the extended, OO host language.

ODL – remind you of anything?

```java
interface Person
{
    (extent People key sin)
    attribute string sin;
    attribute string dept;
    attribute string name;
}
```

```java
interface Course
{
    (extent Crs key cid)
    attribute string cid;
    attribute string cname;
    relationship Person instructor;
    relationship Set<Student> stds inverse takes;
}
```

```java
interface Student extends Person
{
    (extent Students)
    attribute string major;
    relationship Set<Course> takes inverse stds;
}
```

Why did OO Fail?

- Why are relational databases so popular?
  - Very simple abstraction; don’t have to think about programming when storing data.
  - Very well optimized
- Relational db are very well entrenched – not enough advantages, and no good exit strategy (we’ll see more about this)

Merging Relational and OODBs

- Object-oriented models support interesting data types – not just flat files.
- Maps, multimedia, etc.
- The relational model supports very-high-level queries.
- Object-relational databases are an attempt to get the best of both.
- All major commercial DBs today have OR versions – full spec in SQL99, but your mileage may vary.

Outline

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**XML**
- eXtensible Markup Language
- XML 1.0 – a recommendation from W3C, 1998
- Roots: SGML (from document community - works great for them; from db perspective, very nasty).
- After the roots: a format for sharing data

**Why XML is of Interest to Us**
- XML is just syntax for data
- Note: we have no syntax for relational data
- But XML is not relational: *semistructured*
- This is exciting because:
  - Can translate *any* data to XML
  - Can ship XML over the Web (HTTP)
  - Can input XML into any application
- Thus: data sharing and exchange on the Web

**XML Data Sharing and Exchange**

**From HTML to XML**

HTML describes the presentation

**HTML**
```html
<h1>Bibliography</h1>
<p><i>Foundations of Databases</i>  
Abiteboul, Hull, Vianu  
Addison Wesley, 1995</p>
<p><i>Data on the Web</i>  
Abiteboul, Buneman, Suciu  
Morgan Kaufmann, 1999</p>
```

**XML**
```xml
<bibliography>
  <book>
    <title>Foundations…</title>
    <author>Abiteboul</author>
    <author>Hull</author>
    <author>Vianu</author>
    <publisher>Addison Wesley</publisher>
    <year>1995</year>
  </book>
  ...
</bibliography>
```

XML describes the content
XML Document

XML Terminology
- Elements
  - enclosed within tags:
    - <person> ... </person>
  - nested within other elements:
    - <person> <address> ... </address> </person>
  - can be empty
    - <married/> abbreviated as <married/>
  - can have Attributes
    - <person id="0005"> ... </person>
- XML document has as single ROOT element

Buzzwords
- What is XML?
  - W3C data exchange format
  - Hierarchical data model
  - Self-describing
  - Semi-structured

XML as a Tree !!

XML is self-describing
- Schema elements become part of the data
  - In XML <persons>, <name>, <phone> are part of the data, and are repeated many times
- Relational schema: persons(name,phone) defined separately for the data and is fixed
- Consequence: XML is much more flexible

Relational Data as XML

<table>
<thead>
<tr>
<th>name</th>
<th>phone</th>
</tr>
</thead>
<tbody>
<tr>
<td>John</td>
<td>3634</td>
</tr>
<tr>
<td>Sue</td>
<td>6343</td>
</tr>
<tr>
<td>Dick</td>
<td>6363</td>
</tr>
</tbody>
</table>
XML is semi-structured

- Missing elements:
  ```xml
  <person>
    <name>John</name>
    <phone>1234</phone>
  </person>
  <person>
    <name>Joe</name>
  </person>
  ← no phone!
  ```

- Could represent in a table with nulls
<table>
<thead>
<tr>
<th>name</th>
<th>phone</th>
</tr>
</thead>
<tbody>
<tr>
<td>John</td>
<td>1234</td>
</tr>
<tr>
<td>Joe</td>
<td>-</td>
</tr>
</tbody>
</table>

XML is semi-structured

- Repeated elements
  ```xml
  <person>
    <name>Mary</name>
    <phone>2345</phone>
    <phone>3456</phone>
  </person>
  ← two phones!
  ```

- Impossible in tables:
<table>
<thead>
<tr>
<th>name</th>
<th>phone</th>
</tr>
</thead>
<tbody>
<tr>
<td>Mary</td>
<td>2345</td>
</tr>
<tr>
<td></td>
<td>3456</td>
</tr>
<tr>
<td></td>
<td>??</td>
</tr>
</tbody>
</table>

XML is semi-structured

- Elements with different types in different objects
  ```xml
  <person>
    <name>John</name><first>John</first>
    <last>Smith</last>
    <phone>1234</phone>
  </person>
  ← structured name!
  ```

- Heterogeneous collections:
  - `<persons>` can contain both `<person>`s and `<customer>`s

So how would you store XML in a relational db?

Summarizing XML

- XML has first class elements and second class attributes
- XML is semi-structured
- XML is nested
- XML is a tree
- XML is a huge buzzword

Will XML replace relational databases?

Outline

- Entity Relationship (ER) diagrams
- Relational databases
- Object Oriented Databases (OODBs)
- XML
- Other data types
- Database internals (Briefly)

Other data formats

- Makefiles
- Forms
- Application code

What format is your data in?
Outline

- Entity Relationship (ER) diagrams
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- Database internals (Briefly)
  - Query Optimization & Execution
  - Transaction Processing

Overview of Query Optimization

- Plan: Tree of ordered Relational Algebra operators and choice of algorithm for each operator
- Two main issues:
  - For a given query, what plans are considered?
  - Algorithm to search plan space for cheapest (estimated) plan.
- How is the cost of a plan estimated?
- Ideally: Want to find best plan.
  - Practically: Avoid worst plans.
- Some tactics
  - Do selections early
  - Use materialized views
  - Use Indexes

How SQL Gets Executed: Query Execution Plans

Select Name, Price
From Product, Company
Where Manufacturer = Cname
AND Price <= 200
AND Country = 'Japan'

Query optimization also specifies the algorithms for each operator; then queries can be executed

Tree-Based Indexes

- "Find all students with gpa > 3.0"
  - If data is in sorted file, do binary search to find first such student, then scan to find others.
  - Cost of binary search can be quite high.
- Simple idea: Create an 'index' file.

Can do binary search on (smaller) index file!

Tree-Based Indexes (2)

B+ Tree: The Most Widely Used Index

- Insert/delete at \( \log_F N \) cost; keep tree height-balanced. (\( F = \) fanout, \( N = \) # leaf pages)
- Minimum 50% occupancy (except for root). Each node contains \( d \leq \frac{m}{2} \leq 2d \) entries. The parameter \( d \) is called the order of the tree.
Example B+ Tree

- Search begins at root, and key comparisons direct it to a leaf.
- Search for 5*, 15*, all data entries >= 24*

Query Execution

- Now that we have the plan, what do we do with it?
  - How do joins work?
  - How do deal with paging in data, etc.
  - New research covers new paradigms where interleaved with optimization

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Transactions

- Address two issues:
  - Access by multiple users
    - Remember the “client-server” architecture: one server with many clients
  - Protection against crashes

Transaction States

- A transaction can be in one of the following states:
  - active:
    - makes progress or waits for resources; the initial state
  - failed:
    - normal execution cannot continue; it may occur because of:
      - system crash
      - cancellation by the user
  - aborted:
    - DBMS cancels it due to problem in execution (e.g., Consistency)
  - committed:
    - after successful completing a “commit” command
    - to undo its effects we need to run a compensating transaction
- Two options for aborted transactions:
  - restart it as a new transaction later (e.g., system failures)
  - kill it (e.g., internal logical errors)
- Failed transactions are eventually aborted

Transactions

- Transaction = group of statements that must be executed atomically

Transaction Properties: ACID

- Atomicity: either all or none of the operations are completed
- Consistency: preserves database integrity
- Isolation: concurrent transactions must not interfere with each other
- Durability: changes from successful transactions must persist through failures

Search for 5*, 15*, all data entries >= 24*
Transaction Example

Consider two transactions:

T1: READ(A) A=A+100 WRITE(A) READ(B) B=B+100 WRITE(B)

T2: READ(A) A=A+100 WRITE(A) READ(B) B=B+100 WRITE(B)

Intuitively, the first transaction transfers $100 to A's account from B's account. The second credits both accounts with a 10% interest payment.

No guarantee that T1 executes before T2 or vice-versa. However, the end effect must be equivalent to these two transactions running serially in some order:

T1, T2 or T2, T1

Transactions: Serializability

Serializability = the technical term for isolation

An execution is serial if it is completely before or completely after any other function's execution

An execution is serializable if it equivalent to one that is serial

DBMS can offer serializability guarantees

Serializability Example

Enforced with locks, like in Operating Systems!

But this is not enough:

User 1

LOCK A [write A=1]
UNLOCK A

...

LOCK B [write B=2]
UNLOCK B

User 2

LOCK A [write A=3]
UNLOCK A

...

LOCK B [write B=4]
UNLOCK B

Okay, but what if it crashes?

Enforcing Atomicity & Durability

Atomicity:

Transactions may abort; Need to rollback changes

Durability:

What if DBMS stops running? Need to "remember" committed changes.

Basic idea: use steal and no-force

What to do?

Basic idea: use steal and no-force

Keep a log that talks about what's happened

Make checkpoints where write down everything that's actually happened

After a crash: assure Atomicity and Durability by keeping all committed transactions and getting rid of actions of uncommitted transactions

Handling the Buffer Pool

Transactions modify pages in memory buffers

Writing to disk is more permanent

When should updated pages be written to disk?

Force every write to disk?

- Poor response time.
- But provides durability.

Steal buffer-pool frames from uncommitted Xacts?
(resulting in write to disk)

- If not, poor throughput.
- If so, how can we ensure atomicity?

What to do?

Basic idea: use steal and no-force

Keep a log that talks about what's happened

Make checkpoints where write down everything that's actually happened

After a crash: assure Atomicity and Durability by keeping all committed transactions and getting rid of actions of uncommitted transactions
### Crash Recovery: Big Picture

- Start from a checkpoint (Three phases. Need to:
  - Figure out which Xacts committed since checkpoint, which failed (Analysis)
  - REDO all actions
  - UNDO effects of failed Xacts

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### Now what?

- Time to read papers
- Prepare paper responses – it’ll help you focus on the paper, and allow for the discussion leader to prepare better discussion
- You all have different backgrounds, interests, and insights. Bring them into class!