Community Search and Cocktail Party Planning

Mauro Sozio and Aris Gionis. The community-search problem and how to plan a successful cocktail party.
KDD 2010.
Planning a cocktail party
Planning a cocktail party
Recipe for a successful party:

- Participants should be “close” to the organizers (e.g., a friend of a friend).
- Everybody should know sufficiently many in the party (on an average?).
- The graph should be connected.
- The number of participants should not be too small but...
- …not too large either!!!
- …
- social distance not too large.

Not an easy task…
The problem: find the community that a given set of users belongs to.

Authors’ formalization: Given a graph and a set of nodes, find a densely connected subgraph containing the set of users given in input.
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Other applications: Tag suggestions, biological data.
Tag suggestion in Flickr

**Tags**
- Dolomites
- Lake

**Sugg.**
- Mountains
- Nature
- Landscape
Tag suggestions

- Graph of tags: tags $t_i$ and $t_j$ connected if they co-occur in many photos.
- Given a new photo (or any resource) and initial set of tags, recommend new tags to add.
- Tags well connected with one another and the initial set of tags — good candidates.
Protein interactions
Protein interactions
Given: Protein-protein interaction network.
A set of proteins that regulate a gene that a biologist wishes to study.
what other proteins should she study?
  those contained in a compact dense subgraph containing the original proteins.
Related Work
Large body of work on finding communities in social networks:

- Agarwal and Kempe (European Physics Journal, 2008)
- S. White and P. Smyth. (SDM, 2005)
- Y. Dourisboure et al. (WWW, 2007)
- D. Gibson, R. Kumar, and A. Tomkins (VLDB, 2005)

This paper: Query-dependent variant of the problem.

Other related work:

- Lappas et al. (KDD, 2009): team formation.
- FOCS, ICALP, APPROX
Problem Definition
Abstract problem definition

- Input: Undirected graph $G = (V,E)$; a query set of nodes $Q \subseteq V$ and a “goodness” function $f$ that says how good an answer is.
- Find a connected subgraph $H = (V_H, E_H)$ s.t.:
  - $Q \subseteq V_H$ and
  - $f(H)$ is the maximum possible among all connected subgraphs $H$ containing $Q$.

What are some good choices for $f$?
Want $f$ to capture density.
Some choices of density measure

\[ n \equiv \# \text{nodes}; \quad m \equiv \# \text{edges}. \text{Only undirected graphs in this paper.} \]

**Good properties:** small distance, large density, good connectedness.

Two definitions of density of a graph

- \( d(G) = \# \text{of edges in } G / \text{max } \# \text{possible} \)
  
  Formally,
  \[ m / \left[ n(n - 1)/2 \right] \]

- \( D(G) = \# \text{of edges in } G / \# \text{of vertices in } G \)
  
  Formally
  \[ m \quad \text{\_<— average degree}/2. \]
Claim 1: Computing a subgraph $H$ with maximum density $d(H)$ is NP-hard.

Proof Sketch: By reduction from Max Clique.
Fact 2: Computing a subgraph H with maximum density \( D(H) \) can be done in polynomial time but avg. degree based f can lead to counterintuitive results.

\[
D_Q(H) := \max_{v \in V_H} \left( \sum_{q \in Q} d^2(v, q) \right) \leq \Delta
\]

Free riders problem.

=> choose \textit{minimum} degree instead.

Do any problems persist?

Additionally impose a bound on max. distance of nodes in H to query nodes.
Final problem definition

- Input: An undirected graph $G = (V, E)$; query nodes $Q \subseteq V$; distance bound $\Delta$.
- Find a connected subgraph $H = (V_H, E_H)$ s.t.:
  - $Q \subseteq V_H$;
  - $D_Q(H) \leq \Delta$;
  - and $f(H) := \text{min. degree of } H, \text{ is maximized.}$

Good news: The optimal solution can be found in poly time!
The algorithms
A greedy algorithm

1. Let $G_0 = G$. 
2. At each step $t$ if there is a node $v$ in $G_{t-1}$ violating the distance constraint, then remove $v$ and all its edges; 
3. otherwise remove the node with minimum degree in $G_{t-1}$. 
4. Let $G_t$ the graph so obtained, upon saturation. 
5. Among all the graphs $G_0, G_1, \ldots, G_T$ constructed during the execution of the algorithm return the graph $G_i$
   - containing the query nodes; 
   - satisfying the distance constraint; 
   - with maximum minimum degree.

- No need to iterate once $Q$ is no longer contained or connected.
A greedy algorithm

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   - containing the query nodes;
   - satisfying the distance constraint;
   - with maximum minimum degree.

**Theorem**: The greedy algorithm computes an optimum solution for the community-search problem.
Optimality of Greedy (w/o distance constraint)

- Let $G = G_0, G_1, \ldots, G_T$ be the series of graphs obtained from $G$ by removing the min. deg. node and its incident edges, until that min. deg. node is in $Q$ or its removal disconnects $Q$.

- Let $G^*$ be an optimal solution.

- Let $t$ be the smallest number for which the min. deg. node $v$ in $G_t$, is in $G^*$.

- \[ G^* \subseteq G_t' \subseteq G_t, \text{ where } G_t' \text{ is a connected component of } G_t. \]

- $\deg_{G^*}(v) \leq \deg_{G_t'}(v)$.

- $v$ is the min. deg. node in $G_t$ and hence of $G_t'$, so $G_t'$ is an optimal solution! QED

- w/o distance constraint, can be implemented in $O(n+m)$ time (see paper).
Optimality — general case

- Paper claims same logic holds for any monotone constraints.
- However, there are some issues to be resolved there.
- Here is the essence of monotonicity: $G=(V,E)$ and $H=(V',E')$, an induced subgraph. $f$ maps graphs to reals is monotone if for every graph $G$ and induced subgraph $H$,  
  $$f(H) \leq f(G).$$
- Or $f$ could be monotone non-decreasing instead:  
  $$f(H) \geq f(G).$$
- When $f$ is boolean, you get a property (or constraint) instead.

Examples:
- $D_Q(.) \leq \Delta$, i.e, the max. aggregate distance of any node to the query nodes is bounded, is a monotone constraint.
- If $G$ satisfies it, so will any induced subgraph containing $Q$.
- The distance bound constraint remains monotone if distances to query nodes aggregated using max instead.
Optimality in the general case

- \( f(G) = 1 \) iff \( G \) contains \( Q \) and is connected, is monotone. If \( G \) fails, so will any induced subgraph.

- Unfort., bound on min. degree (Ex. 2 in paper) is **not** monotone.

- Requiring nodes of a graph to cover a given set of skills (a la Team Formation paper) is monotone.

- See paper for similar def. of node-monotone, a finer grained notion of monotonicity.

- **General Cocktail Party Problem:** Given query nodes \( Q \) and graph \( G \), you want to find a connected subgraph \( H \) containing \( Q \) that maximizes \( f(\cdot) \), among all such subgraphs which satisfy given monotone properties: say \( \Pi_1, \ldots, \Pi_k \).
  - paper claims an obvious generalization of greedy for this setting is optimal.
The size of the community shouldn’t be too large:

- If we are to organize a party we might not have place for 1M people.
- Humans should be able to analyze the result.

**Bad news**: Adding an upper bound on the number of nodes makes the problem NP-hard even w/o a distance constraint (reduction from Steiner Tree) but...

**Theorem**: Let $H$ and $H'$ be two graphs obtained by executing the greedy algorithm with distance constraint $\Delta$ and $\Delta'$, respectively (the other input parameters are the same).

Then, $\Delta' \leq \Delta$ implies $|V(H')| \leq |V(H)|$. 
Intuition: Bound the size of the graph by making the distance constraint tighter.

GreedyDist:
- solve the problem w/o the cardinality constraint on #nodes.
- if size <= bound, report;
- else successively try with tighter distance constraints (can use binary search!).
  - report any small (i.e., size <= bound) connected subgraph containing Q, if found.
  - else report smallest connected subgraph found that contains Q.
Intuition: Nodes that are far away from the query nodes are most probably not related to them.

GreedyFast:

- Let $k$ be an upperbound on the number of vertices and let $\Delta$ be a distance constraint (i.e., bound).
- Preprocessing: consider only the $k'$ closest nodes to the query nodes, where $k'$ is the smallest number that ensures the resulting graph is connected and contains $k$ nodes.
- Run Greedy with the subgraph induced by these query nodes, as input
Evaluation
Evaluation

Algorithms evaluated on three different datasets:
- DBLP (226k nodes and 1.4M edges);
- Flickr tag graph (38k nodes and 1.3M edges);
- Bio data (16K nodes and 491k nodes).

Queries are generated randomly.

We vary
- Number of query nodes;
- Distance between query nodes;
- Upper bound on the number of nodes.

We measure
- Minimum degree and average degree;
- Size of the output graph;
- Running time.
We consider an approach where at each step we add one node (in contrast with all previous approaches).

A pseudocode:

1. Connect the query nodes: by means of a Steiner Tree algo. (we use a 2-approximation algorithm for this problem);
2. Let $G_t$ be the graph at step $t$;
3. Add the node $v$ with maximum degree in $G_t \cup v$;
   1. Break ties using distance to $Q$ and further ties arbitrarily.
4. Among all the graph $G_0, \ldots, G_T$ constructed, return the one with maximum minimum degree.
Minimum degree vs Size (Flickr)
Average deg. vs. Size (Flickr)
Running time vs Size (Flickr)
Generalization to monotone functions
Generalized Community-Search Problem

Input:
- An undirected graph $G=(V,E)$;
- A set $Q$ of query nodes;
- Integer parameters $k,t$;
- A set of skills $T_v$ associated to every node $v$;
- A required set of skills $\overline{T}$.

Goal: Find an induced subgraph $H$ of $G$ s.t.
- $G$ is connected and contains $Q$;
- The number of vertices of $H$ is $\geq t$;
- The set of skills of $H$ contains $\overline{T}$ $\left( \bigcup_{v \in H} T_v \supseteq \overline{T} \right)$;
- Any node is at distance at most $k$ from the query nodes;
- The minimum degree is maximized.
Generalized Community-Search Problem

Input:
- An undirected graph $G=(V,E)$;
- A set $Q$ of query nodes;
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- Any node is at distance at most $k$ from the query nodes;
- The minimum degree is maximized.

The last one is not monotone but poses no problem. Skill containment — how do you incorporate that in a node elimination paradigm?
Monotone function: \( f(H) \leq f(G) \), if \( H \) is a subgraph of \( G \).

Theorem: There is an **optimum greedy** algorithm for the problem when all constraints are monotone functions.

**Running time**: Depends on the time to evaluate the function \( f_1, \ldots, f_k \), formally \( O(m + \sum n \cdot T_i) \) where \( T_i \) is the time to evaluate the monotone function \( f_i \).
Conclusions
Conclusions and Future Work

Contributions:

- Proposed a novel combinatorial approach for finding the community of a given set of users in input.
- Distance constraints proved to be effective in limiting the size of the output graph.
- Defined a class of functions that can be optimized efficiently.

Questions:

- Are there other useful monotone functions?
- Can we find all communities of a given set of users?
- Community search via Map-Reduce?
- What about other dense subgraphs such as k-core, quasi-clique, k-plex, containing given query nodes?