

# Reasoning Under Uncertainty: Variable Elimination

CPSC 322 – Uncertainty 7

Textbook §10.5

# Lecture Overview

- 1 Recap
- 2 Variable Elimination
- 3 Variable Elimination Example

# Factors and Assigning Variables

- A **factor** is a representation of a function from a tuple of random variables into a number.
- A factor denotes a distribution over the given tuple of variables in some (unspecified) context
- We can make new factors out of an existing factor
- Our first operation: **assign some or all of the variables** of a factor.
  - $f(X_1 = v_1, X_2, \dots, X_j)$ , where  $v_1 \in \text{dom}(X_1)$ , is a factor on  $X_2, \dots, X_j$ .
  - $f(X_1 = v_1, X_2 = v_2, \dots, X_j = v_j)$  is a number that is the value of  $f$  when each  $X_i$  has value  $v_i$ .
- The former is also written as
$$f(X_1, X_2, \dots, X_j)_{X_1 = v_1, \dots, X_j = v_j}$$

# Summing out variables

Our second operation: we can **sum out** a variable, say  $X_1$  with domain  $\{v_1, \dots, v_k\}$ , from factor  $f(X_1, \dots, X_j)$ , resulting in a factor on  $X_2, \dots, X_j$  defined by:

$$\begin{aligned} & \left( \sum_{X_1} f \right) (X_2, \dots, X_j) \\ &= f(X_1 = v_1, \dots, X_j) + \dots + f(X_1 = v_k, \dots, X_j) \end{aligned}$$

# Multiplying factors

- Our third operation: factors can be multiplied together.
- The **product** of factor  $f_1(\overline{X}, \overline{Y})$  and  $f_2(\overline{Y}, \overline{Z})$ , where  $\overline{Y}$  are the variables in common, is the factor  $(f_1 \times f_2)(\overline{X}, \overline{Y}, \overline{Z})$  defined by:

$$(f_1 \times f_2)(\overline{X}, \overline{Y}, \overline{Z}) = f_1(\overline{X}, \overline{Y})f_2(\overline{Y}, \overline{Z}).$$

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# Probability of a conjunction

- Suppose the variables of the belief network are  $X_1, \dots, X_n$ .
- What we **want to compute**: the factor  $P(X_q, X_{o_1} = v_1, \dots, X_{o_j} = v_j)$
- We can compute  $P(X_q, X_{o_1} = v_1, \dots, X_{o_j} = v_j)$  by summing out the variables  $X_{s_1}, \dots, X_{s_k} = \{X_1, \dots, X_n\} \setminus \{X_q, X_{o_1}, \dots, X_{o_j}\}$ .
- We sum out these variables one at a time
  - the order in which we do this is called our **elimination ordering**.

$$\begin{aligned} &P(X_q, X_{o_1} = v_1, \dots, X_{o_j} = v_j) \\ &= \sum_{X_{s_k}} \cdots \sum_{X_{s_1}} P(X_1, \dots, X_n)_{X_{o_1} = v_1, \dots, X_{o_j} = v_j} \end{aligned}$$

# Probability of a conjunction

- What we **know**: the factors  $P(X_i|pX_i)$ .
- Using the chain rule and the definition of a belief network, we can write  $P(X_1, \dots, X_n)$  as  $\prod_{i=1}^n P(X_i|pX_i)$ . Thus:

$$\begin{aligned}
 & P(X_q, X_{o_1} = v_1, \dots, X_{o_j} = v_j) \\
 &= \sum_{X_{s_k}} \cdots \sum_{X_{s_1}} P(X_1, \dots, X_n)_{X_{o_1} = v_1, \dots, X_{o_j} = v_j} \cdot \\
 &= \sum_{X_{s_k}} \cdots \sum_{X_{s_1}} \prod_{i=1}^n P(X_i|pX_i)_{X_{o_1} = v_1, \dots, X_{o_j} = v_j} \cdot
 \end{aligned}$$



# Computing sums of products

Computation in belief networks thus reduces to computing the sums of products.

- It takes 14 multiplications or additions to evaluate the expression  $ab + ac + ad + aeh + afh + agh$ . How can this expression be evaluated more efficiently?

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- How can we compute  $\sum_{X_{s_1}} \prod_{i=1}^n P(X_i|pX_i)$  efficiently?
- Factor out those terms that don't involve  $X_{s_1}$ :

$$\left( \prod_{i|X_{s_1} \notin \{X_i\} \cup pX_i} P(X_i|pX_i) \right) \left( \sum_{X_{s_1}} \prod_{i|X_{s_1} \in \{X_i\} \cup pX_i} P(X_i|pX_i) \right)$$

*(terms that do not involve  $X_{s_1}$ )*
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# Summing out a variable efficiently

To **sum out a variable**  $X_{s_j}$  from a product  $f_1, \dots, f_k$  of factors:

- Partition the factors into
  - those that don't contain  $X_{s_j}$ , say  $f_1, \dots, f_i$ ,
  - those that contain  $X_{s_j}$ , say  $f_{i+1}, \dots, f_k$

We know:

$$\sum_{X_{s_j}} f_1 \times \dots \times f_k = (f_1 \times \dots \times f_i) \left( \sum_{X_{s_j}} f_{i+1} \times \dots \times f_k \right).$$

- $\left( \sum_{X_{s_j}} f_{i+1} \times \dots \times f_k \right)$  is a new factor; let's call it  $f'$ .
- Now we have  $\sum_{X_{s_j}} f_1 \times \dots \times f_k = f_1 \times \dots \times f_i \times f'$ .
- **Store  $f'$  explicitly, and discard  $f_{i+1}, \dots, f_k$ .**
  - Now we've summed out  $X_{s_j}$ .

# Variable elimination algorithm

To compute  $P(X_q | X_{o_1} = v_1 \wedge \dots \wedge X_{o_j} = v_j)$ :

- **Construct a factor** for each conditional probability.
- Set the **observed variables** to their observed values.
- For each of the other variables  $X_{s_i} \in \{X_{s_1}, \dots, X_{s_k}\}$ , **sum out**  $X_{s_i}$
- **Multiply** the remaining factors.
- **Normalize** by dividing the resulting factor  $f(X_q)$  by  $\sum_{X_q} f(X_q)$ .

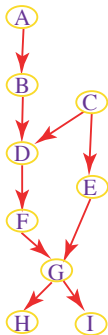
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# Variable elimination example

Compute  $P(G|H = h_1)$ . Elimination order:  $A, C, E, H, I, B, D, F$

- $P(G, H) = \sum_{A, B, C, D, E, F, I} P(A, B, C, D, E, F, G, H, I)$
- $P(G, H) = \sum_{A, B, C, D, E, F, I} P(A) \cdot P(B|A) \cdot P(C) \cdot P(D|B, C) \cdot P(E|C) \cdot P(F|D) \cdot P(G|F, E) \cdot P(H|G) \cdot P(I|G)$

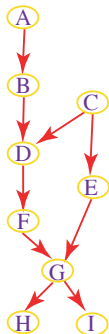




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- $P(G, H) = \sum_{A,B,C,D,E,F,I} P(A) \cdot P(B|A) \cdot P(C) \cdot P(D|B, C) \cdot P(E|C) \cdot P(F|D) \cdot P(G|F, E) \cdot P(H|G) \cdot P(I|G)$
- **Eliminate A:**  $P(G, H) = \sum_{B,C,D,E,F,I} f_1(B) \cdot P(C) \cdot P(D|B, C) \cdot P(E|C) \cdot P(F|D) \cdot P(G|F, E) \cdot P(H|G) \cdot P(I|G)$

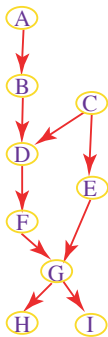


- $f_1(B) := \sum_{a \in \text{dom}(A)} P(A = a) \cdot P(B|A = a)$

# Variable elimination example

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- $P(G, H) = \sum_{B, C, D, E, F, I} f_1(B) \cdot P(C) \cdot P(D|B, C) \cdot P(E|C) \cdot P(F|D) \cdot P(G|F, E) \cdot P(H|G) \cdot P(I|G)$
- **Eliminate  $C$ :**  $P(G, H) = \sum_{B, D, E, F, I} f_1(B) \cdot f_2(B, D, E) \cdot P(F|D) \cdot P(G|F, E) \cdot P(H|G) \cdot P(I|G)$



- $f_1(B) := \sum_{a \in \text{dom}(A)} P(A = a) \cdot P(B|A = a)$
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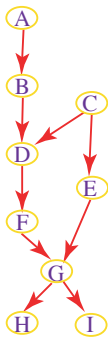
# Variable elimination example

Compute  $P(G|H = h_1)$ . Elimination order:  $A, C, E, H, I, B, D, F$

- $P(G, H) =$   
 $\sum_{B,D,E,F,I} f_1(B) \cdot f_2(B, D, E) \cdot P(F|D) \cdot P(G|F, E) \cdot P(H|G) \cdot P(I|G)$

- **Eliminate  $E$ :**

$$P(G, H) = \sum_{B,D,F,I} f_1(B) \cdot f_3(B, D, F, G) \cdot P(F|D) \cdot P(H|G) \cdot P(I|G)$$



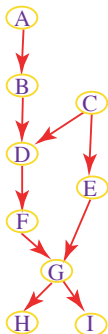
- $f_1(B) := \sum_{a \in \text{dom}(A)} P(A=a) \cdot P(B|A=a)$
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- $f_3(B, D, F, G) := \sum_{e \in \text{dom}(E)} f_2(B, D, E=e) \cdot P(G|F, E=e)$

# Variable elimination example

Compute  $P(G|H = h_1)$ . Elimination order:  $A, C, E, H, I, B, D, F$

- $P(G, H) = \sum_{B, D, F, I} f_1(B) \cdot f_3(B, D, F, G) \cdot P(F|D) \cdot P(H|G) \cdot P(I|G)$
- Observe  $H = h_1$ :

$$P(G, H = h_1) = \sum_{B, D, F, I} f_1(B) \cdot f_3(B, D, F, G) \cdot P(F|D) \cdot f_4(G) \cdot P(I|G)$$



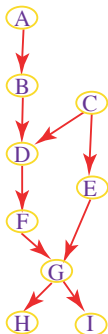
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# Variable elimination example

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- **Eliminate  $I$ :**  

$$P(G, H = h_1) = \sum_{B,D,F} f_1(B) \cdot f_3(B, D, F, G) \cdot P(F|D) \cdot f_4(G) \cdot f_5(G)$$



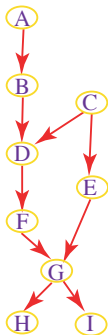
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# Variable elimination example

Compute  $P(G|H = h_1)$ . Elimination order:  $A, C, E, H, I, B, D, F$

- $P(G, H = h_1) = \sum_{B,D,F} f_1(B) \cdot f_3(B, D, F, G) \cdot P(F|D) \cdot f_4(G) \cdot f_5(G)$
- **Eliminate  $B$ :**

$$P(G, H = h_1) = \sum_{D,F} f_6(D, F, G) \cdot P(F|D) \cdot f_4(G) \cdot f_5(G)$$

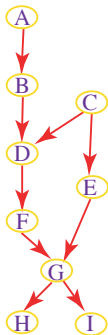


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- $f_6(D, F, G) := \sum_{b \in \text{dom}(B)} f_1(B = b) \cdot f_3(B = b, D, F, G)$

# Variable elimination example

Compute  $P(G|H = h_1)$ . Elimination order:  $A, C, E, H, I, B, D, F$

- $P(G, H = h_1) = \sum_{D, F} f_6(D, F, G) \cdot P(F|D) \cdot f_4(G) \cdot f_5(G)$
- **Eliminate  $D$ :**  $P(G, H = h_1) = \sum_F f_7(F, G) \cdot f_4(G) \cdot f_5(G)$

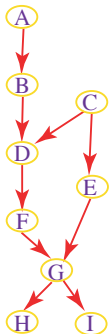


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- $f_6(D, F, G) := \sum_{b \in \text{dom}(B)} f_1(B = b) \cdot f_3(B = b, D, F, G)$
- $f_7(F, G) := \sum_{d \in \text{dom}(D)} f_6(D = d, F, G) \cdot P(F|D = d)$

# Variable elimination example

Compute  $P(G|H = h_1)$ . Elimination order:  $A, C, E, H, I, B, D, F$

- $P(G, H = h_1) = \sum_F f_7(F, G) \cdot f_4(G) \cdot f_5(G)$
- **Eliminate  $F$** :  $P(G, H = h_1) = f_8(G) \cdot f_4(G) \cdot f_5(G)$



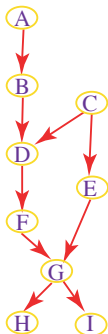
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# Variable elimination example

Compute  $P(G|H = h_1)$ . Elimination order:  $A, C, E, H, I, B, D, F$

- $P(G, H = h_1) = f_8(G) \cdot f_4(G) \cdot f_5(G)$
- **Normalize:**  $P(G|H = h_1) = \frac{P(G, H = h_1)}{\sum_{g \in \text{dom}(G)} P(G, H = h_1)}$



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# What good was Conditional Independence?

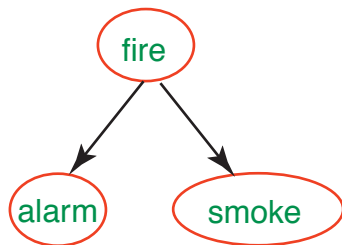
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- And... why did we bother learning **conditional independence**? Does it help us at all?
  - yes—we use the **chain rule decomposition** right at the beginning
- Can we use our knowledge of conditional independence to make this calculation even simpler?
  - yes—there are some variables that we don't have to sum out
  - intuitively, they're the ones that are “pre-summed-out” in our tables
  - example: summing out  $I$  on the previous slide

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One last trick to simplify calculations: we can repeatedly eliminate all **leaf nodes that are neither observed nor queried**, until we reach a fixed point.

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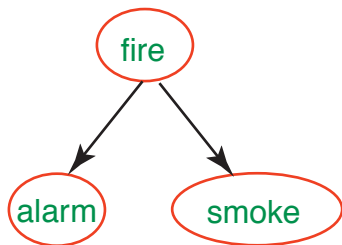


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Can we justify that on a three-node graph—Fire, Alarm, and Smoke—when we ask for:

- $P(\text{Fire})?$
- $P(\text{Fire} \mid \text{Alarm})?$