### Propositional Logic Intro, Syntax

#### CPSC 322 - Logic 1

#### Textbook §5.0 – 5.2

Propositional Logic Intro, Syntax

CPSC 322 - Logic 1, Slide 1

<ロ> (四) (四) (三) (三) (三)

## Lecture Overview



### 2 Logic Intro

#### 3 Propositional Definite Clause Logic: Syntax

#### Propositional Definite Clause Logic: Semantics

Propositional Logic Intro, Syntax

CPSC 322 - Logic 1, Slide 2

3

→ 同 → → 目 → → 目 →

 Recap
 Logic Intro
 PDC: Syntax
 PDC: Semantics

 Planning as a CSP
 Planning as a CSP
 Planning as a CSP
 Planning as a CSP

- We don't have to worry about searching forwards if we set up a planning problem as a CSP
- To do this, we need to "unroll" the plan for a fixed number of steps
  - this is called the horizon
- To do this with a horizon of k:
  - construct a variable for each feature at each time step from 0 to k
  - construct a boolean variable for each action at each time step from 0 to k-1.

個 と く ヨ と く ヨ と …

# CSP Planning: Constraints

As usual, we have to express the preconditions and effects of actions:

- precondition constraints
  - $\bullet\,$  hold between state variables at time t and action variables at time t
  - specify when actions may be taken
- effect constraints
  - between state variables at time t, action variables at time t and state variables at time t+1
  - explain how state variables at time t+1 are affected by the action taken at time t
  - this includes both causal and frame axioms
    - basically, it goes back to the feature-centric representation the book discusses before STRIPS
    - of course, solving the problem this way doesn't mean we can't *encode* the problem using STRIPS

|▲□ ▶ ▲ 目 ▶ ▲ 目 → ○ ○ ○

# CSP Planning: Constraints

Other constraints we must/may have:

- $\bullet\,$  initial state constraints constrain the state variables at time 0
- $\bullet$  goal constraints constrain the state variables at time k
- action constraints
  - specify which actions cannot occur simultaneously
    - note that without these constraints, there's nothing to stop the planner from deciding to take several actions simultaneously
    - when the order between several actions doesn't matter, this is a good thing
  - these are sometimes called mutual exclusion (mutex) constraints
- state constraints
  - hold between variables at the same time step
  - they can capture physical constraints of the system
  - they can encode maintenance goals

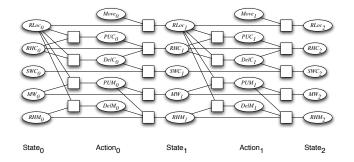
3

回 と く ヨ と く ヨ と …

Recap

**PDC: Semantics** 

## CSP Planning: Robot Example



The constraints shown represent the preconditions of actions and the effects of actions.

◆□ > ◆□ > ◆三 > ◆三 > 三 の < ⊙

### Lecture Overview





#### 3 Propositional Definite Clause Logic: Syntax

#### Propositional Definite Clause Logic: Semantics

Propositional Logic Intro, Syntax

CPSC 322 - Logic 1, Slide 7

3

- 4 回 2 - 4 □ 2 - 4 □

## Logic: A more general framework for reasoning

- Let's now think about how to represent a world about which we have only partial (but certain) information
- Our tool: propositional logic
- General problem:
  - tell the computer how the world works
  - tell the computer some facts about the world
  - ask a yes/no question about whether other facts must be true

A D > A D > ...

Recap Logic Intro PDC: Syntax PDC: Semantics

Why Propositions?

We'll be looking at problems that could still be represented using CSPs. Why use propositional logic?

- Specifying logical formulae is often more natural than constructing arbitrary constraints
- It is easier to check and debug formulae than constraints
- We can exploit the Boolean nature for efficient reasoning
- We need a language for asking queries that may be more complicated than asking for the value of one variable
- It is easy to incrementally add formulae
- Logic can be extended to infinitely many variables (using logical quantification)
- This is a starting point for more complex logics (e.g., first-order logic) that do go beyond CSPs.

個人 くほん くほん 一足

### Representation and Reasoning System

### Definition (RSS)

A Representation and Reasoning System (RRS) is made up of:

- syntax: specifies the symbols used, and how they can be combined to form legal sentences
- semantics: specifies the meaning of the symbols
- reasoning theory or proof procedure: a (possibly nondeterministic) specification of how an answer can be produced.

Using an RRS

- Begin with a task domain.
- Oistinguish those things you want to talk about (the ontology).
- Ochoose symbols in the computer to denote propositions
- It is a system knowledge about the domain.
- Ask the system whether new statements about the domain are true or false.

## Propositional Definite Clauses

- Propositional Definite Clauses: our first representation and reasoning system.
- Two kinds of statements:
  - that a proposition is true
  - that a proposition is true if one or more other propositions are true
- To define this RSS, we'll need to specify:
  - syntax
  - semantics
  - proof procedure

2

## Lecture Overview





#### Oppositional Definite Clause Logic: Syntax

#### 4 Propositional Definite Clause Logic: Semantics

Propositional Logic Intro, Syntax

3

- 4 同 ト 4 臣 ト 4 臣 ト

	Recap	Logic Intro		PDC: Syntax	PDC: Semantics
F	Propositional I	Definite	Clauses:	Syntax	

#### Definition (atom)

An atom is a symbol starting with a lower case letter

Propositional Logic Intro, Syntax

- < ≣ →

Image: A = A = A

Recap	Logic Intro		PDC: Syntax	<b>PDC: Semantics</b>
Propositional	Definite	Clauses:	Syntax	

#### Definition (atom)

An atom is a symbol starting with a lower case letter

#### Definition (body)

A body is an atom or is of the form  $b_1 \wedge b_2$  where  $b_1$  and  $b_2$  are bodies.

▲□ ▶ ▲ □ ▶ ▲ □ ▶ …

Recap			

## Propositional Definite Clauses: Syntax

Definition (atom)

An atom is a symbol starting with a lower case letter

Definition (body)

A body is an atom or is of the form  $b_1 \wedge b_2$  where  $b_1$  and  $b_2$  are bodies.

Definition (definite clause)

A definite clause is an atom or is a rule of the form  $h \leftarrow b$  where h is an atom and b is a body. (Read this as "h if b.")

▲□ ▶ ▲ □ ▶ ▲ □ ▶ …

Recan		

## Propositional Definite Clauses: Syntax

Definition (atom)

An atom is a symbol starting with a lower case letter

Definition (body)

A body is an atom or is of the form  $b_1 \wedge b_2$  where  $b_1$  and  $b_2$  are bodies.

### Definition (definite clause)

A definite clause is an atom or is a rule of the form  $h \leftarrow b$  where h is an atom and b is a body. (Read this as "h if b.")

#### Definition (knowledge base)

A knowledge base is a set of definite clauses

▶ < ≞ ▶

The following are syntactically correct statements in our language:

 $\bullet \ ai\_is\_fun$ 

æ

- 4 回 2 4 三 2 4 三 2 4

The following are syntactically correct statements in our language:

- $ai_is_fun$
- $ai\_is\_fun \leftarrow get\_good\_grade$

æ

・回 ・ ・ ヨ ・ ・ ヨ ・

 Recap
 Logic Intro
 PDC: Syntax
 PDC: Semantics

 Syntax:
 Example

The following are syntactically correct statements in our language:

- $ai_is_fun$
- $ai\_is\_fun \leftarrow get\_good\_grade$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work$

2

프 🖌 🛪 프 🛌

< 🗗 > <

The following are syntactically correct statements in our language:

- $ai_is_fun$
- $ai\_is\_fun \leftarrow get\_good\_grade$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work \land remain\_awake$

2

물 🕨 🗶 물 🕨 👘

The following are syntactically correct statements in our language:

- $\bullet ~ai\_is\_fun$
- $ai\_is\_fun \leftarrow get\_good\_grade$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work \land remain\_awake$

The following statements are syntactically incorrect:

•  $ai\_is\_fun \lor ai\_is\_boring$ 

The following are syntactically correct statements in our language:

- $\bullet ~ai\_is\_fun$
- $ai\_is\_fun \leftarrow get\_good\_grade$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work \land remain\_awake$

The following statements are syntactically incorrect:

- $ai\_is\_fun \lor ai\_is\_boring$
- $ai\_is\_fun \land relaxing\_term \leftarrow get\_good\_grade \land not\_too\_much\_work$

The following are syntactically correct statements in our language:

- $\bullet ~ai\_is\_fun$
- $ai\_is\_fun \leftarrow get\_good\_grade$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work$
- $ai\_is\_fun \leftarrow get\_good\_grade \land not\_too\_much\_work \land remain\_awake$

The following statements are syntactically incorrect:

- $ai\_is\_fun \lor ai\_is\_boring$
- $ai\_is\_fun \land relaxing\_term \leftarrow get\_good\_grade \land not\_too\_much\_work$

Do any of these statements *mean* anything? Syntax doesn't answer this question.

## Lecture Overview





#### 3 Propositional Definite Clause Logic: Syntax

#### Propositional Definite Clause Logic: Semantics

Propositional Logic Intro, Syntax

3

- - 4 回 ト - 4 回 ト

Recap

### Propositional Definite Clauses: Semantics

Semantics allows you to relate the symbols in the logic to the domain you're trying to model.

Definition (interpretation)

An interpretation I assigns a truth value to each atom.

Propositional Logic Intro, Syntax

토 🕨 🗶 토 🕨 👘

Recap

## Propositional Definite Clauses: Semantics

Semantics allows you to relate the symbols in the logic to the domain you're trying to model.

#### Definition (interpretation)

An interpretation I assigns a truth value to each atom.

We can use the interpretation to determine the truth value of clauses and knowledge bases:

#### Definition (truth values of statements)

- A body  $b_1 \wedge b_2$  is true in I if and only if  $b_1$  is true in I and  $b_2$  is true in I.
- A rule h ← b is false in I if and only if b is true in I and h is false in I.
- A knowledge base *KB* is true in *I* if and only if every clause in *KB* is true in *I*.

## Models and Logical Consequence

### Definition (model)

A model of a set of clauses is an interpretation in which all the clauses are *true*.

Propositional Logic Intro, Syntax

個 と く ヨ と く ヨ と …

## Models and Logical Consequence

### Definition (model)

A model of a set of clauses is an interpretation in which all the clauses are *true*.

#### Definition (logical consequence)

If KB is a set of clauses and g is a conjunction of atoms, g is a logical consequence of KB, written  $KB \models g$ , if g is *true* in every model of KB.

- we also say that g logically follows from KB, or that KB entails g.
- In other words,  $KB \models g$  if there is no interpretation in which KB is *true* and g is *false*.

< □ > < □ > < □ > < □ > < □ > < Ξ > = Ξ

Recan			

Logic Intro

PDC: Syntax

PDC: Semantics

### Example: Models

$$KB = \begin{cases} p \leftarrow q, \\ q, \\ r \leftarrow s. \end{cases}$$

	p	q	r	s
$I_1$	true	true	true	true
$I_2$	false	false	false	false
$I_3$	true	true	false	false
$I_4$	true	true	true	false
$I_5$	true	true	false	true

Which interpretations are models?

イロト イヨト イヨト イヨト

æ

**PDC: Semantics** 

### Example: Models

÷.

$$KB = \begin{cases} p \leftarrow q, \\ q, \\ r \leftarrow s. \end{cases}$$

	p	q	r	s
$I_1$	true	true	true	true
$I_2$	false	false	false	false
$I_3$	true	true	false	false
$I_4$	true	true	true	false
$I_5$	true	true	false	true

is a model of KBnot a model of KBis a model of KBis a model of KBnot a model of KB

æ

イロン イヨン イヨン イヨン

### Example: Models

i.

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	p	q	r	s
$I_1$	true	true	true	true
$I_2$	false	false	false	false
$I_3$	true	true	false	false
$I_4$	true	true	true	false
$I_5$	true	true	false	true

is a model of KBnot a model of KBis a model of KBis a model of KBnot a model of KB

Which of the following is true?

• 
$$KB \models q$$
,  $KB \models p$ ,  $KB \models s$ ,  $KB \models r$ 

æ

- - 4 回 ト - 4 回 ト

## Example: Models

i.

$$KB = \begin{cases} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{cases}$$

	p	q	r	s
$I_1$	true	true	true	true
$I_2$	false	false	false	false
$I_3$	true	true	false	false
$I_4$	true	true	true	false
$I_5$	true	true	false	true

is a model of KBnot a model of KBis a model of KBis a model of KBnot a model of KB

Which of the following is true?

• 
$$KB \models q$$
,  $KB \models p$ ,  $KB \models s$ ,  $KB \models r$ 

• 
$$KB \models q$$
,  $KB \models p$ ,  $KB \not\models s$ ,  $KB \not\models r$ 

æ

- - 4 回 ト - 4 回 ト