Local Search

CPSC 322 Lecture 12

February 2, 2007 Textbook §3.8



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Local Search

Lecture Overview



2 Local Search





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Arc Consistency

Definition

An arc $\langle X, r(X, \bar{Y}) \rangle$ is arc consistent if for each value of X in \mathbf{D}_X there is some value \bar{Y} in $\mathbf{D}_{\bar{Y}}$ such that $r(X, \bar{Y})$ is satisfied.

- In symbols, $\forall X \in \mathbf{D}_X, \ \exists \bar{Y} \in \mathbf{D}_{\bar{Y}}$ such that $r(X, \bar{Y})$ is satisfied.
- A network is arc consistent if all its arcs are arc consistent.
- If an arc $\langle X, \bar{Y} \rangle$ is *not* arc consistent, all values of X in \mathbf{D}_X for which there is no corresponding value in $\mathbf{D}_{\bar{Y}}$ may be deleted from \mathbf{D}_X to make the arc $\langle X, \bar{Y} \rangle$ consistent.
 - This removal can never rule out any models

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Arc Consistency Algorithm

- Consider the arcs in turn making each arc consistent.
 - An arc $\left\langle X,r(X,\bar{Y})\right\rangle$ needs to be revisited if the domain of X is reduced.
- Regardless of the order in which arcs are considered, we will terminate with the same result: an arc consistent network.

Revisiting Edges

• When we change the domain of a variable X in the course of making an arc $\langle X,r\rangle$ arc consistent, we add every arc $\langle Z,r'\rangle$ where r' involves X and:

•
$$r \neq r'$$

• $Z \neq X$

- Thus we don't add back the same arc:
 - This makes sense—it's definitely arc consistent.

Revisiting Edges

- When we change the domain of a variable X in the course of making an arc $\langle X, r \rangle$ arc consistent, we add every arc $\langle Z, r' \rangle$ where r' involves X and:
 - $r \neq r'$ • $Z \neq X$

- We don't add back other arcs involving the same variable \boldsymbol{X}
 - We've just *reduced* the domain of X
 - If an arc $\langle X,r\rangle$ was arc consistent before, it will still be arc consistent
 - in the "for all" we'll just check fewer values

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Revisiting Edges

- When we change the domain of a variable X in the course of making an arc $\langle X, r \rangle$ arc consistent, we add every arc $\langle Z, r' \rangle$ where r' involves X and:
 - $r \neq r'$ • $Z \neq X$
- We don't add back other arcs involving the same constraint and a different variable:
 - Imagine that such an arc—involving variable Y—had been arc consistent before, but was no longer arc consistent after X's domain was reduced.
 - This means that some value in $Y{\rm 's}$ domain could satisfy r only when X took one of the dropped values
 - But we dropped these values precisely because there were no values of Y that allowed r to be satisfied when X takes these values—contradiction!

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Arc Consistency Outcomes

- Three possible outcomes (when all arcs are arc consistent):
 - One domain is empty \Rightarrow no solution
 - Each domain has a single value \Rightarrow unique solution
 - $\bullet\,$ Some domains have more than one value $\Rightarrow\,$ may or may not be a solution
 - in this case, arc consistency isn't enough to solve the problem: we need to perform search

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Local Search

Local Search

- Many search spaces are too big for systematic search.
- A useful method in practice for some consistency and optimization problems is local search
 - idea: consider the space of complete assignments of values to variables
 - neighbours of a current node are similar variable assignments
 - move from one node to another according to a function that scores how good each assignment is

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Local Search

Definition

A local search problem consists of a:

- Set of Variables. A node in the search space will be a complete assignment to all of the variables.
- Neighbour relation. An edge in the search space will exist when the neighbour relation holds between a pair of nodes.
- Scoring function. This can be used to incorporate information about how many constraints are violated. It can also incorporate information about the cost of the solution in an optimization context.

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Selecting Neighbours

How do we choose the neighbour relation?

- Usually this is simple: some small incremental change to the variable assignment
 - assignments that differ in one variable's value
 - assignments that differ in one variable's value, by a value difference of one
 - assignments that differ in two variables' values, etc.
- There's a trade-off: bigger neighbourhoods allow more nodes to be compared before a step is taken
 - the best step is more likely to be taken
 - each step takes more time: in the same amount of time, multiple steps in a smaller neighbourhood could have been taken
- Usually we prefer pretty small neighbourhoods

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Lecture Overview



2 Local Search





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Local Search

Hill Climbing

Hill climbing means selecting the neighbour which best improves the scoring function.

• For example, if the goal is to find the highest point on a surface, the scoring function might be the height at the current point.

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Gradient Ascent

What can we do if the variable(s) are continuous?

- With a constant step size we could overshoot the maximum.
- Here we can use the scoring function h to determine the neighbourhood dynamically:
 - Gradient ascent: change each variable proportional to the gradient of the heuristic function in that direction.
 - The value of variable X_i goes from v_i to $v_i + \eta \frac{\partial h}{\partial X_i}$.
 - η is the constant of proportionality that determines how big steps will be
 - Gradient descent: go downhill; v_i becomes $v_i \eta \frac{\partial h}{\partial X_i}$.
 - these partial derivatives may be estimated using finite differences

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Problems with Hill Climbing

Foothills local maxima that are not global maxima

Plateaus heuristic values are uninformative

Ridge foothill where a larger neighbour relation would help

Ignorance of the peak no way of detecting a global maximum

