Let's Make Block Coordinate Descent Go Fast!

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EUROPT Workshop on Advances in Continuous Optimization Montreal, Canada July 12th, 2017 • Block coordinate descent methods are key tools in large-scale optimization.

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- \rightarrow Cheap iteration costs.
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- Used for almost two decades to solve LASSO and SVMs.
- → **Any** improvements on convergence will affect many applications.

- We propose 4 ways to speed up Block Coordinate Descent (BCD) methods:
 - 1. New greedy block selection rules.
 - 2. New second-order update rule.
 - 3. New exact update rule for LASSO and SVMs.
 - 4. New exact update rule for graph-structured problems.

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 \rightarrow E.g., gradient descent update $d^k = -\alpha_k \nabla_{b_k} f(x^k)$ for some $\alpha_k > 0$.

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- Assume that f is L_b -block-wise Lipschitz continuous,

$$\|\nabla_b f(x+U_b d) - \nabla_b f(x)\| \le L_b \|d\|, \text{ for all } d.$$

 \rightarrow If f is twice-differentiable, this is equivalent to $\nabla^2_{bb} f(x) \preceq L_b \mathbb{I}$ for each block b.

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- \rightarrow Incorporates Lipschitz information in the rule.
- \rightarrow Equivalent to MI for quadratics.

 As an obvious extension of the GSL rule to the block setting, we propose the Block Gauss-Southwell-Lipschitz (BGSL) rule:

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- Derived by minimizing quadratic bound from block-wise Lipschitz continuity.
- Guarantees more progress than the block GS rule.
- \rightarrow Unlike GSL, not equivalent to the MI rule for quadratic functions.

Experiment: L2-Regularized Logistic Regression

• Comparing block selection rules using fixed blocks.



$$\|\nabla_b f(x+U_b d) - \nabla_b f(x)\|_{H_b^{-1}} \le \|d\|_{H_b} = \sqrt{d^T H_b d},$$

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- \rightarrow May be difficult to find Hessian bounds H_b , depends on how we define blocks.

Blocking Strategy

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• Fixed blocks we could use Lipschitz constants to help determine the partition.

Experiment: L2-Regularized Logistic Regression

• Comparing block partitioning strategies using BGSD rule.



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Do better updates exist? Yes!

• Why do we expect to develop better updates than the Hessian-bound update?

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 \rightarrow We consider a Newton-style method based on a cubic regularization framework.

• While gradient-style methods are based on a quadratic upper-bound,

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 - → Guaranteed to decrease the objective without needing extra objective function evaluations required for a line search.

Experiment: Multi-class Logistic Regression

• Comparing update rules using variable blocks with greedy block selection.



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 - E.g., 2-variable non-separable quadratic
- \rightarrow Possible to get superlinear convergence for problems with certain structures.

• Consider minimizing a differentiable function *f* with L1-regularization,

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• E.g., LASSO: $F(x) = \frac{1}{2!} ||Ax - b||^2 + \lambda ||x||_1$ for $\lambda > 0$.

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→ SUPERLINEAR CONVERGENCE!

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\rightarrow FINITE TERMINATION!

Experiment: Dual SVM

• Comparing update methods using variable blocks with greedy selection.



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 \rightarrow Exploit connection to Gaussian Markov random fields, update tree-structured blocks in O(M) using Gaussian belief propagation.
Message-Passing for Sparse Quadratics

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- For lattice-structured graphs, can use blocks of size n/2 in O(n).
- Maintains modelling dependencies.



Experiment: Sparse Quadratic Problem

• Comparing exact updates using variable blocks with greedy selection.



- Exact solver uses M = 8, Gaussian belief propagation method uses $M = 8^3$.
- NP-hard to choose best "tree-structure" block.
 - → Use approximation method that performs substantially better than BGSD.

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 - Greedy block updates have "active set" identification property for LASSO, SMVs.
 - Superlinear convergence with variable blocks and higher order updates.
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 - Superlinear convergence with variable blocks and higher order updates.
 - Finite convergence with variable blocks and exact updates.
- Propose optimal block update strategy for sparse quadratic problems.
 - Use "tree-structured" blocks.
 - Exploits Gaussian belief propagation algorithm developed for GMRFs.
 - Requires linear time in block size.