# HOMEWORK 6 SOLUTIONS, MATH 340, FALL 2015 

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(1) The dictionary consists of

$$
\begin{array}{rll}
\vec{x}_{B} & =A_{B}^{-1} \vec{b} & -A_{B}^{-1} A_{N} \vec{x}_{N} \\
z & =\vec{c}_{B}^{\mathrm{T}} A_{B}^{-1} \vec{b} & +\left(\vec{c}_{N}^{\mathrm{T}}-\vec{c}_{B}^{\mathrm{T}} A_{B}^{-1} A_{N}\right) \vec{x}_{N}
\end{array}
$$

Since there are always $m$ basic variables, and $A_{B}$ is $m \times m$, naively inverting $A_{B}$ will take order $m^{3}$ FLOPs. Multiplying $\vec{c}_{B}^{\mathrm{T}} A_{B}^{-1}$ will take order $m^{2}$ FLOPs. Also the cost of computing $A_{B}^{-1} A_{N} \vec{x}$ is mostly not needed: we only compute $A_{B}^{-1}$ times the $A_{N}$ row corresponding to the entering variable, which again costs order $m^{2}$ FLOPs. The only operation that can be expensive is multiplying $\vec{c}_{B}^{T} A_{B}^{-1}$ by $A_{N}$ to compute the $z$ row.

If $A_{N}$ is sparse, then although $\vec{c}_{B}^{\top} A_{B}^{-1}$ is $m$ dimensional and $A_{N}$ is $m \times n$, the fact that $A_{N}$ is mostly zeros means that will only need a number of multiplications and additions each equal to the number of nonzero entries of $A_{N}$. Hence this step will require $m n$ times $5 \%$; an ordinary pivot (at least when the dictionary is dense) requires $m n$ additions and $m n$ multiplications. Hence the revised simplex method requires order $m^{3}$ plus $2 m n / 20$ FLOPs (one addition and one multiplication for each of $m n$ times $5 \%$ entries of $A_{N}$ ), as opposed to $2 m n$ for a standard simplex method pivot. So you save roughly a factor of 20 for each pivot (neglecting the order $m^{3}$ terms, since $m n=m^{5}$ is much larger).
(2) One again, the standard simplex method requires roughly $2 m n$ operations to update each dictionary. And once again, the revised simplex method takes order $m^{3}$ operations plus the cost of taking the product of $\vec{c}_{B}^{\top} A_{B}^{-1}$ with $A_{N}$. However, rather than multiply all of $A_{N}$ you only have to do the columns of $A_{N}$ corresponding to the non-basic variables that might enter. So if you can identify $n / 10$ variables from which to choose, you only need to compute $\vec{c}_{B}^{\mathrm{T}} A_{B}^{-1}$ times $n / 3$ of the $n$ columns of $A_{N}$. Hence you save a factor of 3 for each pivot where this is possible. In general, if $n / 3$ is replaced with $n^{\prime}$, then you save a factor of $n / n^{\prime}$; so for $n^{\prime}=n / 30$ you save a factor of 30 , and for $n^{\prime}=\sqrt{n}$ you save a factor of $\sqrt{n}$.

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