

Space Bounded Randomized Complexity Classes

one-sided error, log space bounded classes

handy techniques for probabilistic reasoning

RLP

- A language L is in RLP if there is an log-space *and poly-time* PTM M such that
 - if $x \in L$ then $\Pr[M \text{ accepts } x] \geq 2/3$ and
 - if $x \notin L$ then $\Pr[M \text{ accepts } x] = 0$

UPATH is in RLP

- $UPATH = \{ (G,s,t) \mid \text{node } t \text{ can be reached from node } s \text{ in an } \textit{undirected} \text{ graph } G \}$

UPATH is in RLP

UPATH Algorithm:

- On input (G,s,t) , follow a random path from s
 - If t is reached at some step, halt and accept
 - If t is not reached within $6e(n-1)$ steps, halt and reject
- The algorithm is correct if t is not reachable from s , since it must reject
- What if t is reachable from s ?

Analysis of UPATH

- Let $T(G,s,t)$ be the number of edges of G that are traversed on a random walk from s to t
- Claim: Expected value of $T(G,s,t) \leq 2e(n-1)$
- We can use the claim and Markov's inequality to show that the random walk algorithm accepts yes instances of UPATH with probability at least $2/3$

Analysis of UPATH

- *Markov's Inequality*: If X is a nonnegative random variable and k is a positive real then

$$\text{Prob}[X \geq k E[X]] \leq 1/k$$

- Let X be $T(G,s,t)$, i.e., the number of edges of G that are traversed on a random walk from s to t
- Since the expected value of $T(G,s,t) \leq 2e(n-1)$, then the probability that a random walk from s takes at least $6e(n-1)$ steps to reach t is at most $1/3$

Analysis of UPATH

- Let $T(G,s,t)$ be the number of edges of G that are traversed on a random walk from s to t
- We still need to prove the claim that the expected value of $T(G,s,t) \leq 2e(n-1)$
- We need some background on *Markov Chains*

Background on Markov Chains

A finite *Markov chain* with discrete time and stationary transition probabilities is an infinite sequence of random variables over some state space S

$X_0, X_1, X_2, \dots,$

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such that for all i, j in S , $\text{Prob}[X_k = j \mid X_{k-1} = i] = P_{ij}$,
where P_{ij} may depend on i and j but not on k

The matrix P is called the *transition matrix*

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P_{ij}^m is the probability of reaching j from i in exactly m steps (there is an easy proof by induction on m)

Background on Markov Chains

A Markov Chain is *irreducible* if for all states i and j , there exists k such that

$$\text{Prob}[X_k = j \mid X_0 = i] > 0$$

Theorem: Let P be the transition matrix of an irreducible Markov Chain. Then $\boldsymbol{\pi}P = \boldsymbol{\pi}$ has a unique solution $\boldsymbol{\pi}$ up to constant multiplicative factors

If $\sum_i \boldsymbol{\pi}_i = 1$, $\boldsymbol{\pi}$ is sometimes referred to as the *stationary distribution* of the Markov chain

Markov Chains and Random Walks on Graphs

- For an undirected connected graph $G = (V, E)$, let
 - $N(u)$ be the set of neighbours of node u
 - $d(u) = |N(u)|$ be the degree of node u
- A random walk on G is an irreducible Markov chain with state space equal to V and with
 - $P_{uv} = 1/d(u)$, if $\{u, v\}$ is in E
 - $P_{uv} = 0$, if $\{u, v\}$ is not in E
- The stationary distribution $\boldsymbol{\pi}$ of this random walk is such that $\boldsymbol{\pi}_u = d(u)/2e$

Analysis of UPATH: Random Commutes

A random commute from i to j in G is a random walk starting at i that ends the first time it returns to i after having at some point visited j

Let θ_{ijuv} be the expected number of times edge $\{u,v\}$ is visited from u to v on a random commute from i to j

Analysis of UPATH: Random Commutes

Claim: θ_{ijuv} is independent of v : for all v' in $N(u)$,
 $\theta_{ijuv} = \theta_{ijuv'}$

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Proof follows from fact that each time u is visited, it is equally likely that v or v' is visited next

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Proof: The following identity holds for any u in V :

$$d(u) \theta_{iju} = \sum_{v \in N(u)} \theta_{ijv}$$

expected number of times a random commute from i to j leaves u

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So, the vector of terms $d(u) \theta_{iju}$ is a constant times the stationary distribution $\boldsymbol{\pi}$, where the constant is independent of u (but depends on i and j).

Analysis of UPATH: Back to $T(G,s,t)$

For an edge $\{i,j\}$ of G , let T_{ij} be the expected time to reach j from i on a random walk starting at i

Claim: $T_{ij} \leq 2e$.

Proof: $T_{ij} \leq \sum_{\{u,v\} \in E} (\theta_{ijuv} + \theta_{ijvu}) = 2e \theta_{ijij} \leq 2e$

Analysis of UPATH: Back to $T(G,s,t)$

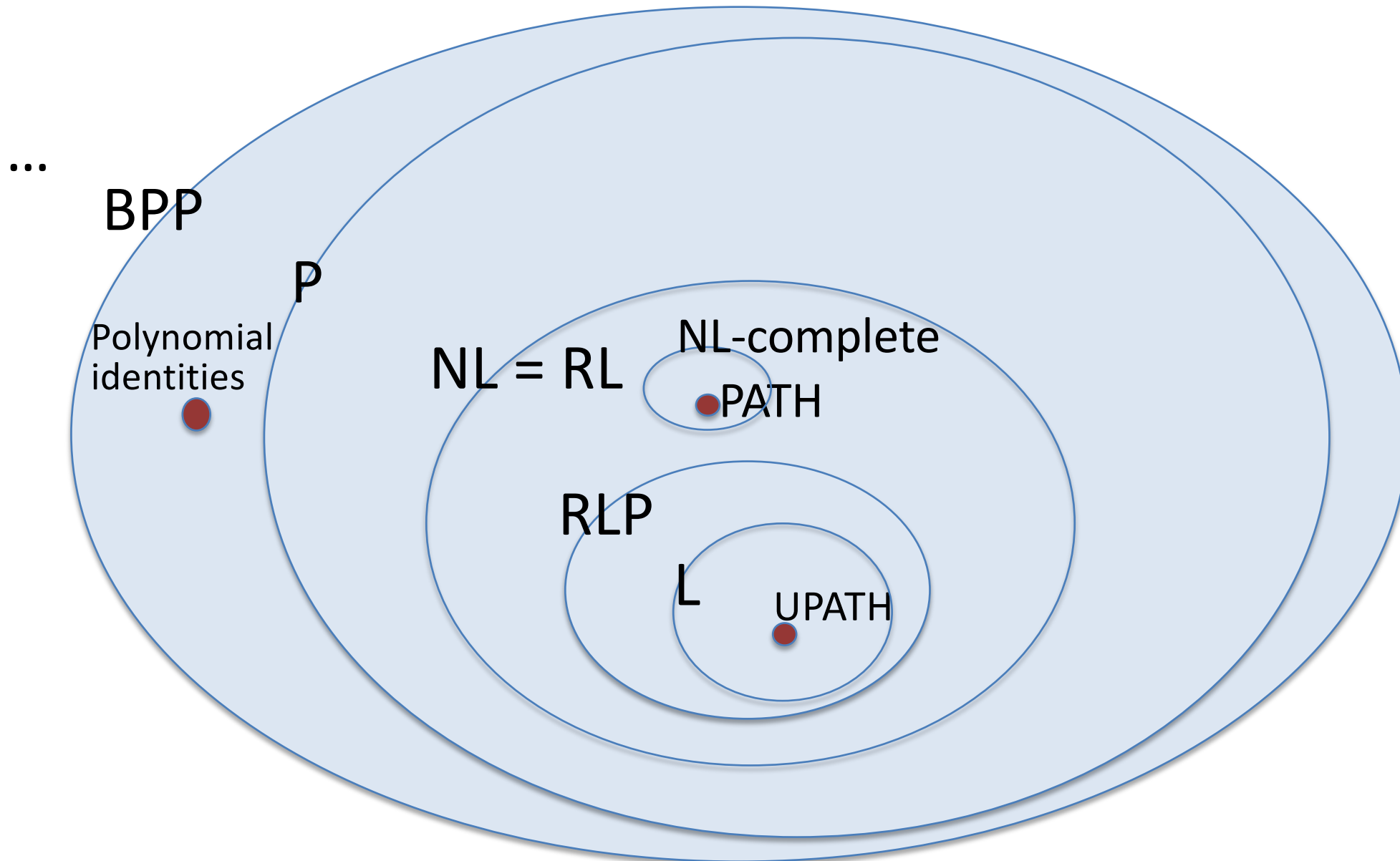
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Finally: if p is a path of length at most $n-1$ from s to t , then $T(G,s,t) \leq \sum_{\{i,j\} \in p} T_{ij} \leq 2e(n-1)$

Summary



In fact, UPATH is in Log Space: Shown by Omer Reingold, 2004.

Summary

- Many conjecture that
 - $BPP = P$
 - $RLP = L$ (here L is “log space”, see Reingold, Trevisan, Vadhan 2004)
- Reingold’s proof uses theory of graph expanders
- There’s an extensive body of work on pseudorandom generators, motivated in part by the goal of resolving these conjectures (see appendix of Arora-Barak)