Intelligent Systems (AI-2)

Computer Science cpsc422, Lecture 11

Oct, 3, 2016



CPSC 422, Lecture 11

422 big picture: Where are we?

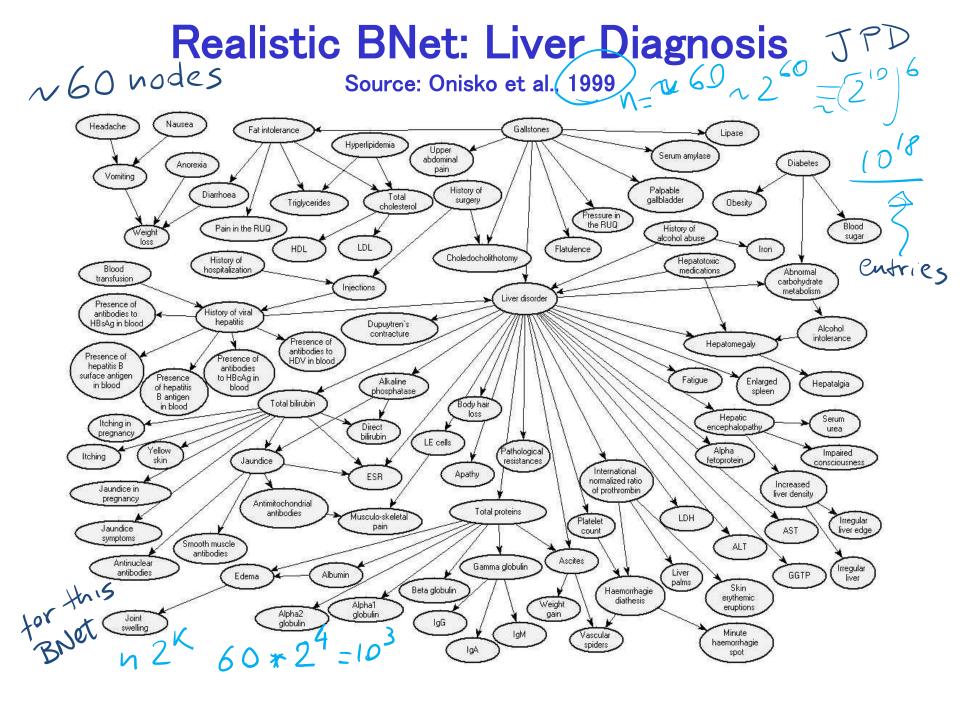
StarAI (statistical relational AI)

Hybrid: Det +Sto Prob CFG Prob Relational Models Markov Logics

	Deterministic	Stochastic Markov I	Logics
Query	Logics First Order Logics Ontologies Temporal rep. • Full Resolution • SAT	Belief Nets Approx. : Gibbs Markov Chains and HMMs Forward, Viterbi···. Approx. : Particle Filtering Undirected Graphical Models Markov Networks Conditional Random Fields Markov Decision Processes and Partially Observable MDP • Value Iteration	
		Approx. Inference <i>Reinforcement Learning</i>	Representation
	Applicatio	ons of AI	Reasoning Technique

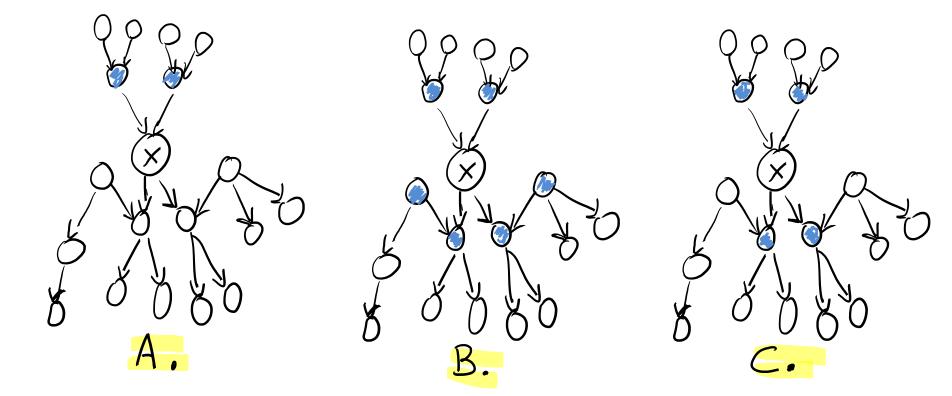
Lecture Overview

- **Recap of BNs** Representation and Exact Inference
- Start Belief Networks Approx. Reasoning
 - Intro to Sampling
 - First Naïve Approx. Method: Forward Sampling
 - Second Method: Rejection Sampling



Revise (in)dependencies

Independence (Markov Blanket)



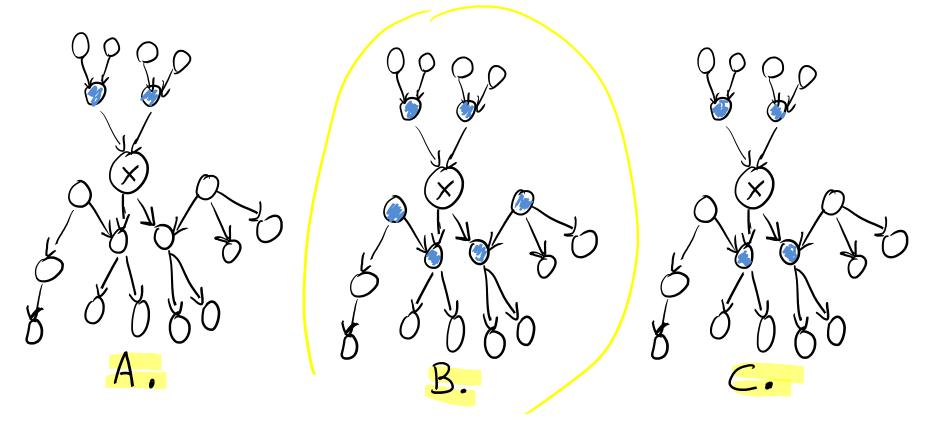
What is the minimal set of nodes that must be observed in order to make **node X** independent from all the non-observed nodes in the network



Slide 8

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Independence (Markov Blanket)



A node is conditionally independent from all the other nodes in the network, given its parents, children, and children's parents (i.e., its **Markov Blanket**) Configuration B

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Variable elimination algorithm: Summary

 $P(Z, Y_1, Y_i, Z_1, Z_i)$

To compute $P(Z|Y_1=v_1, \cdots, Y_j=v_j)$:

- 1. Construct a factor for each conditional probability.
- 2. Set the observed variables to their observed values.
- 3. Given an <u>elimination ordering</u>, <u>simplify/decompose</u> sum of products
 - For all Z_i : Perform products and sum out Z_i
- 4. Multiply the remaining factors (all in ? Z)
- 5. Normalize: divide the resulting factor f(Z) by $\sum_{Z} f(Z)$.

Variable elimination ordering

 $P(G,D=t) = \sum_{A,B,C} f(A,G) f(B,A) f(C,G,A) f(B,C)$ $\sum_{A} f(A,G) \sum_{B} f(B,A) \sum_{C} f(C,G,A) f(B,C)$ CBA BCA $\sum_{A} f(A,G) \sum_{C} f(C,G,A) \sum_{B} f(B,C) f(B,A)$

Complexity: Just Intuition...

- Tree-width of a network given an elimination ordering: max number of variables in a factor created while running VE.
- Tree-width of a belief network : min tree-width over all elimination orderings (only on the graph structure and is a measure of the sparseness of the graph)

- The complexity of VE is exponential in the tree-width 🛞 and linear in the number of variables.
- Also, finding the elimination ordering with minimum tree-width is NP-hard 🟵 (but there are some good elimination ordering heuristics)

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Approximate Inference

Basic idea:

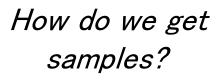
- Draw N samples from known prob. distributions
- Use those samples to estimate unknown prob. distributions

Why sample?

 Inference: getting a sample is faster than computing the right answer (e.g. with variable elimination)

We use *Sampling*

Sampling is a process to obtain samples adequate to estimate an unknown probability



Samples



Known prob. distribution(s)

Estimates for unknown (hard to compute) distribution(s)

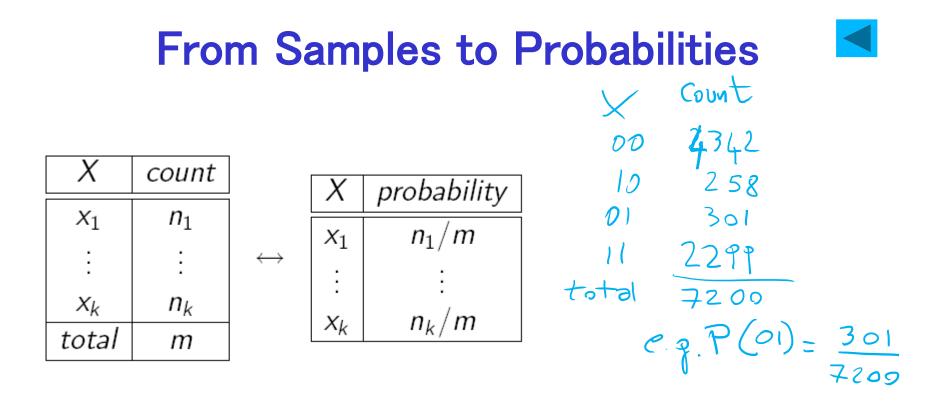
Generating Samples from a Known Distribution

For a random variable X with

- values $\{x_1, \dots, x_k\}$
- Probability distribution $P(X) = \{P(x_1), \dots, P(x_k)\}$
- Partition the interval [0, 1] into k intervals p_i , one for each x_i , with length $P(x_i)$
- To generate one sample
 - ✓ Randomly generate a value y in [0, 1] (i.e. generate a value from a uniform distribution over [0, 1].
 - ✓ Select the value of the sample based on the interval p_i that includes y

From probability theory: $P(y \subset p_i) = Length(p_i) = P(x_i)$

$$29, 5, c^{2}, 5, 6$$



Count total number of samples *m*

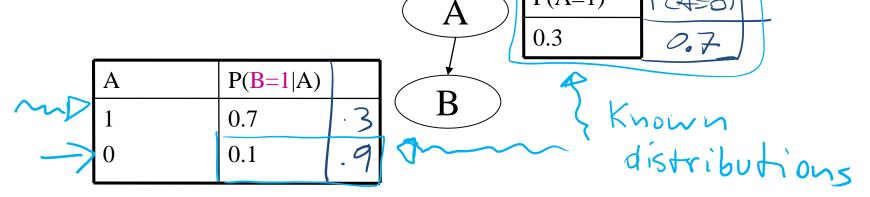
Count the number n_i of samples x_i

Generate the frequency of sample x_i as n_i / m

This frequency is your estimated probability of x_i

Sampling for Bayesian Networks (N)

Suppose we have the following BN with two binary variables P(A=1)P(A=0)



 \succ It corresponds to the joint probability distribution

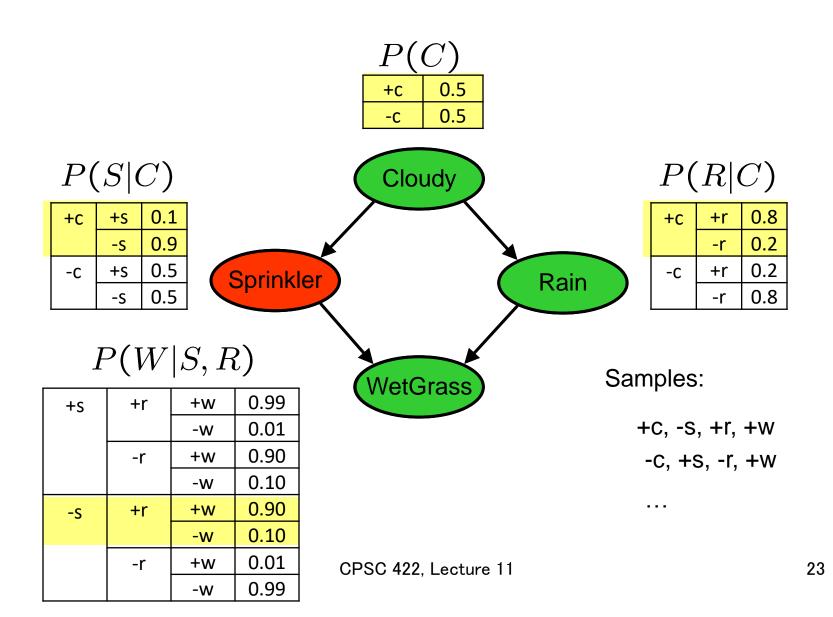
P(A,B) = P(B|A)P(A)

• P(A,B) = P(B|A)F(A)> To sample from P(A,B) i.e., unknown distribution

- we first sample from P(A). Suppose we get A = 0.
- In this case, we then sample from $\mathbb{P}(B | A = 0)$ ٠
- If we had sampled (A = 1) then in the second step we would have sampled • from

B -1

Prior (Forward) Sampling



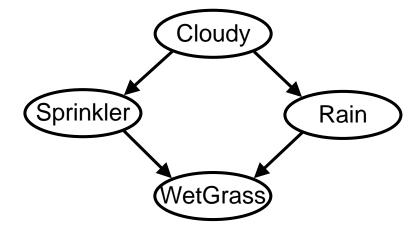
Example

We'll get a bunch of samples from the BN:

+c, -s, +r, +w +c, +s, +r, +w -c, +s, +r, -w +c, -s, +r, +w -c, -s, -r, +w

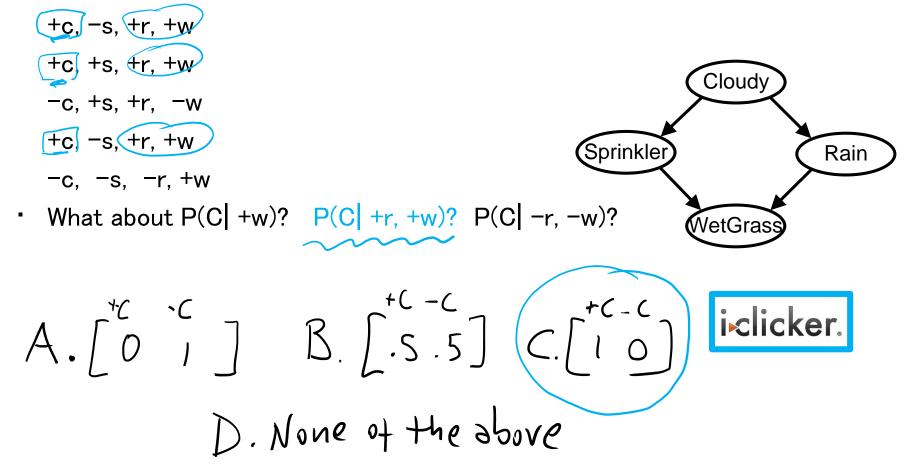
If we want to know P(W)

- We have counts <+w:4, -w:1>
- Normalize to get $P(W) = \langle +W : 2 \rangle$
- This will get closer to the true distribution with more samples



Example

Can estimate anything else from the samples, besides P(W), P(R), etc:



Can use/generate fewer samples when we want to estimate a probability conditioned on evidence?

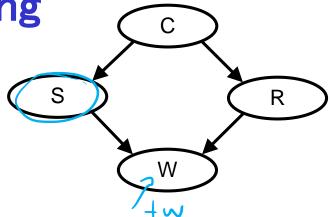
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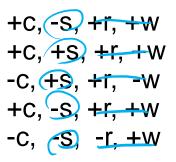
Rejection Sampling

Let's say we want P(S +w)

- Ignore (reject) samples which don't have W=+w
- This is called rejection sampling
- It is also consistent for conditional probabilities (i.e., correct in the limit)

See any problem as the number of evidence vars increases?





Hoeffding's inequality

- Suppose *p* is the true probability and *s* is the sample average from *n* independent samples. $P(|s-p| > \varepsilon) \le 2e^{-2n\varepsilon^2}$
- > p above can be the probability of any event for random variable $X = {X_1, \dots, X_n}$ described by a Bayesian network
- > If you want an infinitely small probability of having an error greater than $\mathcal{E}_{,}$ you need infinitely many samples
- > But if you settle on something less than infinitely small, let's say δ , then you just need to set

$$2e^{-2n\varepsilon^2} < \delta$$

So you pick

- the error \mathcal{E} you can tolerate,
- the frequency δ with which you can tolerate it
- And solve for *n*, i.e., the number of samples that can ensure this performance $1 = \delta$

$$n > \frac{-\ln\frac{\delta}{2}}{2\varepsilon^2} \qquad (1)$$

Hoeffding's inequality

> Examples:

• You can tolerate an error greater than 0.1 only in 5% of your cases

 $n > \frac{-\ln \frac{o}{2}}{2\varepsilon^2}$

con rewrite (

- Set ε =0.1, δ = 0.05
- Equation (1) gives you n > 184

- If you can tolerate the same error (0.1) only in 1% of the cases, then you need 265 samples
- If you want an error greater than 0.01 in no more than 5% of the cases, you need 18,445 samples
 so it should be clear that
 J goes down

Learning Goals for today's class

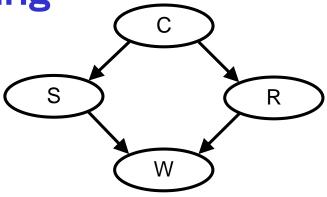
≻You can:

- Motivate the need for approx inference in Bnets
- Describe and compare Sampling from a single random variable
- Describe and Apply Forward Sampling in BN
- Describe and Apply Rejection Sampling
- Apply Hoeffding's inequality to compute number of samples needed

TODO for Wed

- Read textbook 6.4.2
- Assignment-2 will be out today: Start working on it
- Next research paper will be this coming Fri

Rejection Sampling



Let's say we want P(C)

- No point keeping all samples around
- Just tally counts of C as we go

Let's say we want P(C|+s)

- Same thing: tally C outcomes, but ignore (reject) samples which don't have S=+s
- This is called rejection sampling
- It is also consistent for conditional probabilities (i.e., correct in the limit)
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+C, -S, +r, +W +C, +S, +r, +W -C, +S, +r, -W +C, -S, +r, +W -C, -S, -r, +W