

# Finish Search

Computer Science cpsc322, Lecture 10

*(Textbook Chpt 3.6)*

January, 28, 2008



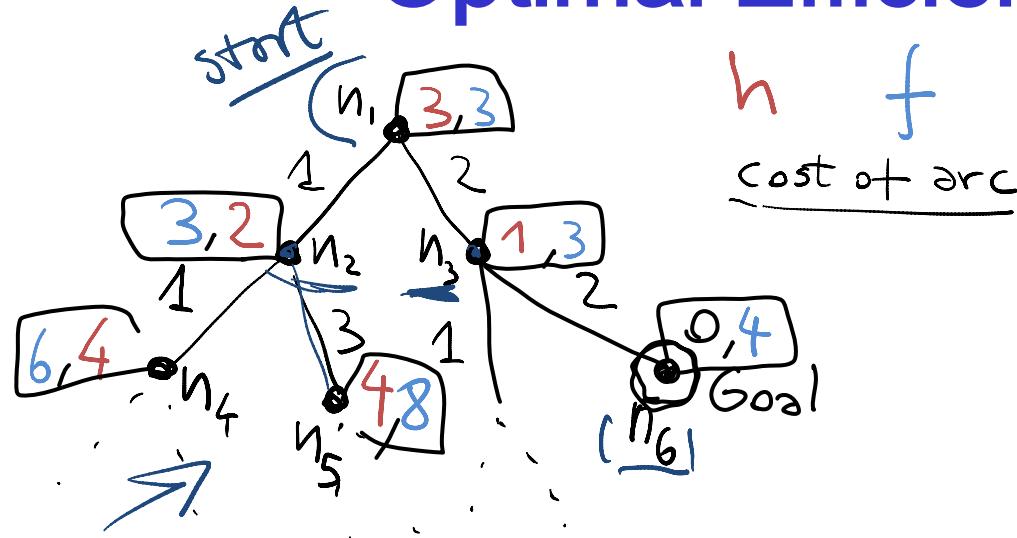
# Announcements

- Another practice exercise has been posted. These exercises can really help you with the assignment. Please do check them out!
- New textbook pdf should be online. With parts you need to read clearly marked
- Branch and Bound on Aispace is buggy ☹

# Lecture Overview

- Optimal Efficiency Example
- Pruning Cycles and Repeated states Examples
- Dynamic Programming
- 8-puzzle Applet
- Search Recap

# Optimal Efficiency Example

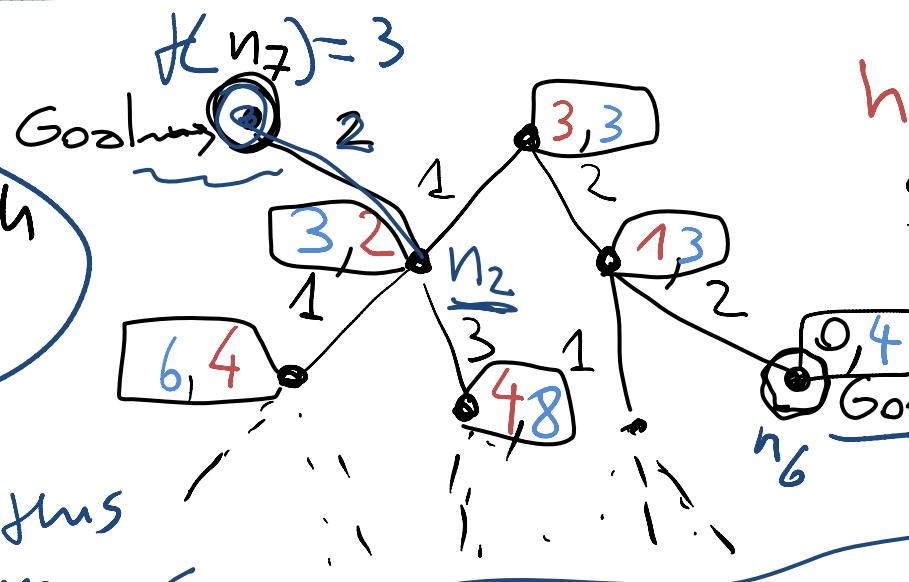


$$t = c + h$$

$A^*$

$$f < f^* = 4$$

$A'$   
 is not expansion  
 $n_2$



New Search Problem

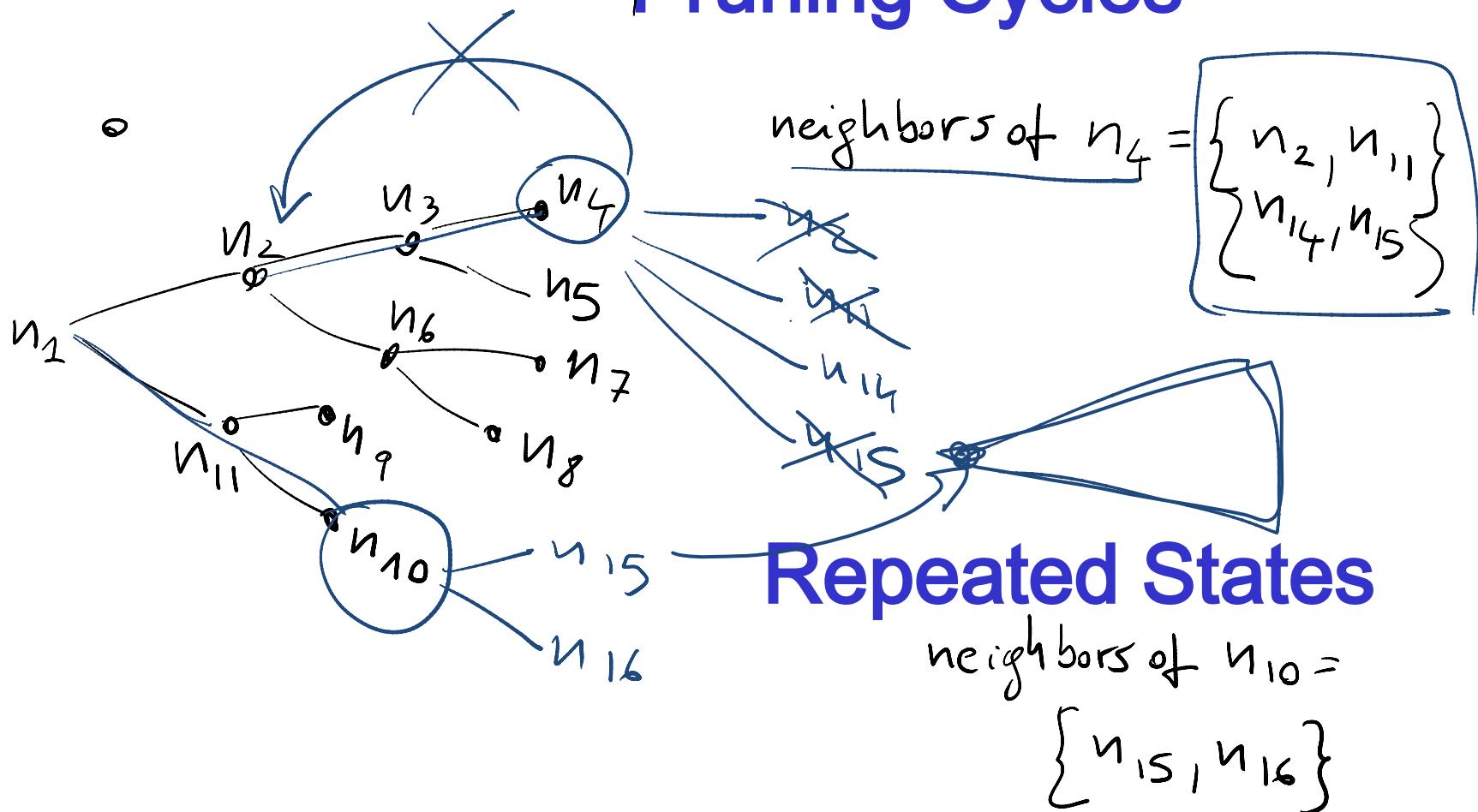
So  $A'$  not optimal on this problem

$A'$   
 solution =  $n_6$   
 $f(\text{solution}) = 4$

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# Pruning Cycles



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# Dynamic Programming

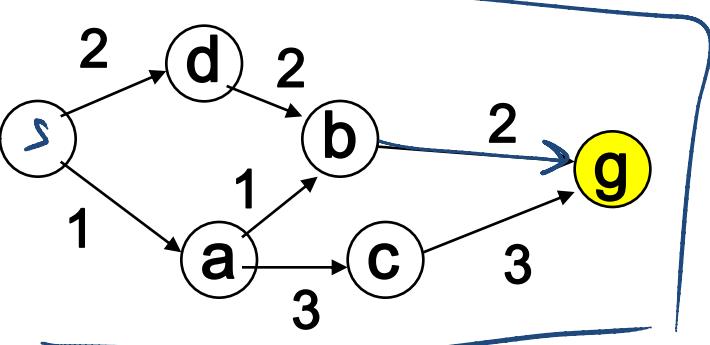
Idea: for statically stored graphs, build a table of  $dist(n)$  the actual distance of the shortest path from node  $n$  to a goal.

This is the perfect heuristic

This can be built backwards from the goal:

$$\underline{dist}(n) = \begin{cases} 0 & \text{if } is\_goal(n), \leftarrow \\ \min_{\langle n,m \rangle \in A} (\cancel{(n,m)} + \underline{dist}(m)) & \text{otherwise} \end{cases}$$

m all neighbors of n



$\cancel{cost}(n, m)$

$dist(n)$

g  
0

b  $\min((2+0)) = 2$

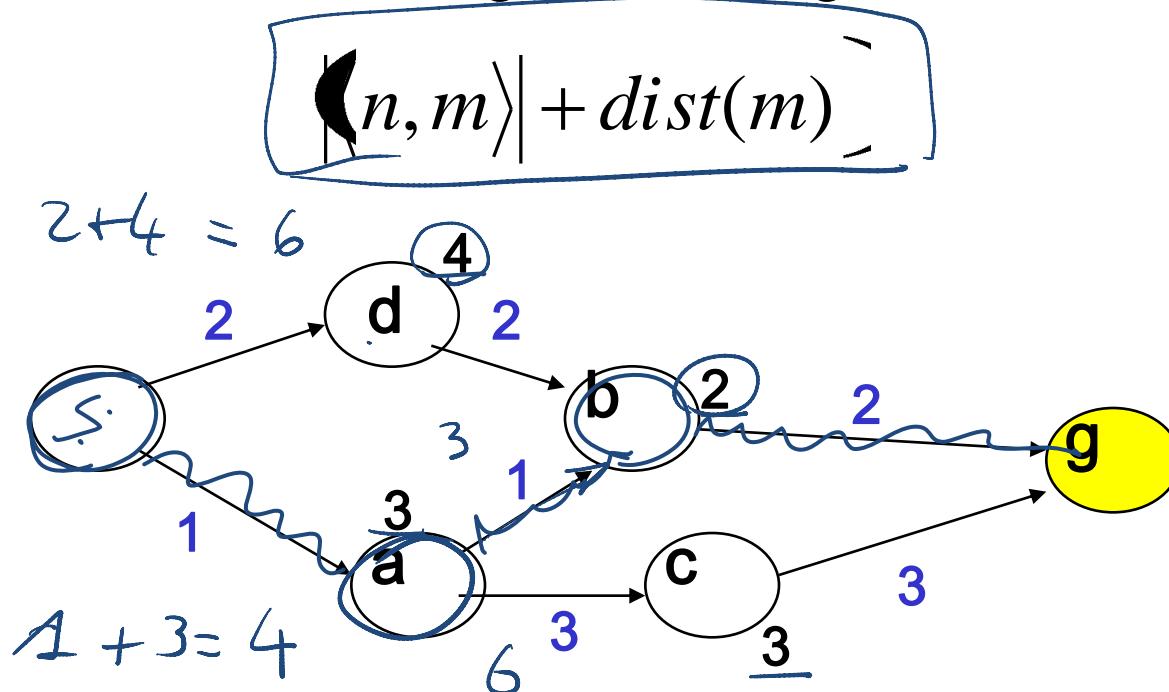
c  $\min((3+0)) = 3$

a  $\min((1+2), (3+3)) = 3$

# Dynamic Programming

This can be used locally to determine what to do.

From each node  $n$  go to its neighbor which minimizes



But there are at least two main problems:

- You need enough space to store the graph. ↴
- The  $dist$  function needs to be recomputed for each goal ↴

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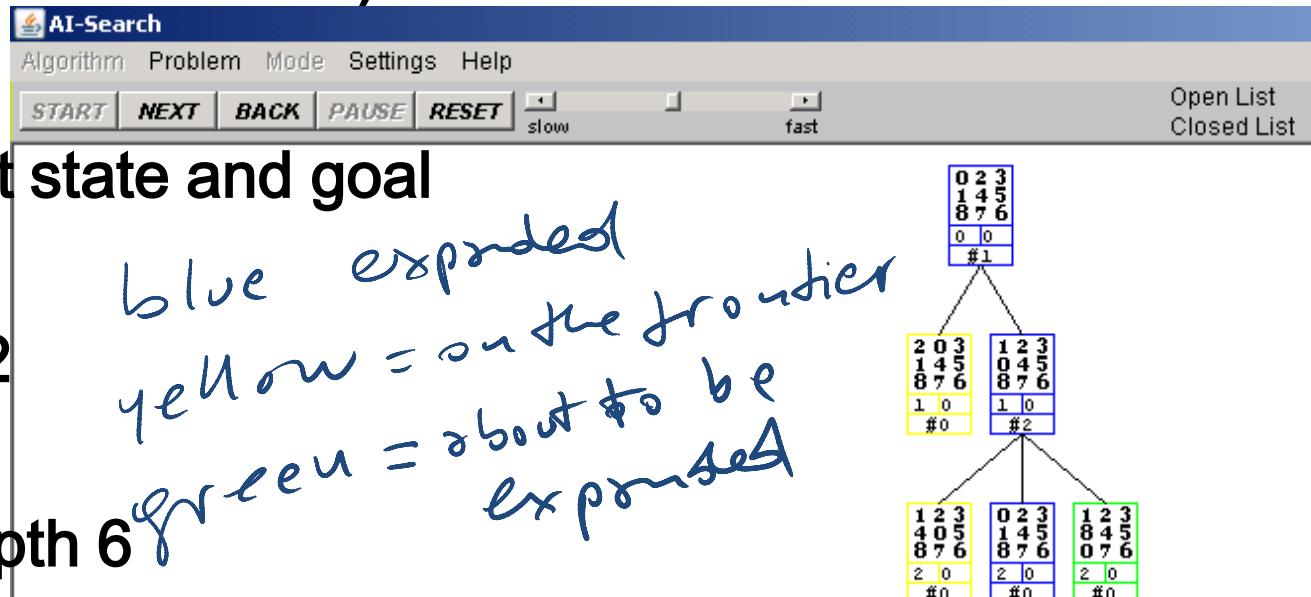
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# DFS, BFS, A\* Animation Example

- The AI-Search animation system

<http://www.cs.rmit.edu.au/AI-Search/Product/>

- To examine Search strategies when they are applied to the 8puzzle
- Compare only DFS, BFS and A\* (with only the two heuristics we saw in class )



- With default start state and goal

- DFS will find

Solution at depth 32

- BFS will find

Optimal solution depth 6

- A\* will also find opt. sol. expanding much less nodes

# nPuzzles are not always solvable

Half of the starting positions for the  $n$ -puzzle are impossible to resolve (for more info on 8puzzle)  
<http://www.isle.org/~sbay/ics171/project/unsolvable>



- So experiment with the AI-Search animation system with the default configurations.
- If you want to try new ones keep in mind that you may pick unsolvable problems

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# Recap Search

	<u>Selection</u>	<u>Complete</u>	<u>Optimal</u>	<u>Time</u>	<u>Space</u>
DFS	LIFO	N	N	$O(b^m)$	$O(mb)$
BFS	FIFO	Y	Y	$O(b^m)$	$O(b^m)$
IDS(C)	LIFO	Y	Y	$O(b^m)$	$O(mb)$
LCFS	min cost	Y	Y	$O(b^m)$	$O(b^m)$
BFS	min h	N	N	$O(b^m)$	$O(b^m)$
A*	min $f = c + h$	Y	Y	$O(b^m)$	$O(b^m)$
B&B	LIFO + pruning	N	Y	$O(b^m)$	$O(mb)$
IDA*	LIFO	Y	Y	$O(b^m)$	$O(mb)$
MBA*	min f	N	N	$O(b^m)$	$O(b^m)$

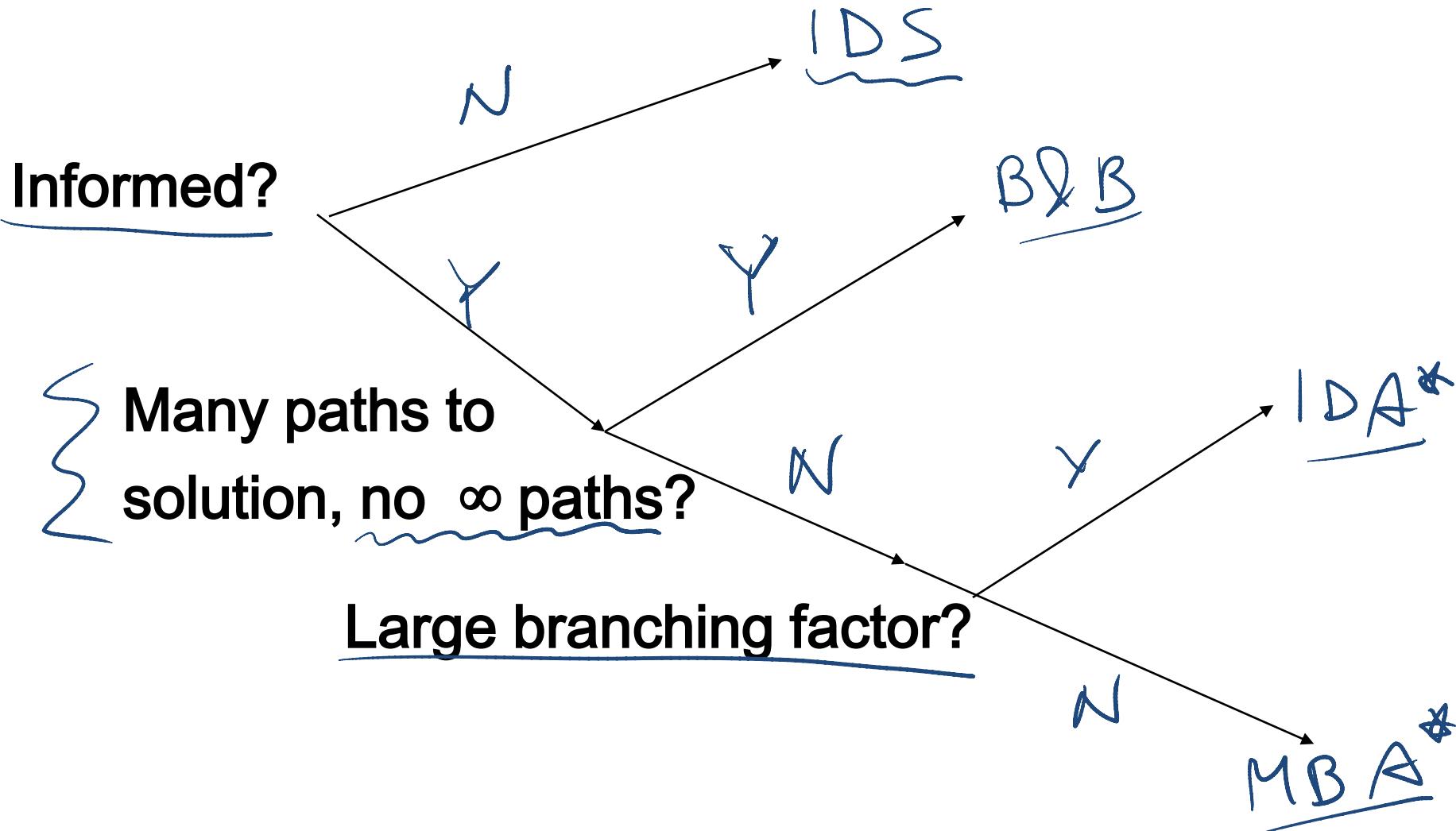
# Recap Search (some qualifications)

	Complete	Optimal	Time	Space
DFS	N	N	$O(b^m)$	$O(mb)$
BFS	Y	Y	$O(b^m)$	$O(b^m)$
IDS(C)	Y	Y	$O(b^m)$	$O(mb)$
LCFS	Y	Y ? <i>&lt; &gt; O</i>	$O(b^m)$	$O(b^m)$
BFS	N	N	$O(b^m)$	$O(b^m)$
A*	Y	Y ? <i>had to use</i> $O(b^m)$	$O(b^m)$	$O(b^m)$
B&B	N	Y ?	$O(b^m)$	$O(mb)$
IDA*	Y	Y	$O(b^m)$	$O(mb)$
MBA*	N	N	$O(b^m)$	$O(b^m)$

# Search in Practice

	Complete	Optimal	Time	Space
DFS	N	N	$O(b^m)$	$O(mb)$
BFS	Y	Y	$O(b^m)$	$O(b^m)$
<b>IDS(C)</b>	Y	Y	$O(b^m)$	$O(mb)$
LCFS	Y	Y	$O(b^m)$	$O(b^m)$
BFS	N	N	$O(b^m)$	$O(b^m)$
A*	Y	Y	$O(b^m)$	$O(b^m)$
<b>B&amp;B</b>	N	Y	$O(b^m)$	$O(mb)$
<b>IDA*</b>	Y	Y	$O(b^m)$	$O(mb)$
<b>MBA*</b>	N	N	$O(b^m)$	$O(b^m)$
BDS	Y	Y	$O(b^{m/2})$	$O(b^{m/2})$

# Search in Practice (cont')



# (Adversarial) Search: Chess

Deep Blue's Results in the second tournament:

- second tournament: won 3 games, lost 2, tied 1
- 30 CPUs + 480 chess processors
- Searched 126.000.000 nodes per sec
- Generated 30 billion positions per move reaching depth 14 routinely



- Iterative Deepening with evaluation function (similar to a heuristic) based on 8000 features (e.g., sum of worth of pieces: pawn 1, rook 5, queen 10)

# Modules we'll cover in this course: R&Rsys

Problem	Environment	
	Deterministic	Stochastic
Static	Arc Consistency Vars + Constraints Search	Belief Nets Var. Elimination
Inference	Logics Search	
Sequential	STRIPS actions Search	Decision Nets Var. Elimination
Planning		Markov Processes Value Iteration
Representation		
Reasoning Technique		

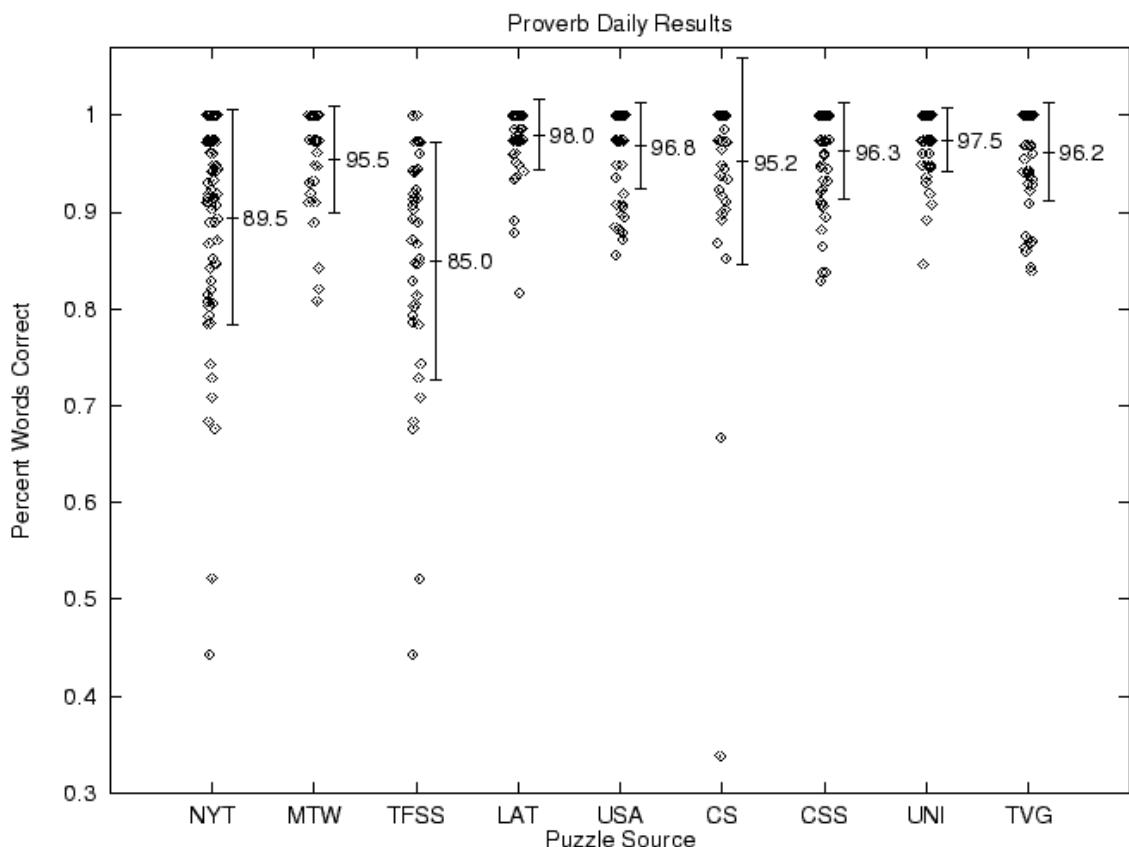
# CSPs: Crossword Puzzles

## Daily Puzzles

370 puzzles from 7 sources.

Summary statistics:

- 95.3% words correct (miss three or four words per puzzle)
- 98.1% letters correct
- 46.2% puzzles completely correct



Source: *Michael Littman*

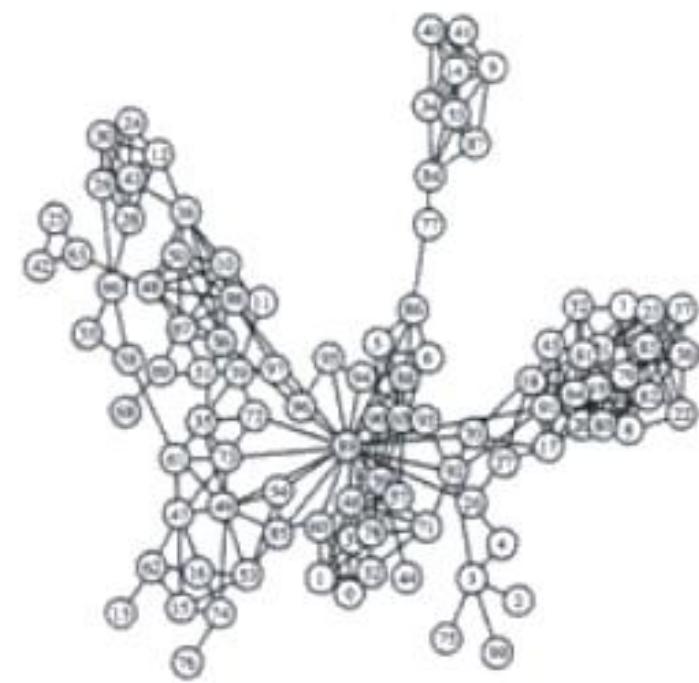
# CSPs: Radio link frequency assignment

Assigning frequencies to a set of radio links defined between pairs of sites in order to avoid interferences.

Constraints on frequency depend on position of the links and on physical environment .

Source: INRIA

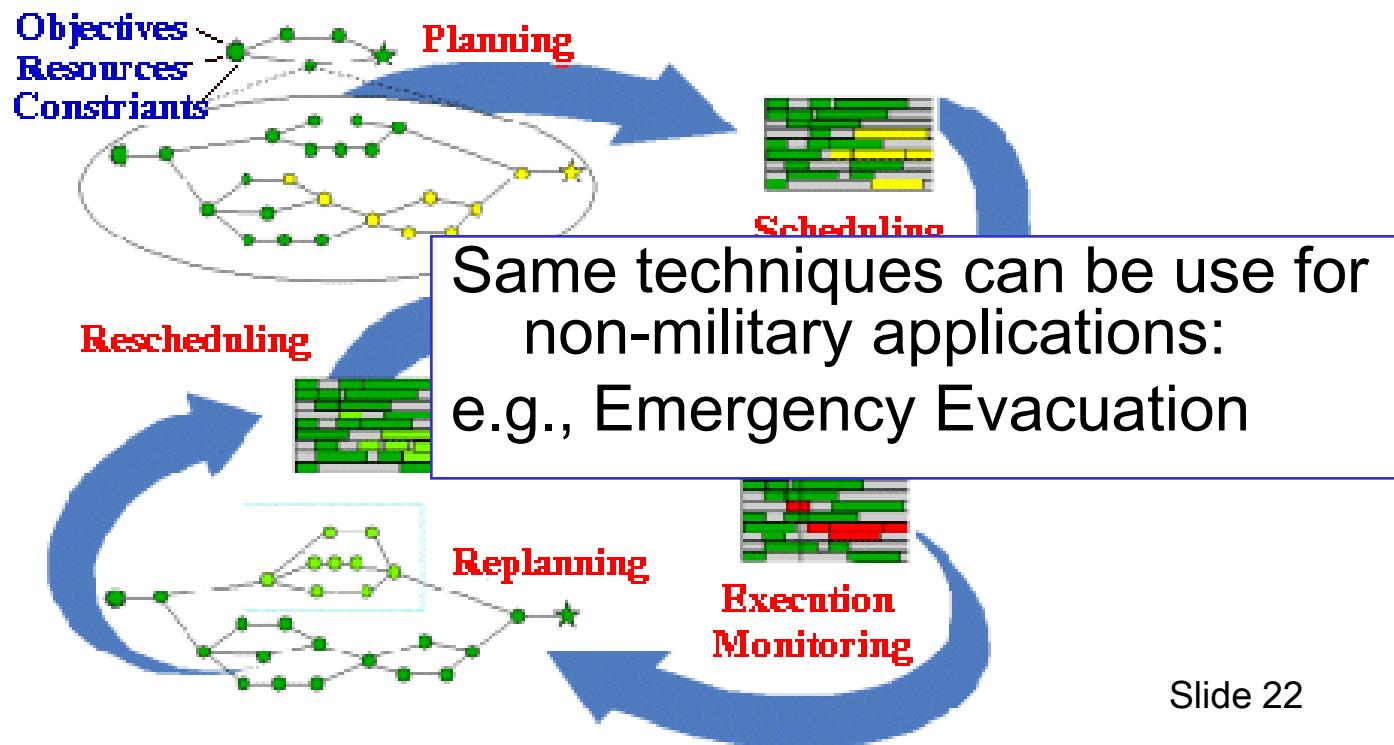
Sample Constraint network



# Planning & Scheduling: Logistics

Dynamic Analysis and Replanning Tool (Cross & Walker)

- logistics planning and scheduling for military transport
- used in the 1991 Gulf War by the US
- problems had 50,000 entities (e.g., vehicles); different starting points and destinations



# Next class

Start Constraint Satisfaction Problems (CSPs)

Textbook 4.1-4.3