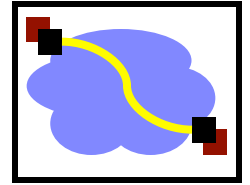


416 Distributed Systems

Errors and Failures, part 2

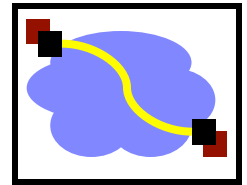
Feb 3, 2016

Options in dealing with failure

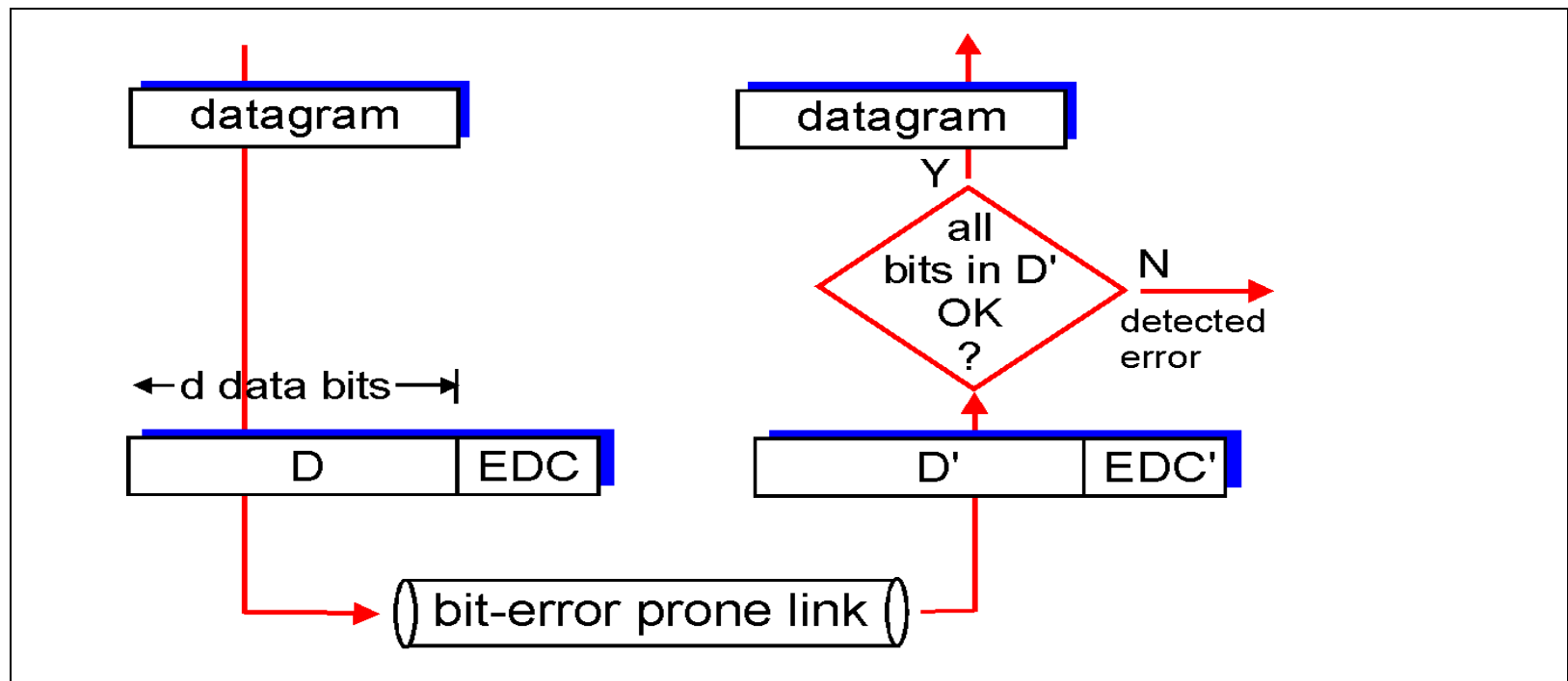


1. Silently return the wrong answer.
2. Detect failure.
3. Correct / mask the failure

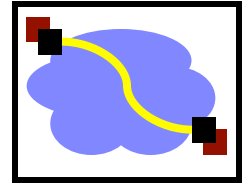
Block error detection/correction



- EDC= Error Detection and Correction bits (redundancy)
- D = Data protected by error checking, may include header fields
- Error detection not 100% reliable!
 - Protocol may miss some errors, but rarely
 - Larger EDC field yields better detection and correction



Parity Checking



Single Bit Parity:

Detect single bit errors

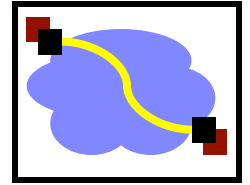
← d data bits → | parity
bit

0111000110101011	0
------------------	---

Calculated using XOR over data bits:

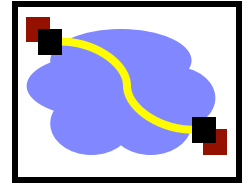
- 0 bit: even number of 0s
- 1 bit: odd number of 0s

Error Detection - Checksum



- Used by TCP, UDP, IP, etc..
- Ones complement sum of all 16-bits in packet
- Simple to implement
 - Break up packet into 16-bits strings
 - Sum all the 16-bit strings
 - Take complement of sum = checksum; add to header
 - One receiver, compute same sum, add sum and checksum, check that the result is 0 (no error)
- Relatively weak detection
 - Easily tricked by typical loss patterns

Example: Internet Checksum



- Goal: detect “errors” (e.g., flipped bits) in transmitted segment

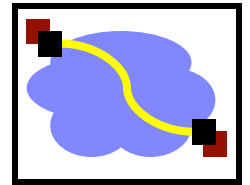
Sender

- Treat segment contents as sequence of 16-bit integers
- Checksum: addition (1's complement sum) of segment contents
- Sender puts checksum value into checksum field in header

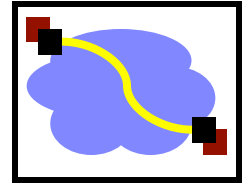
Receiver

- Compute checksum of received segment
- Check if computed checksum equals checksum field value:
 - NO - error detected
 - YES - no error detected. But maybe errors nonetheless?

Error Detection – Cyclic Redundancy Check (CRC)

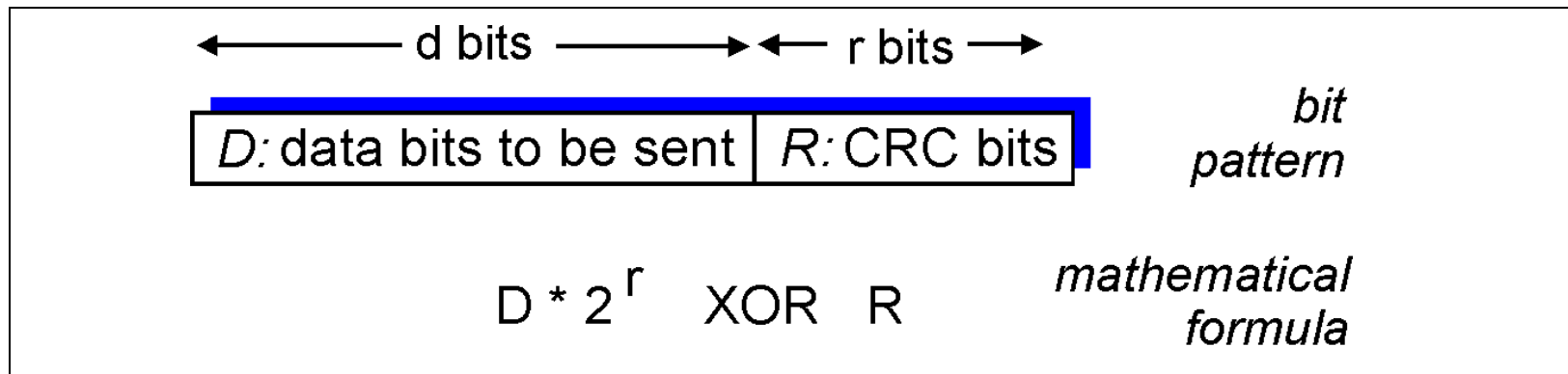


- Polynomial code
 - Treat packet bits as coefficients of n -bit polynomial
 - Choose $r+1$ bit generator polynomial (well known – chosen in advance)
 - Add r bits to packet such that message is divisible by generator polynomial
- Better loss detection properties than checksums
 - Cyclic codes have favorable properties in that they are well suited for detecting burst errors
 - Therefore, used on networks/hard drives

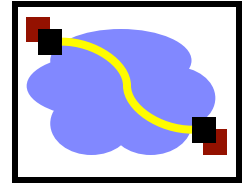


Error Detection – CRC

- View data bits, **D**, as a binary number
- Choose $r+1$ bit pattern (generator), **G**
- Goal: choose r CRC bits, **R**, such that
 - $\langle D, R \rangle$ exactly divisible by G (modulo 2)
 - Receiver knows G , divides $\langle D, R \rangle$ by G . If non-zero remainder: error detected!
 - Can detect all burst errors less than $r+1$ bits
- Widely used in practice



CRC Example



Want:

$$D \cdot 2^r \text{ XOR } R = nG$$

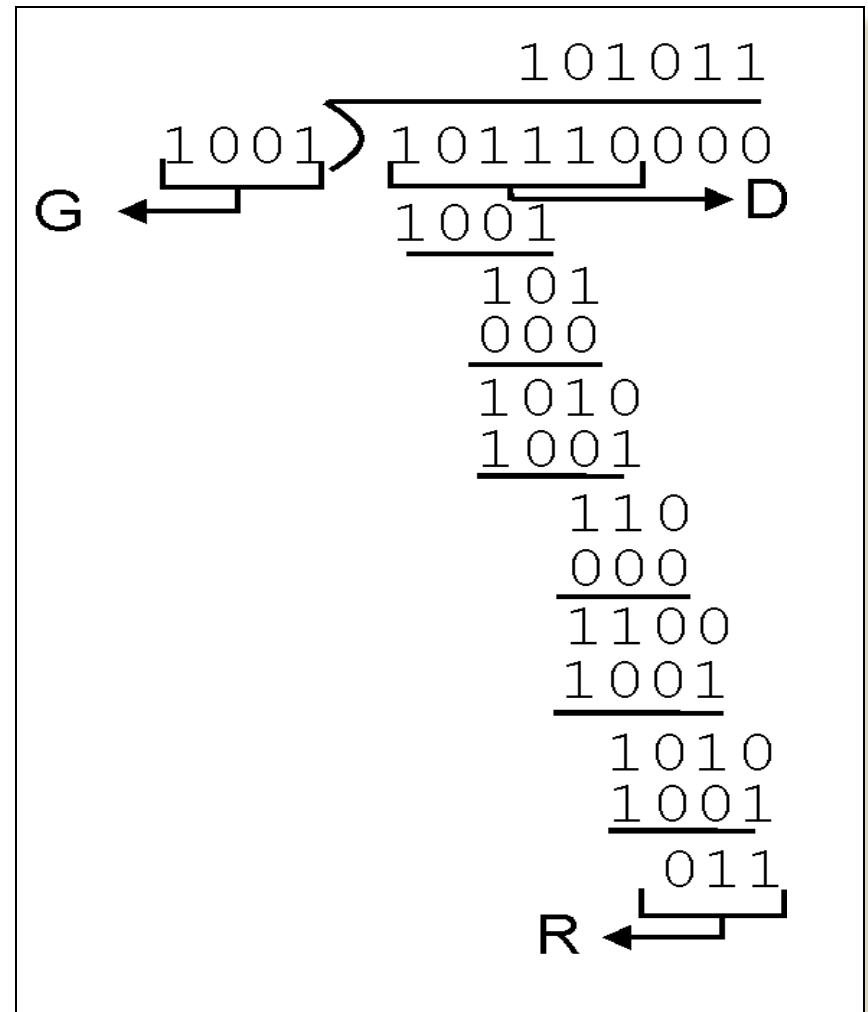
equivalently:

$$D \cdot 2^r = nG \text{ XOR } R$$

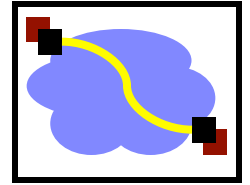
equivalently:

if we divide $D \cdot 2^r$ by G ,
want remainder R

$$R = \text{remainder} \left[\frac{D \cdot 2^r}{G} \right]$$

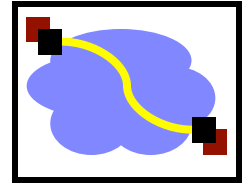


Options in dealing with failure



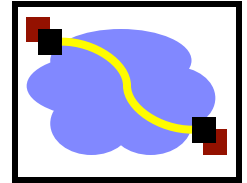
1. Silently return the wrong answer.
2. Detect failure.
3. Correct / mask the failure

Error Recovery



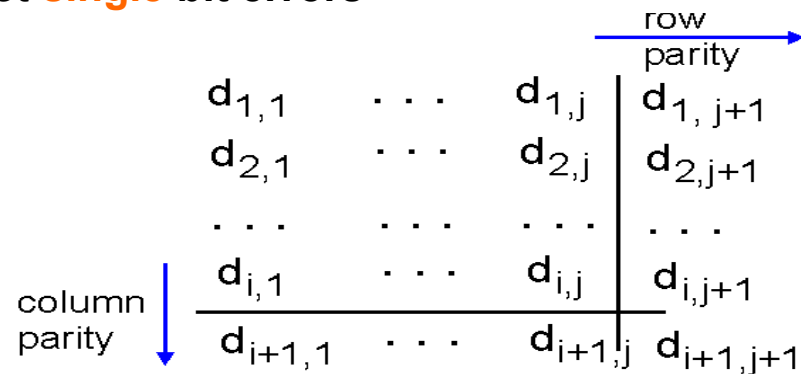
- Two forms of error recovery
 - Redundancy
 - Error Correcting Codes (ECC)
 - Replication/Voting
 - Retry
- ECC
 - Keep encoded redundant data to help repair losses
 - Forward Error Correction (FEC) – send bits in advance
 - Reduces latency of recovery at the cost of bandwidth

Error Recovery – Error Correcting Codes (ECC)



Two Dimensional Bit Parity:

Detect and correct **single** bit errors



1	0	1	0	1	1
1	1	1	1	0	0
0	1	1	1	0	1
0	0	1	0	1	0

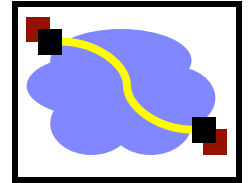
no errors

1	0	1	0	1	1
1	0	1	1	0	0
0	1	1	1	0	1
0	0	1	0	1	0

parity
error

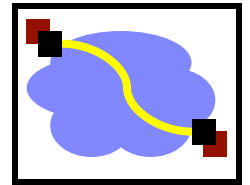
*correctable
single bit error*

Replication/Voting

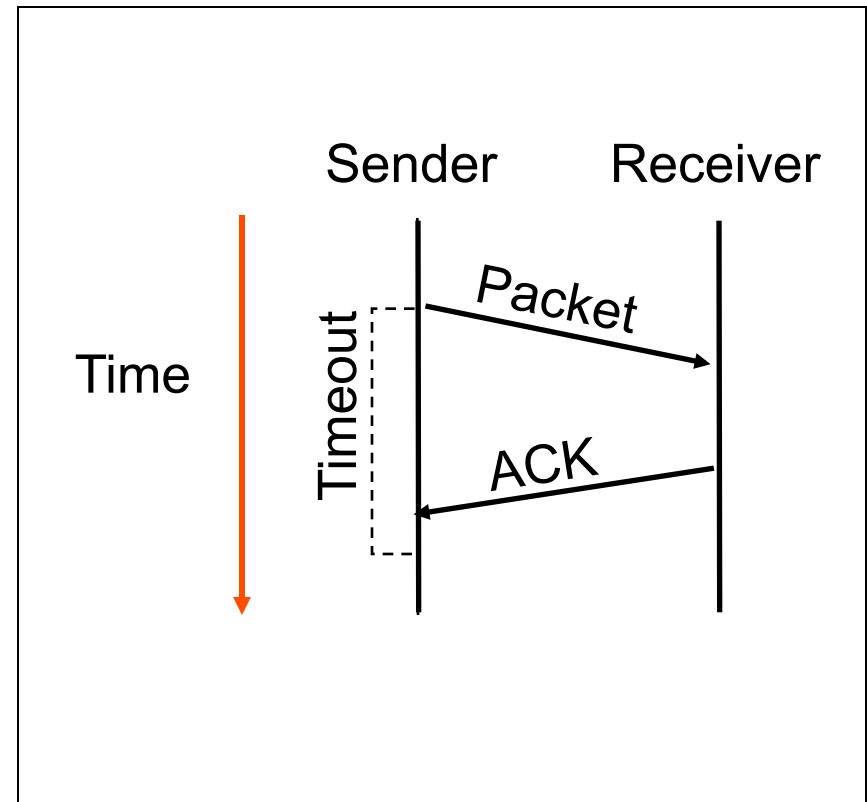


- If you take this to the extreme, three software versions:
[r1] [r2] [r3]
- Send requests to all three versions of the software: Triple modular redundancy
 - Compare the answers, take the majority
 - Assumes no error detection
- In practice - used mostly in space applications; some extreme high availability apps (stocks & banking? maybe. But usually there are cheaper alternatives if you don't need real-time)
 - Stuff we cover later: surviving malicious failures through voting (byzantine fault tolerance)

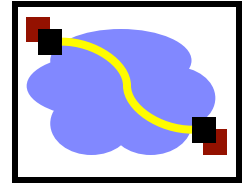
Retry – Network Example



- Sometimes errors are transient
- Need to have error detection mechanism
 - E.g., timeout, parity, checksum
 - No need for majority vote



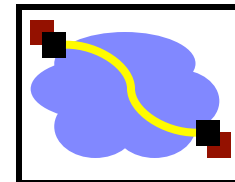
One key question



- How correlated are failures?
- Can you assume independence?
 - If the failure probability of a computer in a rack is p ,
 - What is $p(\text{computer 2 failing}) \mid \text{computer 1 failed}$?
 - Maybe it's p ... or maybe they're both plugged into the same UPS...
- Why is this important?

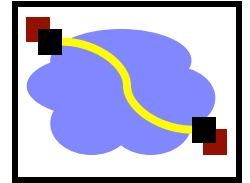
Back to Disks...

What are our options?



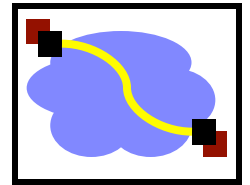
1. Silently return the wrong answer.
2. Detect failure.
 - Every sector has a header with a checksum. Every read fetches both, computes the checksum on the data, and compares it to the version in the header. Returns error if mismatch.
3. Correct / mask the failure
 - Re-read if the firmware signals error (may help if transient error, may not)
 - Use an error correcting code (what kinds of errors do they help?)
 - Bit flips? Yes. Block damaged? No
 - Have the data stored in multiple places (RAID)

Fail-fast disk



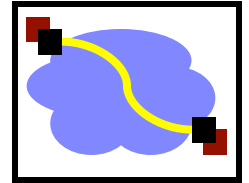
```
failfast_get (data, sn) {  
    get (s, sn);  
    if (checksum(s.data) = s.cksum) {  
        data ← s.data;  
        return OK;  
    } else {  
        return BAD;  
    }  
}
```

Careful disk (try 10 times on error)



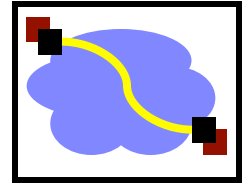
```
careful_get (data, sn) {  
    r ← 0;  
    while (r < 10) {  
        r ← failfast_get (data, sn);  
        if (r = OK) return OK;  
        r++;  
    }  
    return BAD;  
}
```

“RAID”



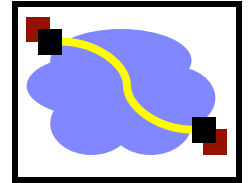
- Redundant Array of {Inexpensive, Independent} disks
- Replication! Idea: Write everything to two disks (“RAID-1”)
 - If one fails, read from the other
- `write(sector, data) ->`
 - `write(disk1, sector, data)`
 - `write(disk2, sector, data)`
- `read(sector, data)`
 - `data = read(disk1, sector)`
 - if error
 - `data = read(disk2, sector)`
 - if error, return error
 - return data
- Not perfect, though... doesn't solve all uncaught errors.

Durable disk (RAID 1)



```
 durable_get (data, sn) {  
     r ← disk1.careful_get (data, sn);  
     if (r = OK) return OK;  
     r ← disk2.careful_get (data, sn);  
     signal(repair disk1);  
     return r;  
 }
```

Summary



- Definition of MTTF/MTBF/MTTR: Understanding availability in systems.
- Failure detection and fault masking techniques
- Engineering tradeoff: Cost of failures vs. cost of failure masking.
 - At what level of system to mask failures?
 - Leading into replication as a general strategy for fault tolerance (more RAID next time)
- Thought to leave you with:
 - What if you have to survive the failure of entire computers? Of a rack? Of a datacenter?