## **Situation Calculus**

- State-based representation where the states are denoted by terms.
- > A situation is a term that denotes a state.
- > There are two ways to refer to states:
  - $\succ$  *init* denotes the initial state
  - > do(A, S) denotes the state resulting from doing action A in state S, if it is possible to do A in S.
- > A situation also encodes how to get to the state it denotes.

# **Example Situations**

🕨 init

 $\blacktriangleright$  do(pickup(rob, k1),

*do(move(rob, o*103, *mail)*, *do(move(rob, o*109, *o*103), *init))).* 

# Using the Situation Terms

Add an extra term to each dynamic predicate indicating the situation.

Example Atoms:

at(rob, o109, init)
at(rob, o103, do(move(rob, o109, o103), init))
at(k1, mail, do(move(rob, o109, o103), init))

## Axiomatizing using the Situation Calculus

- You specify what is true in the initial state using axioms with *init* as the situation parameter.
  - Primitive relations are axiomatized by specifying what is true in situation do(A, S) in terms of what holds in situation *S*.
  - Derived relations are defined using clauses with a free variable in the situation argument.

Static relations are defined without reference to the situation.

**Initial Situation** 

*sitting\_at(rob, o109, init).* 

sitting\_at(parcel, storage, init).

sitting\_at(k1, mail, init).

## **Derived Relations**

 $adjacent(P_1, P_2, S) \leftarrow$   $between(Door, P_1, P_2) \land$  unlocked(Door, S). adjacent(lab2, o109, S).

#### When are actions possible?

poss(A, S) is true if action A is possible in situation S.

 $poss(putdown(Ag, Obj), S) \leftarrow carrying(Ag, Obj, S).$ 

 $poss(move(Ag, Pos_1, Pos_2), S) \leftarrow$ autonomous(Ag)  $\land$ adjacent(Pos\_1, Pos\_2, S)  $\land$ sitting\_at(Ag, Pos\_1, S).

#### **Axiomatizing Primitive Relations**

**Example:** Unlocking the door makes the door unlocked:

 $unlocked(Door, do(unlock(Ag, Door), S)) \leftarrow poss(unlock(Ag, Door), S).$ 

Frame Axiom: No actions lock the door:

 $unlocked(Door, do(A, S)) \leftarrow$ 

 $unlocked(Door, S) \land$ 

poss(A, S).

Example: axiomatizing *carried* 

Picking up an object causes it to be carried:

 $carrying(Ag, Obj, do(pickup(Ag, Obj), S)) \leftarrow poss(pickup(Ag, Obj), S).$ 

**Frame Axiom:** The object is being carried if it was being carried before unless the action was to put down the object:

 $carrying(Ag, Obj, do(A, S)) \leftarrow$ 

 $carrying(Ag, Obj, S) \land$ 

 $poss(A, S) \land$ 

 $A \neq putdown(Ag, Obj).$ 



An object is sitting at a location if:

▶ it moved to that location:

sitting\_at(Obj, Pos, do(move(Obj, Pos<sub>0</sub>, Pos), S)) <
poss(move(Obj, Pos<sub>0</sub>, Pos).

▶ it was put down at that location:

sitting\_at(Obj, Pos, do(putdown(Ag, Obj), S))  $\leftarrow$ poss(putdown(Ag, Obj), S)  $\land$ at(Ag, Pos, S).

it was at that location before and didn't move and wasn't picked up.

#### More General Frame Axioms

The only actions that undo *sitting\_at* for object *Obj* is when *Obj* moves somewhere or when someone is picking up *Obj*.

sitting\_at(Obj, Pos, do(A, S))  $\leftarrow$ poss(A, S)  $\land$ sitting\_at(Obj, Pos, S)  $\land$   $\forall Pos_1 \ A \neq move(Obj, Pos, Pos_1) \land$  $\forall Ag \ A \neq pickup(Ag, Obj).$ 

The last line is equivalent to:

$$\sim \exists Ag A = pickup(Ag, Obj)$$

which can be implemented as

sitting\_at(Obj, Pos, do(A, S))  $\leftarrow$ ... $\land$  ... $\land$  ... $\land$ ~is\_pickup\_action(A, Obj).

with the clause:

$$is\_pickup\_action(A, Obj) \leftarrow$$
  
 $A = pickup(Ag, Obj).$ 

which is equivalent to:

is\_pickup\_action(pickup(Ag, Obj), Obj).

## **STRIPS** and the Situation Calculus

- Anything that can be stated in STRIPS can be stated in the situation calculus.
- The situation calculus is more powerful. For example, the "drop everything" action.
- To axiomatize STRIPS in the situation calculus, we can use holds(C, S) to mean that C is true in situation S.





 $holds(C, do(A, W)) \leftarrow$ The preconditions of preconditions $(A, P) \land$ of A all hold in W.  $holdsall(P, W) \land$  $add\_list(A, AL) \land$ C is on the addlist of A. member(C, AL). $holds(C, do(A, W)) \leftarrow$ preconditions(A, P)  $\land$ The preconditions of of A all hold in W. holdsall(P, W)  $\land$ C isn't on the  $delete\_list(A, DL) \land$  $notin(C, DL) \land$ deletelist of A. C held before A. holds(C, W).

-