

# CILOG User Manual

## Version 0.14

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### Abstract

This manual describes CILOG, a simple representation and reasoning system based on the book *Computational Intelligence: A Logical Approach* [1]. CILOG provides:

- a definite clause representation and reasoning system
- a simple tell-ask user interface, where the user can tell the system facts and ask questions of the system
- explanation facilities to explain how a goal was proved, why an answer couldn't be found, why a question was asked, why an error-producing goal was called, and why the depth-bound was reached
- knowledge-level debugging tools, that let the user debug incorrect answers, missing answers, why the system asks a question, system errors, and possible infinite loops
- depth-bounded search, that can be used to investigate potential infinite loops and used to build an iterative-deepening search procedure
- sound negation-as-failure, that interacts appropriately with the depth-bound
- ask-the-user facilities
- assumables

CILOG is a purely declarative representation and reasoning system. It is intended as a pedagogical tool to present a simple logic that can be used for AI problems. It is meant for programming in the small, where you can axiomatize a domain, ask questions and debug the knowledge base, without knowing how answers are produced. There may be sophisticated problem-solving techniques used in finding answers<sup>1</sup>.

We have many of the [examples from the book](#) available.

## 1 Getting and Starting CILOG

To start CILOG you find the cilog executable (`~cs322/cilog/cilog`) and run it. The following should work on the UBC undergraduate machines (or a similar setup at other locations). User input is in bold:

---

<sup>1</sup>There weren't any when this manual was written, but one of the design goals is that we should be able to use, for example, constraint satisfaction techniques to solve goals, without needing to change the user-level view of the system. It may just run faster (or not, as the case may be).

```
pender.ugrad.cs.ubc.ca% ~cs322/cilog/cilog
CILOG Version 0.14. Copyright 1998-2004, David Poole.
CILOG comes with absolutely no warranty.
All inputs end with a period. Type "help." for help.
cilog:
```

Alternatively, you can load the file *cilog.pl* and type “start.” Currently we have:

- *cilog\_swi.pl* for **SWI Prolog**, version 5.4 or later (or *cilog\_swi\_old.pl* for older versions of SWI Prolog). SWI Prolog runs on Windows and most Unix platforms (including Linux), and is free for non-commercial uses. SWI Prolog is available from <http://www.swi-prolog.org/>.
- *cilog\_sics.pl* for **Sicstus Prolog**. Loading this file should just start CILOG. This version is no longer supported.

The following shows a trace of starting CILOG assuming “pl” starts SWI Prolog:

```
[pender: cs322/cilog] 33 % pl -f ~cs322/cilog/cilog_sics.pl

CILOG Version 0.14. Copyright 1998-2004, David Poole.
CILOG comes with absolutely no warranty.
All inputs end with a period. Type "help." for help.
cilog:
```

## Version Changes

The following are the recent changes:

- 1.4 added *bagof*.
- 1.3 updated to latest version of SWI prolog, including delaying inequality.

## 2 Help and Exiting

The basic system prompt is:

```
cilog:
```

This means that CILOG is waiting for input from the user. All user inputs end with a period (“.”) and a return (enter).

At any stage you can ask for help by issuing a help command:

```
cilog: help.
```

This gives a brief description of the available commands.

To quit CILOG, you can issue the command:

```
cilog: quit.
```

CILOG is written in Prolog. To exit CILOG to the underlying Prolog, you can use the command: .

```
cilog: prolog.
```

To restart CILOG, you should use the command:

```
| ?- start.
```

at the prolog prompt. Note that this does not clear the knowledge base.

To turn on reporting the run time of queries, you can do

```
ciLOG: stats runtime.
```

On some systems, this is the elapsed time since the query was asked or the user asked for more answers, which may be a very inaccurate of run time if there was user interaction. To turn off reporting of run time, you can do:

```
ciLOG: stats none.
```

### 3 Syntax

The syntax of logical terms is based on Prolog's syntax in its convention for variables, but uses a different syntax for operators (because Prolog's are so confusing and to emphasize it is not Prolog).

A **variable** is a sequence of alphanumeric characters (possibly including “\_”) that starts with an upper case letter or “\_”. For example, X, Letter, Any\_cat, A\_big\_dog are all variables. The variable “\_” in an anonymous variable which means that all occurrences are assumed to be different variables. If a variable occurs only once in a clause, it should probably be written as “\_”.

A **constant** is either:

- a sequence of alphanumeric characters (possibly including “\_”) starting with a lower case letter, such as: david, comp\_intell, ciLOG, and a45\_23
- an integer or float, such as 123, -5, 1.0, -3.14159, 4.5E7, -0.12e+8, and 12.0e-9. There must be a decimal point in floats written with an exponent and at least one digit before and after a decimal point.
- any sequence of characters delimited by single quotes, such as 'X', '2b~2b', ' ../ch2/foo.pl ', 'A Tall Person'

A **term** is either a variable, a constant, or of the form  $f(t_1, \dots, t_n)$ , where  $f$ , a function symbol, is a sequence of alphanumeric characters (possibly including “\_”) starting with a lower case letter and the  $t_i$  are terms.

An **atom** is either of the form  $p$  or  $p(t_1, \dots, t_n)$ , where  $p$ , a predicate symbol, is a sequence of alphanumeric characters (including \_) starting with a lower case letter and the  $t_i$  are terms.

An **body** is either an atom, of the form  $\alpha\&\beta$ , where  $\alpha$  and  $\beta$  are bodies, or of the form  $\sim\alpha$ , where  $\alpha$  is a body.

A **clause** is either an atom or is a rule of the form  $h\leftarrow b$  where  $h$  is an atom (the head of the clause) and  $b$  is a body.

Some predicate and function symbols can be written using infix notation. For example “X is 4+3\*Y” means the same as “is(X, +(4, \*(3, Y)))”, where “is” is a predicate symbol and “+” and “\*” are function symbols. The operator precedence follows Prolog's conventions.

The symbol “<=” is defined to be infix, but there are no clauses defining it. This is designed to be used for meta-programming where “<=” can be used as a meta-level predicate symbol defining the object-level implication. (See *Computational Intelligence* [1], Chapter 6).

## 4 Ask and Tell

The general interaction with CILOG is to **tell** it clauses and to **ask** it queries (which are bodies). It replies with answers, which are instances of the query that are consequences of its knowledge base. The user can then either explore how the system derived an answer or why the system didn't find another answer.

To add a clause to the knowledge base you issue the command:

```
cilog: tell clause.
```

**Example 1** You could tell the system some genealogy knowledge:

```
cilog: tell parent(X,Y) <- mother(X,Y) .
cilog: tell parent(X,Y) <- father(X,Y) .
cilog: tell grandfather(X,Y) <- father(X,Z) & parent(Z,Y) .
cilog: tell grandmother(X,Y) <- mother(X,Z) & parent(Z,Y) .
cilog: tell father(randy,sally) .
cilog: tell father(randy,joe) .
cilog: tell mother(sally,mary) .
cilog: tell father(joe,sue) .
```

You can ask a query to determine if some expression is a consequence of the knowledge base:

```
cilog: ask query.
```

where *query* is an expression. The system will respond with either an instance of the query that is a consequence of the knowledge base or “No. *query* doesn't follow from the knowledge base” if there are no answers. When the system have given an answer you can reply with:

**ok.** to indicate that you don't want any more answers

**more.** to ask for more answers

**how.** to investigate how that answer was produced (see Section 6.1)

**help.** to get a menu of available answers (this option is available at all points where the system is asking for a response)

If you reply with **more**, the system either finds another answer or reports there are no more answers.

**Example 2** Given the knowledge base after Example 1, you can ask queries of the knowledge base:

```
cilog: ask grandfather(randy,mary) .
Answer: grandfather(randy,mary) .
[ok,more,how,help]: ok.
```

It has thus told you that *grandfather(randy,mary)* is a consequence of the knowledge base.

```
cilog: ask grandfather(joe,mary) .
No. grandfather(joe,mary) doesn't follow from the knowledge base.
```

which means *grandfather(joe,mary)* is not a consequence of the knowledge base.

You can ask queries with free variables:

```

cilog: ask grandfather(G,C).
Answer: grandfather(randy,mary).
      [ok,more,how,help]: more.
Answer: grandfather(randy,sue).
      [ok,more,how,help]: more.
No more answers.

```

This means there are two instances of the query *grandfather(G, C)* that are consequences of the knowledge base.

## 5 Loading and Listing Knowledge Bases

You can also tell the system a set of clauses from a file. The command:

```

cilog: load 'filename'.

```

where *'filename'* is the name (enclosed in single quotes) of a file that contains a sequence of clauses (each clause ending with a period), adds the clauses to the knowledge base as though they had been told to the system. Note that the file does *not* contain **tell** commands; it just contains the clauses.

Within files (and at any prompt), **comments** can be included between a “%” and the end of a line, or between “/\*” and “\*/”.

Loading a file adds the clauses to the database. It does not replace definitions. To clear the knowledge base you can issue the command:

```

cilog: clear.

```

To remove all clauses with head *atom* from the knowledge base, you issue the command:

```

cilog: clear atom.

```

To list the contents of the knowledge base you can issue the command:

```

cilog: listing.

```

This gives a listing of the whole knowledge base that is suitable for copying to a file and loading in a subsequent CILOG session.

To get a listing of all clauses in the knowledge base whose head unifies with *atom*, you can issue the command:

```

cilog: listing atom.

```

Note that in listings, and any time atoms with free variables are written, the variables are renamed. The system does not remember the names you gave the variables. It just calls them *A*, *B*, *C*, etc.

**Example 3** Suppose the file “**genealogy.cil**” contains the rules from Example 1. The following shows how the file can be loaded and listed.

```

cilog: load 'genealogy.cil'.
CILOG theory genealogy.cil loaded.
cilog: listing.
parent(A,B) <- mother(A,B).
parent(A,B) <- father(A,B).

```

```
grandfather(A,B) <- father(A,C)&parent(C,B).
grandmother(A,B) <- mother(A,C)&parent(C,B).
father(randy,sally).
father(randy,joe).
mother(sally,mary).
father(joe,sue).
cilog:
```

To load a file that uses Prolog syntax, you can issue the command:

```
cilog: prload 'filename'.
```

This assumes the clauses in the file use the Prolog implication :- and use a comma for conjunction. This does not allow the use of Prolog extra-logical predicates.

One useful command to check your knowledge base is the command:

```
cilog: check.
```

This lists the rules in the knowledge base with atoms in the body which don't unify with the head of any clause in the knowledge base (or any askable or assumable). This is a simple check that the rule can never be used in a proof. This is useful as this is usually an indication that there is a problem with the knowledge base.

This static check replaces the dynamic “undefined predicate” exception of Prolog, but is more useful in that it can be done before asking any queries, it finds clauses that contain atoms that must immediately fail (even if there are other clauses with the same predicate symbol). The warning can be ignored if that is what you intended (for example if you intend to add appropriate clauses later).

## 6 Explanation and Debugging

The main power of CILOG is in the explanation and debugging facilities. These fall into three classes:

- determining how an answer was proved
- determining why a potential answer wasn't produced
- determining why the search is in a particular state

Each of these is described in the following sections.

### 6.1 How? Questions

When the system has derived an answer you can ask **how** that answer was produced. This provides a mechanism for interactively traversing a proof tree. At each stage the system either says why the answer was produced immediately (e.g., if it was a fact or a built-in predicate — see Section 8) or produces the instance of the top-level rule that was used to prove the goal.

When you ask how atom  $h$  was proved, it produces the instance of the rule in the knowledge base with  $h$  as the head that succeeded:

```
 $h$  <-
  1:  $a_1$ 
  2:  $a_2$ 
  ...
  k:  $a_k$ 
```

which indicates that the rule  $h \leftarrow a_1 \& a_2 \& \dots \& a_k$  was used to prove  $h$ . You can then give one of:

**how**  $i$ . where  $i$  is an integer  $1 \leq i \leq k$ . This means that you want to see how  $a_i$  was proved.

**up**. to go back to see the rule with  $h$  in the body. If  $h$  is the initial query it goes back to the answer interaction.

**retry**. to ask for a different proof tree.

**ok**. to exit the how traversal and go back to the answer prompt.

**help**. for a list of the available commands.

If  $a_i$  was not proved via a rule (for example, if it was a fact or a built-in predicate), the reason that  $a_i$  was proved is printed and then the rule with  $a_i$  in the body is printed again.

**Example 4** Given the knowledge base after Example 1, you can ask how a particular instance of the query was proved, as in:

```

cilog: ask grandfather(G,C).
Answer: grandfather(randy,mary) .
      [ok,more,how,help]: how.
      grandfather(randy,mary) <-
        1: father(randy,sally)
        2: parent(sally,mary)
      How? [number,up,retry,ok,help]: how 2.
      parent(sally,mary) <-
        1: mother(sally,mary)
      How? [number,up,retry,ok,help]: how 1.
      mother(sally,mary) is a fact
      parent(sally,mary) <-
        1: mother(sally,mary)
      How? [number,up,retry,ok,help]: up.
      grandfather(randy,mary) <-
        1: father(randy,sally)
        2: parent(sally,mary)
      How? [number,up,retry,ok,help]: how 1.
      father(randy,sally) is a fact
      grandfather(randy,mary) <-
        1: father(randy,sally)
        2: parent(sally,mary)
      How? [number,up,retry,ok,help]: up.
Answer: grandfather(randy,mary) .
      [ok,more,how,help]: more.
Answer: grandfather(randy,sue) .
      [ok,more,how,help]: ok.
cilog:

```

## 6.2 Whynot? Questions

Just as how questions let you explore the proof tree for a particular derivation, **whynot** questions let you explore the search tree. This facility is mainly used for determining why there was no proof for a particular query. As such the documentation is written assuming that you are trying to determine why a query failed, when you thought that it should have succeeded.

If a query fails when it should have succeeded, either there was a rule defining the query whose body fails when it should have succeeded, or else there is a missing rule for that query.

You can ask a query to determine why some query failed using:

```
cilog: whynot query.
```

where *query* is an expression.

If the query is an atom, you can examine each rule whose head unifies with the query. For each such rule, the system asks whether you want to trace this rule. You can reply with:

**yes.** to determine why this rule failed. You should give this answer when this rule should have succeeded for this query.

**no.** to ask for another rule. You give this answer if this rule should have failed for this query.

**up.** to return to a previous choice point (the rule in which the atom appears, or else the top-level if the atom was the initial query).

**ok.** to return to the top-level.

**help.** to get a menu of available answers.

If you answer “**no**” to each rule, this means that all of the rules in the knowledge base should have failed, and so the appropriate rule for the query that should have succeeded is missing.

To determine why a rule failed, we determine why the body failed. If the body is atomic, we use the above *whynot* mechanism to determine why the rule failed. If the body of the rule is a conjunction,  $\alpha \& \beta$ , there are four cases:

- $\alpha$  fails, in which case we should use recursively use the *whynot* mechanism to determine why it failed.
- an instance of  $\alpha$  succeeds, but should not have succeeded. In this case we can use the *how* mechanism to debug this proof. This case arises because  $\beta$  may have rightfully failed on the instance of  $\alpha$  that succeeded.
- an instance of  $\alpha$  succeeds, but this instance should have lead to failure of the conjunction. In this case we should look for another proof for  $\alpha$ .
- an instance of  $\alpha$  succeeds that should lead to the success of the conjunction, in which case we need to determine why  $\beta$  failed on this instance.

Thus, when there is a conjunctive body, we first try to prove the leftmost conjunct. If it fails, we use the above *whynot* mechanism to determine why it failed. If it succeeds, the user is asked *Should this answer lead to a successful proof?* The user can reply:

**yes.** this instance should have made the body succeed. Thus you need to debug the rest of the conjunction.

**no.** this instance should lead to a failure of the body. Thus you need to try another proof for this atom.



**debug.** this instance is false, debug it. This enters the *how* interaction.

**ok.** to return to the top-level.

**help.** to get a menu of available answers.

**Example 5** Suppose we had the knowledge base from Example 1. Suppose that, the user knew that Joe had a child called Jane, and wanted to know why the system didn't think that Randy was Jane's grandfather. We could imagine the following dialogue:

```

cilog: whynot grandfather(randy, jane) .
      grandfather(randy, jane) <- father(randy, A)&parent(A, jane) .
      Trace this rule? [yes,no,up,help]: yes.
      The proof for father(randy, sally) succeeded.
      Should this answer lead to a successful proof?
      [yes,no,debug,help]: no.
      The proof for father(randy, joe) succeeded.
      Should this answer lead to a successful proof?
      [yes,no,debug,help]: yes.
      parent(joe, jane) <- mother(joe, jane) .
      Trace this rule? [yes,no,up,help]: no.
      parent(joe, jane) <- father(joe, jane) .
      Trace this rule? [yes,no,up,help]: yes.
      There is no applicable rule for father(joe, jane) .
      parent(joe, jane) <- father(joe, jane) .
      Trace this rule? [yes,no,up,help]: up.
      grandfather(randy, jane) <- father(randy, joe)&parent(joe, jane) .
      Trace this rule? [yes,no,up,help]: up.
cilog:

```

### 6.3 Depth Bound

The search strategy used by CILOG is a depth-bounded depth-first search. The depth bound is set using the command:

```

cilog: bound n.

```

where *n* is a positive integer. *n* is a bound on the depth of any proof tree. The default value is 30.

When a goal is asked and no more answers can be found, if the depth-bound was not reached, the system reports there are no more answers. If the depth-bound was reached, there still may be some answers, although it is more likely that there is a bug in the knowledge base where some recursions do not terminate. CILOG allows the user to interactively explore the places where the depth bound was reached.

When no more answers can be found and the search was cut off due to hitting the depth-bound, the user is informed of this, as is given the option of one of:

*i.* where *i* is an integer bigger than the current depth-bound. This lets the user explore larger search trees. Repeatedly increasing the depth-bound lets the user simulate an iterative deepening search.

**where.** to let the user explore the place where the depth-bound was reached.

**ok.** to go back to the answer prompt.

**help.** for a list of the available commands.

When the user given the **where** command, the proof is retried, and it halts at a point where the depth-bound was reached, and shows the user the current subgoal,  $g$ . The user can give one of the following commands:

**fail.** to fail  $g$ , and explore other places where the depth-bound was reached.

**succeed.** to say that  $g$  should succeed.

**proceed.** to fail  $g$  and not explore any more subgoals where the depth-bound was reached.

**why.** to let the user explore why  $g$  was called (see Section 6.4).

**ok.** to go back to the answer prompt.

**help.** for a list of the available commands.

## 6.4 Why? Questions

If you find yourself at a particular subgoal, it is often useful to find out **why** that subgoal is being proved.

When you ask why an atom was being asked, CILOG produces the instance of the rule in the knowledge base with the atom in the body, such that the head was trying to be proved. This is of the form:

```

h :-
  1: a1
  2: a2
  ...
** j: aj
  ...
  k: ak

```

This means that  $a_j$  is the atom that is being asked, the atoms  $a_1 \dots a_{j-1}$  have been proved, and  $a_{j+1} \dots a_k$  have still to be proved. When this is presented you can ask one of:

**how  $i$ .** where  $i$  is an integer  $1 \leq i < j$ . This means that you want to see how  $a_i$  was proved. This enters the how dialog.

**how  $j$ .** This means that you want to see how CILOG is trying to prove  $a_j$  (the atom you have previously asked *why* about). This returns to the rule with  $a_j$  in the body.

**why.** to see the rule with  $h$  in the body.

**prompt.** to return to the CILOG prompt.

**help.** for a list of the available commands.

Note that you can't ask **how  $i$**  for  $i > j$  as there is no proof tree for the atoms  $a_{j+1} \dots a_k$ . Also, CILOG does not guarantee that the rule given will have the atoms in the body in the same order as in the knowledge base; CILOG is free to change the order of atoms in a body as long as this doesn't introduce an error.

**Example 6** The following shows a trace of a looping program that reaches the depth bound. We first increase the depth-bound, and then explore where the depth-bound was reached.

```

cilog: tell a <- b & c & d.
cilog: tell b <- e.
cilog: tell e.
cilog: tell c <- f & d & g.
cilog: tell f <- b & h.
cilog: tell h.
cilog: tell d <- a.
cilog: ask a.
Query failed due to depth-bound 30.
  [New-depth-bound,where,ok,help]: 50.
Query failed due to depth-bound 50.
  [New-depth-bound,where,ok,help]: where.
Depth-bound reached. Current subgoal: e
  [fail,succeed,proceed,why,ok,help]: why.
e is used in the rule
b <-
** 1: e
  [Number,why,help,ok]: why.
b is used in the rule
a <-
** 1: b
   2: c
   3: d
  [Number,why,help,ok]: why.
a is used in the rule
d <-
** 1: a
  [Number,why,help,ok]: why.
d is used in the rule
c <-
   1: f
** 2: d
   3: g
  [Number,why,help,ok]: how 1.
f <-
   1: b
   2: h
How? [Number,up,retry,ok,help]: how 2.
h is a fact
f <-
   1: b
   2: h
How? [Number,up,retry,ok,help]: up.
d is used in the rule
c <-
   1: f
** 2: d
   3: g

```

```

[Number, why, help, ok]: why.
c is used in the rule
a <-
  1: b
  ** 2: c
  3: d
[Number, why, help, ok]: ok.
Depth-bound reached. Current subgoal:e
  [fail, succeed, proceed, why, ok, help]: fail.
No more answers.
cilog:

```

## 7 Ask-the-user

The ask-the-user facility exists, but is poorly documented in this manual. The subgoals have to be ground before being asked of the user. This will be fixed in future implementations.

To to make an atom askable, you can issue the command:

```
cilog:askable atom.
```

Whenever a ground instance of the atom is attempted to be proved, the user is asked if it is true. The system may ask:

```
Is g true?
```

The user can reply with one of:

**yes.** if  $g$  is known to be true. The system will not ask this instance again.

**no.** if  $g$  is known to be false. The system will not ask this instance again.

**unknown.** if  $g$  is unknown. In this case any applicable clauses for  $g$  can be used.

**fail.** to fail the subgoal (but not record an answer). This is only used to test you axiomatization.

**why.** to see why the system was asking this question. This then enters the **why** interaction described in Section 6.4.

**prompt.** to return to the cilog prompt.

**help.** to get a menu of available answers.

Note the assumption that CILOG makes about the interaction between asking the user about  $g$  and using clauses for  $g$ . It assumes that if the user knows whether  $g$  is true, the user's answer should be used. Any rules for  $g$  are used only if the user doesn't know whether  $g$  is true or not.

**Example 7** The following gives the example of the [electrical example with askables](#) (see [1], page 214):

```

cilog: load 'cilog_code/ch6/elect_askable.pl'.
CILOG theory cilog_code/ch6/elect_askable.pl loaded.
cilog: ask lit(L).
Is up(s2) true? [yes,no,unknown,why,help]: yes.

```

```

Is up(s1) true? [yes,no,unknown,why,help]: no.
Is down(s2) true? [yes,no,unknown,why,help]: no.
Is up(s3) true? [yes,no,unknown,why,help]: yes.
Answer: lit(l2).
      [ok,more,how,help]: ok.

```

## 8 Built-in predicates

There are a few built-in predicates. These cannot be redefined. The following arithmetic predicates are predefined:

- $X < Y$ , where  $X$  and  $Y$  must be ground arithmetic expressions when called, is true if the value of  $X$  is less than the value of  $Y$ .
- $X > Y$ , where  $X$  and  $Y$  must be ground arithmetic expressions when called, is true if the value of  $X$  is greater than the value of  $Y$ .
- $X \leq Y$ , where  $X$  and  $Y$  must be ground arithmetic expressions when called, is true if the value of  $X$  is less than or equal to the value of  $Y$ .
- $X \geq Y$ , where  $X$  and  $Y$  must be ground arithmetic expressions when called, is true if the value of  $X$  is greater than or equal to the value of  $Y$ .
- $X \neq Y$ , where  $X$  and  $Y$  must be ground arithmetic expressions when called, is true if the value of  $X$  is different to the value of  $Y$ .
- $V \text{ is } E$  where  $E$  must be a ground arithmetic expression when called. This is true if expression  $E$  evaluates to  $V$ .
- $\text{number}(N)$ , where  $N$  is ground is true if  $N$  is a number.

Arithmetic expressions can either be numbers or of the form  $(E_1 + E_2)$ ,  $(E_1 - E_2)$ ,  $(E_1 * E_2)$ ,  $(E_1/E_2)$ ,  $\sin(E)$ ,  $\log(E)$ , etc.

**Example 8** The following shows a simple use of `is` and `>`. It shows how we can use variables in the expressions, as long as they are bound before (to the left) they are used.

```

cilog: ask X is 3+4*5 & X*2 > 11*3.
Answer: 23 is 3+4*5 & 23*2>11*3.

```

The restriction that the arithmetic expressions must be ground when called is severe, and probably won't be fixed in future implementations.

There is one non-arithmetic comparison:

- $X \neq Y$  true if  $X$  and  $Y$  denote different objects under the unique names assumption.

$X$  and  $Y$  must be bound enough to determine whether they can be unified. (Delaying is not implemented). This may be fixed in future versions.

## 9 Assumables

Assumables exist as in Chapter 9 of Computational Intelligence. There are no integrity constraints that are handled specially. The interaction between assumables and negation as failure is unsatisfactory, in that  $\sim g$  succeeds if there is provably no explanation of  $g$ .

To to make an atom assumable, you can issue the command:

```
cilog: assumable atom.
```

When the atom is encountered, it is assumed to be true. The assumptions used to make a proof go though are collected and presented to the user when the answer is returned.

```
cilog: tell a(X) <- b(X) & c(X).
```

```
cilog: assumable c(X).
```

```
cilog: tell b(X) <- d(X) & e(X).
```

```
cilog: askable e(X).
```

```
cilog: tell d(a).
```

```
cilog: tell d(b).
```

```
cilog: tell d(c).
```

```
cilog: ask a(X).
```

```
Is e(a) true? [yes,no,unknown,why,help]: yes.
```

```
Answer: a(a).
```

```
Assuming: [c(a)].
```

```
[more,ok,how,help]: more.
```

```
Is e(b) true? [yes,no,unknown,why,help]: no.
```

```
Is e(c) true? [yes,no,unknown,why,help]: yes.
```

```
Answer: a(c).
```

```
Assuming: [c(c)].
```

```
[more,ok,how,help]: more.
```

```
No more answers.
```

```
cilog:
```

## 10 Negation-as-failure

Negation as finite failure is implemented. Use  $\sim G$  to mean the negation of expression  $G$ .  $G$  must be ground when  $\sim G$  is called; this may be fixed in future implementations. The interaction of negation as failure and the depth-bound is handled correctly.

There is a problem with negation and **why** questions; sometimes the rule written out omits some negation symbols. This only occurs when non-atomic formulae are negated. The interaction of negation as failure with assumables isn't satisfactory. It is recommended that you don't use both assumables and negation as failure.

## 11 Finding All Answers

You can collect all of the answers to a query using:

$$\text{bagof}(X, Q, L)$$

asks where  $X$  is a term made of the free variables in  $Q$ . This is true if  $L$  is a non-empty list of the  $X$ 's for which  $Q$  is true. There is an element of  $L$  for each *proof* of  $Q$ . This fails if  $Q$  has no answers.

```
cilog: tell p(a,b).
cilog: tell p(b,c).
cilog: tell p(a,d).
cilog: tell p(e,f).
cilog: tell p(e,d).
cilog: ask bagof(X,p(a,X),L).
Answer: bagof(A, p(a, A), [b, d]).
```

Note that  $X$  can be any term. For example, you could ask for the list of pairs:

```
cilog: ask bagof(ppair(X,Y),p(X,Y),L).
Answer: bagof(ppair(A, B), p(A, B), [ppair(a, b), ppair(b, c), ppair(a, d),
ppair(e, f), ppair(e, d)]).
```

$Q$  can be any body. For example, continuing our example:

```
cilog: ask bagof(triple(X,Y,Z),p(X,Y)&p(Y,Z),L).
Answer: bagof(triple(A, B, C), p(A, B)&p(B, C), [triple(a, b, c)]).
Runtime since last report: 0 secs.
[ok,more,how,help]: ok.
cilog: tell p(c,e).
cilog: ask bagof(triple(X,Y,Z),p(X,Y)&p(Y,Z),L).
Answer: bagof(triple(A, B, C), p(A, B)&p(B, C),
[triple(a, b, c), triple(b, c, e), triple(c, e, f), triple(c, e,
d)]).
Runtime since last report: 0 secs.
[ok,more,how,help]:
```

If  $Q$  contains a variable that doesn't appear in  $X$ , then this returns a list for each value that has a non-empty list of solutions. For example,

```
cilog: ask bagof(X,p(Y,X),L).
Answer: bagof(A, p(b, A), [c]).
Runtime since last report: 0 secs.
[ok,more,how,help]: more.
Answer: bagof(A, p(a, A), [b, d]).
Runtime since last report: 0 secs.
[ok,more,how,help]: more.
Answer: bagof(A, p(e, A), [f, d]).
Runtime since last report: 0 secs.
[ok,more,how,help]: more.
Answer: bagof(A, p(c, A), [e]).
Runtime since last report: 0 secs.
[ok,more,how,help]: more.
No more answers.
Runtime since last report: 0 secs.
cilog:
```

If you want to ask for the  $X$ 's for which there exists a  $Y$  such that  $P(Y, X)$  (i.e., if you don't want an answer for each  $Y$ , you can use the operator  $V \hat{=} Q$  which means "there exists a  $V$  such that  $Q$  is true". [Note that  $\hat{=}$  is only defined in the context of *bagof*. It doesn't need to be used elsewhere.]

```
cilog: ask bagof(X, Y^p(Y, X), L) .  
Answer: bagof(A, B^p(B, A), [b, c, d, f, d, e]).  
Runtime since last report: 0 secs.  
[ok, more, how, help]: more.  
No more answers.  
Runtime since last report: 0 secs.  
cilog:
```

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## References

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