#### Semantics: General Idea

A semantics specifies the meaning of sentences in the language. An interpretation specifies:

- what objects (individuals) are in the world
- the correspondence between symbols in the computer and objects & relations in world
  - constants denote individuals
  - predicate symbols denote relations

#### Formal Semantics

An interpretation is a triple  $I = \langle D, \phi, \pi \rangle$ , where

- D, the domain, is a nonempty set. Elements of D are individuals.
- $\phi$  is a mapping that assigns to each constant an element of D. Constant c denotes individual  $\phi(c)$ .
- $\pi$  is a mapping that assigns to each *n*-ary predicate symbol a relation: a function from  $D^n$  into  $\{TRUE, FALSE\}$ .

### **Example Interpretation**

Constants: phone, pencil, telephone.

Predicate Symbol: noisy (unary), left\_of (binary).

- $D = \{ \sim, \sim, \infty \}.$
- $\phi(phone) = \mathbf{\hat{a}}$ ,  $\phi(pencil) = \mathbf{\hat{a}}$ ,  $\phi(telephone) = \mathbf{\hat{a}}$ .
- $\pi(noisy)$ :  $\langle \mathcal{F} \rangle$  FALSE  $\langle \mathcal{T} \rangle$  TRUE  $\langle \mathcal{D} \rangle$  FALSE  $\pi(left\_of)$ :

$\langle \gg, \gg \rangle$	FALSE	⟨≫,☎⟩	TRUE	$\langle \mathbf{pprox}, \mathbf{\textcircled{N}} \rangle$	TRUE
⟨☎,≫⟩	FALSE	$\langle \mathbf{\Delta}, \mathbf{\Delta} \rangle$	FALSE	$\langle \mathbf{\Delta}, \mathfrak{D} \rangle$	TRUE
$\langle \mathfrak{D}, \mathfrak{S} \rangle$	FALSE	$\langle \mathfrak{D}, \mathbf{\Delta}  angle$	FALSE	$\langle \mathfrak{D}, \mathfrak{D} \rangle$	FALSE



### Important points to note

- The domain D can contain real objects. (e.g., a person, a room, a course). D can't necessarily be stored in a computer.
- π(p) specifies whether the relation denoted by the n-ary predicate symbol p is true or false for each n-tuple of individuals.
- If predicate symbol p has no arguments, then  $\pi(p)$  is either TRUE or FALSE.

# Truth in an interpretation

A constant c denotes in I the individual  $\phi(c)$ . Ground (variable-free) atom  $p(t_1, \ldots, t_n)$  is

- true in interpretation I if  $\pi(p)(\langle \phi(t_1), \ldots, \phi(t_n) \rangle) = \mathit{TRUE}$  in interpretation I and
- false otherwise.

Ground clause  $h \leftarrow b_1 \wedge \ldots \wedge b_m$  is false in interpretation I if h is false in I and each  $b_i$  is true in I, and is true in interpretation I otherwise.



### **Example Truths**

In the interpretation given before, which of following are true?

```
noisy(phone) \\ noisy(telephone) \\ noisy(pencil) \\ left\_of(phone, pencil) \\ left\_of(phone, telephone) \\ noisy(phone) \leftarrow left\_of(phone, telephone) \\ noisy(pencil) \leftarrow left\_of(phone, telephone) \\ noisy(pencil) \leftarrow left\_of(phone, pencil) \\ noisy(phone) \leftarrow noisy(telephone) \land noisy(pencil) \\ \end{cases}
```

### **Example Truths**

In the interpretation given before, which of following are true?

```
noisy(phone)
                                                            true
noisy(telephone)
                                                            true
noisy(pencil)
                                                            false
left_of(phone, pencil)
                                                            true
left_of(phone, telephone)
                                                            false
noisy(phone) \leftarrow left\_of(phone, telephone)
                                                            true
noisy(pencil) \leftarrow left\_of(phone, telephone)
                                                            true
noisy(pencil) \leftarrow left\_of(phone, pencil)
                                                            false
noisy(phone) \leftarrow noisy(telephone) \land noisy(pencil)
                                                            true
```

# Models and logical consequences (recall)

- A knowledge base, KB, is true in interpretation I if and only if every clause in KB is true in I.
- A model of a set of clauses is an interpretation in which all the clauses are true.
- If KB is a set of clauses and g is a conjunction of atoms, g is a logical consequence of KB, written  $KB \models g$ , if g is true in every model of KB.
- That is,  $KB \models g$  if there is no interpretation in which KB is true and g is false.

#### User's view of Semantics

- 1. Choose a task domain: intended interpretation.
- 2. Associate constants with individuals you want to name.
- 3. For each relation you want to represent, associate a predicate symbol in the language.
- 4. Tell the system clauses that are true in the intended interpretation: axiomatizing the domain.
- 5. Ask questions about the intended interpretation.
- 6. If  $KB \models g$ , then g must be true in the intended interpretation.

# Computer's view of semantics

- The computer doesn't have access to the intended interpretation.
- All it knows is the knowledge base.
- The computer can determine if a formula is a logical consequence of KB.
- If  $KB \models g$  then g must be true in the intended interpretation.
- If  $KB \not\models g$  then there is a model of KB in which g is false. This could be the intended interpretation.



#### Role of Semantics in an RRS

