Heuristic Search: BestFS and A*

Computer Science cpsc322, Lecture 8

(Textbook Chpt 3.6)

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Lecture Overview

Recap / Finish Heuristic Function

- · Best First Search
- A*

How to Combine Heuristics

i⊷licker.

If $h_1(n)$ is admissible and $h_2(n)$ is also admissible then

- A. min($h_1(n)$, $h_2(n)$) is also admissible and dominates its components (doesn't dominate)
- B. sum(h1(n), h2(n)) is also admissible and dominates its components (may not be admissible)
- C. $avg(h_1(n), h_2(n))$ is also admissible and dominates its components (desu' + Jom, met)D. None of the above

Combining Admissible Heuristics

How to combine heuristics when there is no dominance?

If $h_1(n)$ is admissible and $h_2(n)$ is also admissible then $h(n) = \frac{m \delta x}{(h_1 / h_2)}$ is also admissible

··· and dominates all its components



Example Heuristic Functions (2)

 Another one we can use the number of moves between each tile's current position and its position in the solution



= 18



Goal node

Combining Heuristics: Example

In 8-puzzle, solution cost for the 1,2,3,4 subproblem is substantially more accurate than sum of Manhattan distance of each tile from its goal position in some cases

So.....

Admissible heuristic for Vacuum world?



<u>states?</u> Where it is dirty and robot location <u>actions?</u> *Left, Right, Suck* <u>Possible goal test?</u> no dirt at all locations



<u>states?</u> Where it is dirty and robot location <u>actions?</u> <u>*Left*</u>, <u>*Right*</u>, <u>*Suck*</u> <u>Possible goal test?</u> no dirt at all locations

Lecture Overview

· Recap Heuristic Function

Best First Search

• A*

Best-First Search

- Idea: select the path whose end is closest to a goal according to the heuristic function.
- Best-First search selects a path on the frontier with minimal *h*-value (for the end node).
- It treats the frontier as a priority queue ordered by h. (similar to ?) $L \subset F \leq C_{by} \subset s t$
- This is a greedy approach: it always takes the path which appears locally best

Analysis of Best-First Search

 Not Complete : a low heuristic value can mean that a cycle gets followed forever.



Lecture Overview

- · Recap Heuristic Function
- · Best First Search
- A* Search Strategy

How can we effectively use h(n) i∗clicker.

Maybe we should combine it with the cost. How? Shall we select from the frontier the path *p* with:

- A. Lowest cost(p) h(p)
- B. Highest cost(p) h(p)

C. Highest cost(p)+h(p)

D. Lowest cost(p)+h(p)

A* Search Algorithm

Cost

- A* is a mix of:
 - lowest-cost-first and
 - best-first search

• A^* treats the frontier as a priority queue ordered by f(p) = cost(p) + h(p) is an estimate

• It always selects the node on the frontier with the

h estimate of shortest path from end of p to a Goal

Computing f-values



Analysis of A*

i⊳clicker.

If the heuristic is completely uninformative and the edge costs are all the same, A* is equivalent to....



- C. DFS
- D. None of theAbove

Analysis of A* for our states heuristic is equal to 0

 $\forall s h(s) = 0$

Let's assume that arc costs are strictly positive.

- Time complexity is O(b^m)
 - the heuristic could be completely uninformative and the edge costs could all be the same, meaning that A* does the same thing as....
 DFS BFS LCFS
- Space complexity is $O(b^m)$ like A^* maintains a frontier which grows with the size of the tree
- Completeness: yes.
- Optimality: ??

Optimality of A^*

If A^* returns a solution, that solution is guaranteed to be optimal, as long as

When

- the branching factor is finite
- arc costs are strictly positive
- h(n) is an underestimate of the length of the shortest path from *n* to a goal node, and is non-negative

admissible

Theorem

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If A^* selects a path p as the solution,
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p is the shortest (i.e., lowest-cost) path.
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Why is A^* optimal?

A* returns p

Assume for contradiction that some other path p' is actually the shortest p<u>ath t</u>o a goal

 $\operatorname{cost}(p) > \operatorname{cost}(p')$

Consider the moment when p is chosen from the frontier Some part of path p' will also be on the frontier; let's call this partial path start *р*".

p" think path for any path stort to stort here is (of that path) there is (of that path) there is (of that path) or state subpath p' of ways provider G p 203

Why is A* optimal? (cont') p" р $cost(p)+h(p) \leq cost(p')+h$ derestimate of Because p was expanded before p", $f(p) \leq f(p'')$ Because p is a goal, h(p) = 0 Thus $(ost(p) \leq cost(p'') + h(p''))$ Because *h* is admissible, $cost(p'') + h(p'') \le cost(p)$ for any path p'to a goal that extends p''Thus $cost(p) \leq cost(p')$ for any other path p'to a goal. $_{->}$ cost(p') < cost(p) This contradicts our assumption that p' is the shortest path.

Optimal efficiency of A^*

- In fact, we can prove something even stronger about
 A*: in a sense (given the particular heuristic that is available) no search algorithm could do better!
- Optimal Efficiency: Among all optimal algorithms that start from the same start node and use the same heuristic h, A* expands the minimal number of paths.

Samples A* applications

- An Efficient A* Search Algorithm For Statistical Machine Translation. 2001
- The Generalized A* Architecture. Journal of Artificial Intelligence Research (2007)
 - <u>Machine Vision</u> ···· Here we consider a new compositional model for finding salient curves.

Samples A* applications (cont')

Aker, A., Cohn, T., Gaizauskas, R.: Multi-document summarization using A* search and discriminative training. Proceedings of the 2010 Conference on Empirical Methods in Natural Language Processing.. ACL (2010)

Samples A* applications (cont')

EMNLP 2014 A* CCG Parsing with a Supertagfactored Model M. Lewis, M. Steedman

We introduce a new CCG parsing model which is factored on lexical category assignments. Parsing is then simply a deterministic search for the most probable category sequence that supports a CCG derivation. The parser is extremely simple, with a tiny feature set, no POS tagger, and no statistical model of the derivation or dependencies. Formulating the model in this way allows a highly effective heuristic for A* parsing, which makes parsing extremely fast. Compared to the standard C&C CCG parser, our model is more accurate out-of-domain, is four times faster, has higher coverage, and is greatly simplified. We also show that using our parser improves the performance of a state-of-the-art question answering system

Follow up ACL 2017 (main NLP conference – will be in Vancouver in August!)

A* CCG Parsing with a Supertag and Dependency Factored Model Masashi Yoshikawa Hiroshi Noji, Yuji Matsumoto_{Slide 25}

DFS, BFS, A* Animation Example

TheAI-Search animation system

http://www.cs.rmit.edu.au/AI-Search/Product/

DEPRECATED 🔗

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- To examine Search strategies when they are applied to the 8puzzle
- Compare only DFS, BFS andA* (with only the two heuristics we saw in class) AI-Search



nPuzzles are not always solvable

Half of the starting positions for the *n*-puzzle are impossible to solve (for more info on 8puzzle)

So experiment with the AI-Search animation system (DEPRECATED) with the default configurations.

If you want to try new ones keep in mind that you may pick unsolvable problems

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Learning Goals for today's class

- Define/read/write/trace/debug & Compare different search algorithms
 With / Without cost
 Informed / Uninformed
- Formally prove A* optimality.

Next class

