Online Convex Optimization

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Introduction

- Online learning: Make a sequence of accurate predictions given knowledge of the correct answer to previous prediction tasks and possibly additional available information.
- Applications: Online advertisement placement, web ranking, spam filtering, online shortest paths, portfolio selection, recommender systems

Notation

- \mathbf{X}_t Decision/point chosen on timestamp t
- \mathcal{K} Bounded convex, decision set
- $f_t \in \mathcal{F}: \mathcal{K} \mapsto \mathbb{R}$ Bounded convex function available at timestamp t
- T Number of iterations
- \mathcal{A} Online algorithm

Protocol

For t = 1:T

- 1. Learner chooses \mathbf{X}_t
- 2. Environment / Adversary chooses f_t
- 3. Learner suffers loss $f_t(\mathbf{x}_t)$

Aim: To minimize cumulative loss across rounds measured using regret

$$\operatorname{regret}_{T}(\mathcal{A}) = \sup_{\{f_{1}, \dots, f_{t}\} \subseteq \mathcal{F}} \left\{ \sum_{t=1}^{T} f_{t}(\mathbf{x}_{t}) - \min_{\mathbf{x} \in \mathcal{K}} \sum_{t=1}^{T} f_{t}(\mathbf{x}) \right\}$$

Chosen by our online algorithm Chosen by the offline algorithm

Learning from expert advice

- Need to make a decision (eg: whether to invest or not), helped by N experts who make predictions
- Decide which expert to follow at each step
- Receive feedback for each expert

Learning from expert advice

Algorithm 1 Hedge

- 1: Initialize: $\forall i \in [N], W_1(i) = 1$
- 2: for t = 1 to T do
- 3: Pick $i_t \sim_R W_t$, i.e., $i_t = i$ with probability $p_t(i) = \frac{W_t(i)}{\sum_i W_t(j)}$
- 4: Observe loss $\ell_t(i_t)$.
- 5: Update weights $W_{t+1}(i) = W_t(i)e^{-\varepsilon \ell_t(i)}$

6: end for

Theorem 1.5. Let ℓ_t^2 denote the *N*-dimensional vector of square losses, i.e., $\ell_t^2(i) = \ell_t(i)^2$, and let $\varepsilon < \frac{1}{2}$. The Hedge algorithm satisfies for any expert $i^* \in [N]$:

$$\sum_{t=1}^{T} p_t^{\top} \ell_t \le \sum_{t=1}^{T} f_t(i^*) + \varepsilon \sum_t p_t^{\top} \ell_t^2 + \frac{\ln N}{\varepsilon}$$

Asymptotic Regret Bounds

Theorem 3.2. Any algorithm for online convex optimization incurs $\Omega(DG\sqrt{T})$ regret in the worst case. This is true even if the cost functions are generated from a fixed stationary distribution.

 Achieved by first order methods like online gradient descent and second order methods like online newton

α -strongly convex	β -smooth
$\frac{1}{\alpha}\log T^{1}$	\sqrt{T} 2

Smoothness doesn't help (unlike in offline gradient descent), strong convexity does

Follow the leader

Idea: Use loss information in previous iterations to choose point in current iteration.

Follow-The-Leader (FTL)

$$\forall t, \mathbf{w}_t = \operatorname*{argmin}_{\mathbf{w} \in S} \sum_{i=1}^{t-1} f_i(\mathbf{w}) \quad \text{(break ties arbitrarily)}$$

Special case: Obtains log(T) regret with quadratic functions (shooting game) **Problem:** Predictions are not stable and may fluctuate drastically

Example 2.2 (Failure of FTL). Let $S = [-1, 1] \subset \mathbb{R}$ and consider the sequence of linear functions such that $f_t(w) = z_t w$ where

$$z_t = \begin{cases} -0.5 & \text{if } t = 1\\ 1 & \text{if } t \text{ is even}\\ -1 & \text{if } t > 1 \land t \text{ is odd} \end{cases}$$

Then, the predictions of FTL will be to set $w_t = 1$ for t odd and $w_t = -1$ for t even. The cumulative loss of the FTL algorithm will therefore be T while the cumulative loss of the fixed solution $u = 0 \in S$ is 0. Thus, the regret of FTL is T !

Follow the regularized leader

• **Idea:** Add regularization to make prediction more stable.

Algorithm 10 Regularized Follow The Leader

- 1: Input: $\eta > 0$, strongly convex regularization function R, and a convex compact set \mathcal{K} .
- 2: Let $\mathbf{x}_1 = \arg\min_{\mathbf{x}\in\mathcal{K}} \{R(\mathbf{x})\}.$
- 3: for t = 1 to T do
- 4: Predict \mathbf{x}_t .
- 5: Observe the payoff function f_t and let $\nabla_t = \nabla f_t(\mathbf{x}_t)$.
- 6: Update

$$\mathbf{x}_{t+1} = \underset{\mathbf{x}\in\mathcal{K}}{\operatorname{arg\,min}} \underbrace{\left\{ \eta \sum_{s=1}^{t} \nabla_s^{\top} \mathbf{x} + R(\mathbf{x}) \right\}}_{\Phi_t(\mathbf{x})}$$
(5.1)

7: end for

Follow the regularized leader

Theorem 5.1. The RFTL algorithm 10 attains for every $\mathbf{u} \in \mathcal{K}$ the following bound on the regret:



Special Cases:

- Online Gradient Descent with L2 regularization i.e. $R(\mathbf{x}) = \frac{1}{2} \|\mathbf{x}\|_2^2$
- Hedge with negative entropy regularization i.e. $R(\mathbf{x}) = \mathbf{x} \log \mathbf{x}$

Online gradient descent

Algorithm 6 Online Gradient Descent

- 1: Input: convex set $\mathcal{K}, T, \mathbf{x}_1 \in \mathcal{K}$, step sizes $\{\eta_t\}$
- 2: for t = 1 to T do
- 3: Play \mathbf{x}_t and observe cost $f_t(\mathbf{x}_t)$.
- 4: Update and project:



Online gradient descent

Theorem 3.1. ONLINE GRADIENT DESCENT with step sizes $\{\eta_t = \frac{D}{G\sqrt{t}}, t = 1, ..., T\}$ guarantees the following for all $T \ge 1$.

$$\operatorname{regret}_{T} = \sum_{t=1}^{T} f_{t}(\mathbf{x}_{t}) - \min_{\mathbf{x}^{*} \in \mathcal{K}} \sum_{t=1}^{T} f_{t}(\mathbf{x}^{*}) \leq 3GD\sqrt{T}$$

Theorem 3.3. For α -strongly convex loss functions, ONLINE GRADI-ENT DESCENT with step sizes $\eta_t = \frac{1}{\alpha t}$ achieves the following guarantee, for all $T \ge 1$

$$\operatorname{regret}_T \leq \frac{G^2}{2\alpha} (1 + \log T).$$

Online to Batch reduction

Suppose we have a batch of data, can we OGD to look one point at a time and get a rate as good as SGD ?

Algorithm 7 Stochastic Gradient Descent

- 1: Input: $f, \mathcal{K}, T, \mathbf{x}_1 \in \mathcal{K}$, step sizes $\{\eta_t\}$
- 2: for t = 1 to T do
- 3: Let $\tilde{\nabla}_t \leftarrow \mathcal{O}(\mathbf{x}_t)$ and define: $f_t(x) \triangleq \langle \tilde{\nabla}_t, x \rangle$
- 4: Update and project:

$$\mathbf{y}_{t+1} = \mathbf{x}_t - \eta_t \tilde{\nabla}_t$$

$$\mathbf{x}_{t+1} = \prod_{\mathcal{K}} [\mathbf{y}_{t+1}]$$

- 5: end for
- 6: return $\bar{\mathbf{x}}_T \triangleq \frac{1}{T} \sum_{t=1}^t \mathbf{x}_t$

Use linearized functions with noisy gradient as input and return the average point over the iterations

Online to Batch reduction

Regret bound:

$$\mathbf{E}[f(\bar{\mathbf{x}}_T)] \le \min_{\mathbf{x}^* \in \mathcal{K}} f(\mathbf{x}^*) + \frac{3GD}{\sqrt{T}}$$

- For general convex function, achieves 1/sqrt(T) rate
- For strongly convex functions, online to batch conversion can give a log(T)/T rate later improved to 1/T (Hazan'14)

ADAGRAD

 Different regularizations in RTFL lead to different algorithms. Can we learn an optimal regularization to use ?

Algorithm 16 AdaGrad

1: Input: parameters $\eta, \delta > 0, \mathbf{x}_1 \in \mathcal{K}$.

2: Initialize:
$$\mathbf{S}_0 = \mathbf{G}_0 = \delta I$$
,

- 3: for t = 1 to T do
- 4: Predict \mathbf{x}_t , suffer loss $f_t(\mathbf{x}_t)$.
- 5: Update: $\mathbf{S}_t \leftarrow \mathbf{S}_{t-1} + \nabla_t \nabla_t^\top, \ \mathbf{G}_t = \mathbf{S}_t^{1/2}$

$$\mathbf{y}_{t+1} \leftarrow \mathbf{x}_t - \eta \mathbf{G}_t^{-1} \nabla_t$$

$$\mathbf{x}_{t+1} \leftarrow \arg\min_{\mathbf{x}\in\mathcal{K}} \|\mathbf{y}_{t+1} - \mathbf{x}\|_{\mathbf{G}_t}^2$$

6: **end for**

ADAGRAD

 $\forall \mathbf{x} \in \mathcal{K} \ . \ \nabla^2 R(\mathbf{x}) = \nabla^2 \in \mathcal{H} \triangleq \{ X \in \mathbb{R}^{d \times d} \ , \ \mathbf{Tr}(X) \le 1 \ , \ X \succeq 0 \}$

Theorem 5.9. Let $\{\mathbf{x}_t\}$ be defined by Algorithm 16 with parameters $\delta = 1, \eta = \frac{1}{D}$, where

 $D = \max_{\mathbf{u} \in \mathcal{K}} \max_{H \in \mathcal{H}} \|\mathbf{u} - \mathbf{x}_1\|_H$

Then for for any $\mathbf{x}^* \in \mathcal{K}$,

$$\operatorname{regret}_{T}(\operatorname{AdaGrad}) \leq 2D \cdot \sqrt{\min_{H \in \mathcal{H}} \sum_{t} \|\nabla_{t}\|_{H}^{*2}}$$
(5.6)

 $\operatorname{regret}(\operatorname{AdaGrad}) \leq 2D\mathbf{Tr}(\mathbf{G}_T)$

Variants

- Follow the perturbed leader (FPL) form of randomized regularization
- Bandit feedback (observe loss only for the selected point) – can be analyzed using OCO by estimating the gradient at each point
- Use Frank Wolfe when projection is expensive