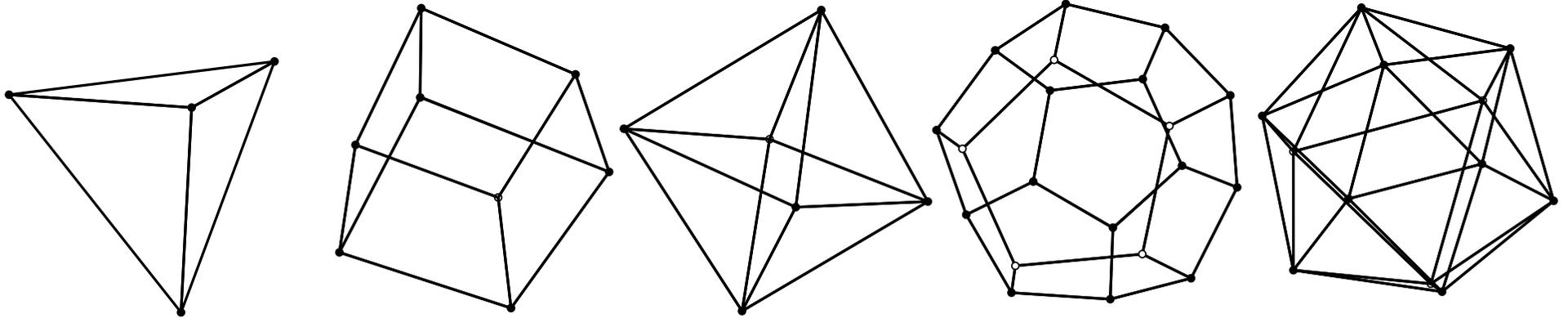
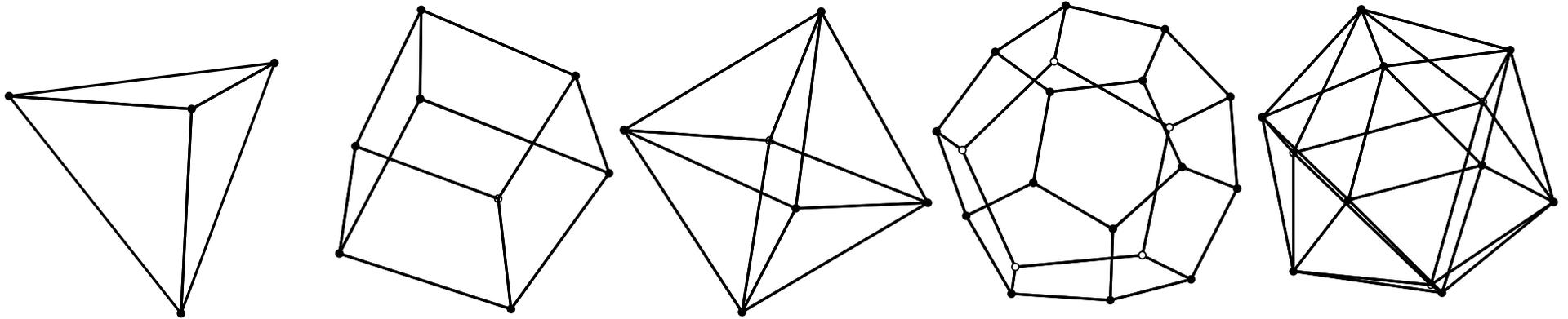


Euler's formula $n - e + f = 2$ and Platonic solids

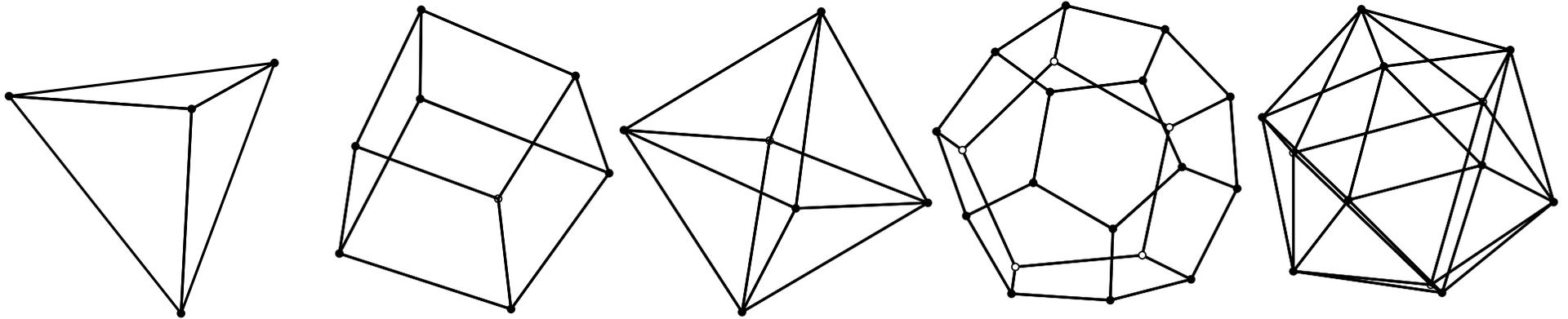


Euler's formula $n - e + f = 2$ and Platonic solids



spherical projection makes these planar graphs

Euler's formula $n - e + f = 2$ and Platonic solids



spherical projection makes these planar graphs

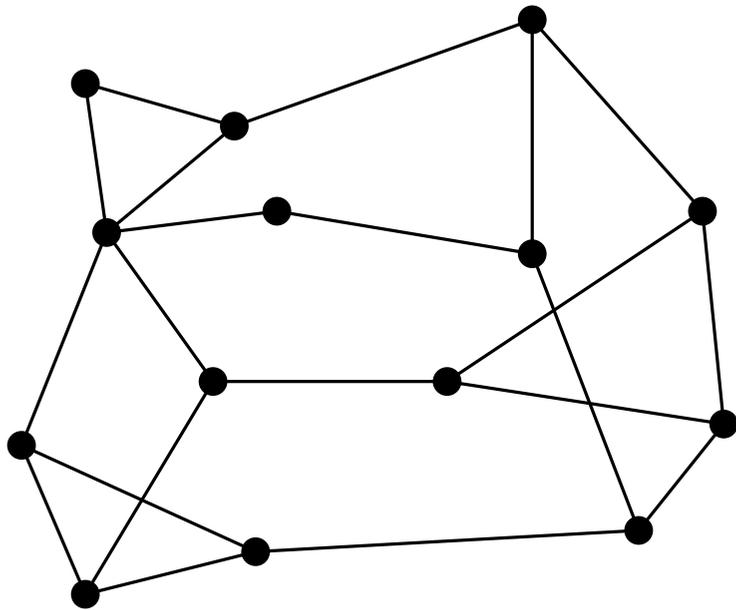
$p = \text{vertex degree} \geq 3, q = \text{face degree} \geq 3$

$pn = 2e \quad qf = 2e \quad \text{Euler implies } \frac{2e}{p} - e + \frac{2e}{q} = 2$

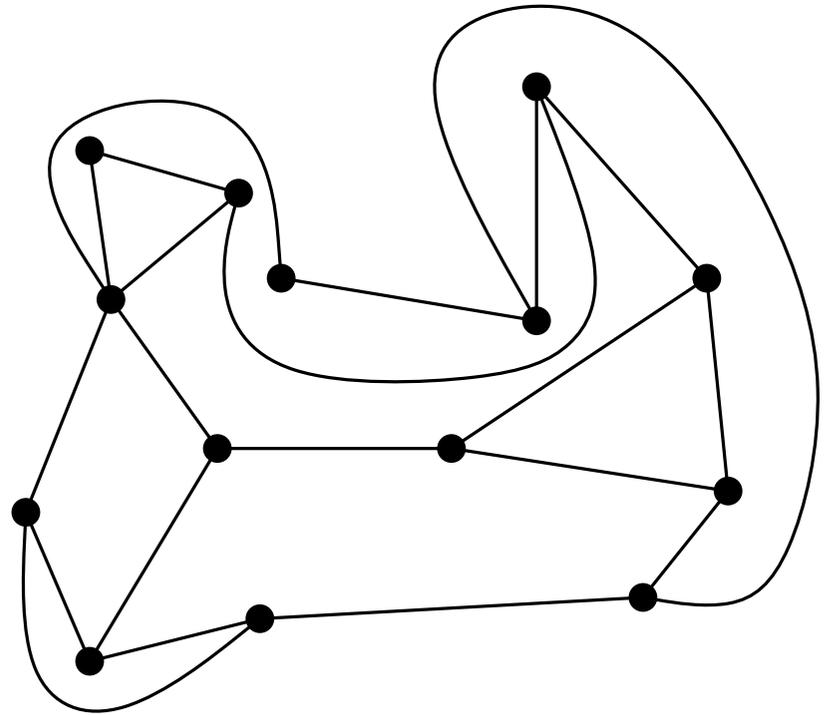
so $\frac{2}{p} + \frac{2}{q} > 1$

Thus $(p, q) \in \{(3, 3), (3, 4), (4, 3), (3, 5), (5, 3)\}$

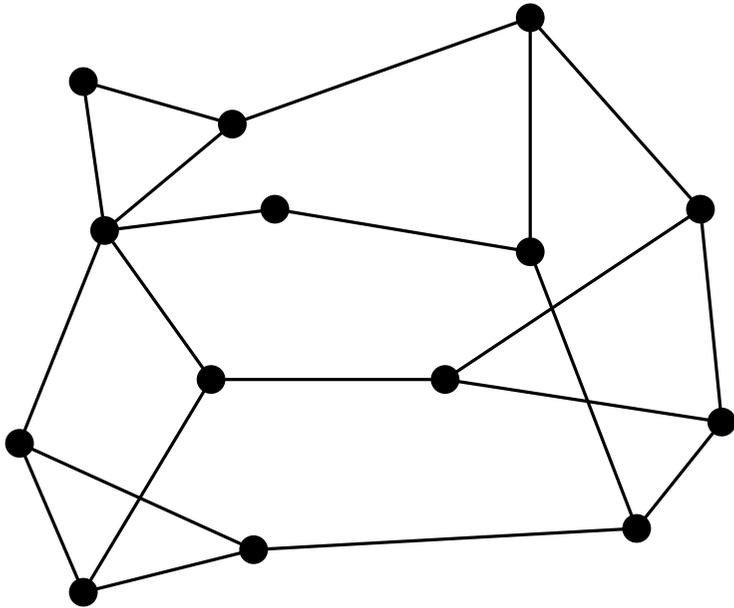
Planar graph



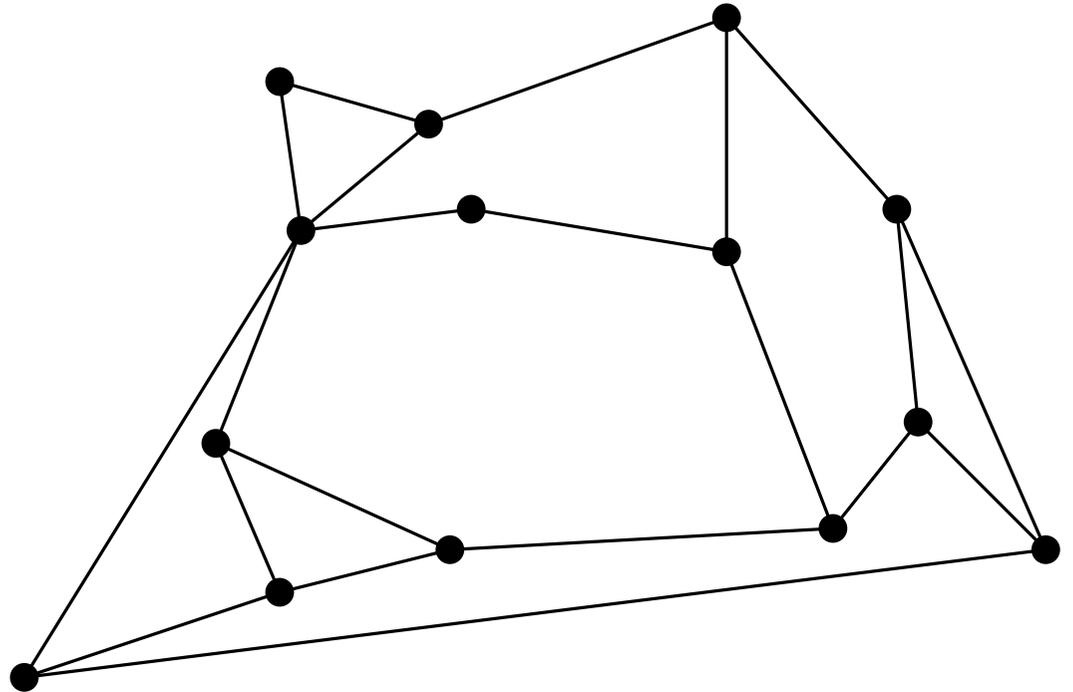
Planar drawing



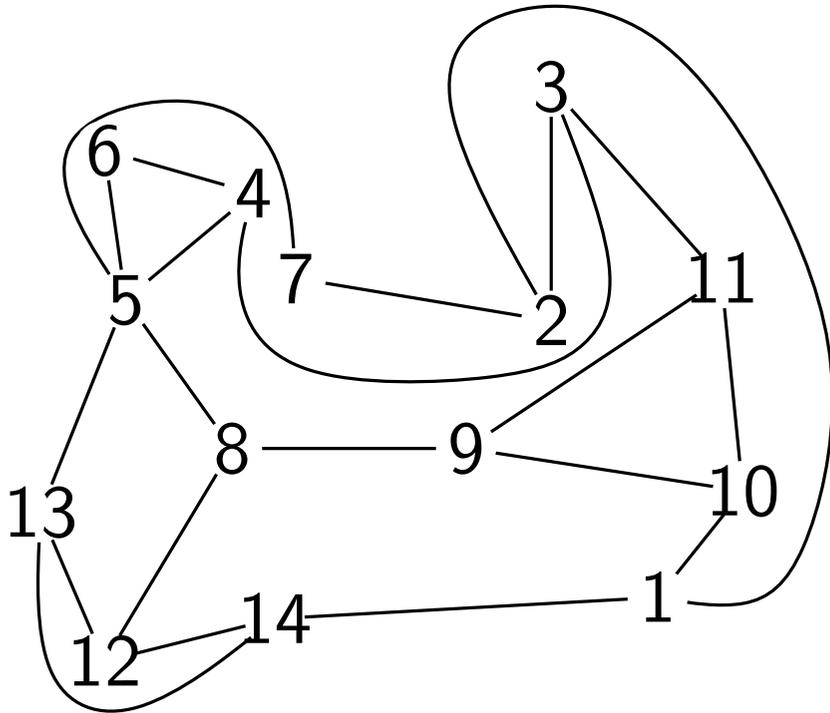
Planar graph



Planar (straight-line) drawing



Planar drawing



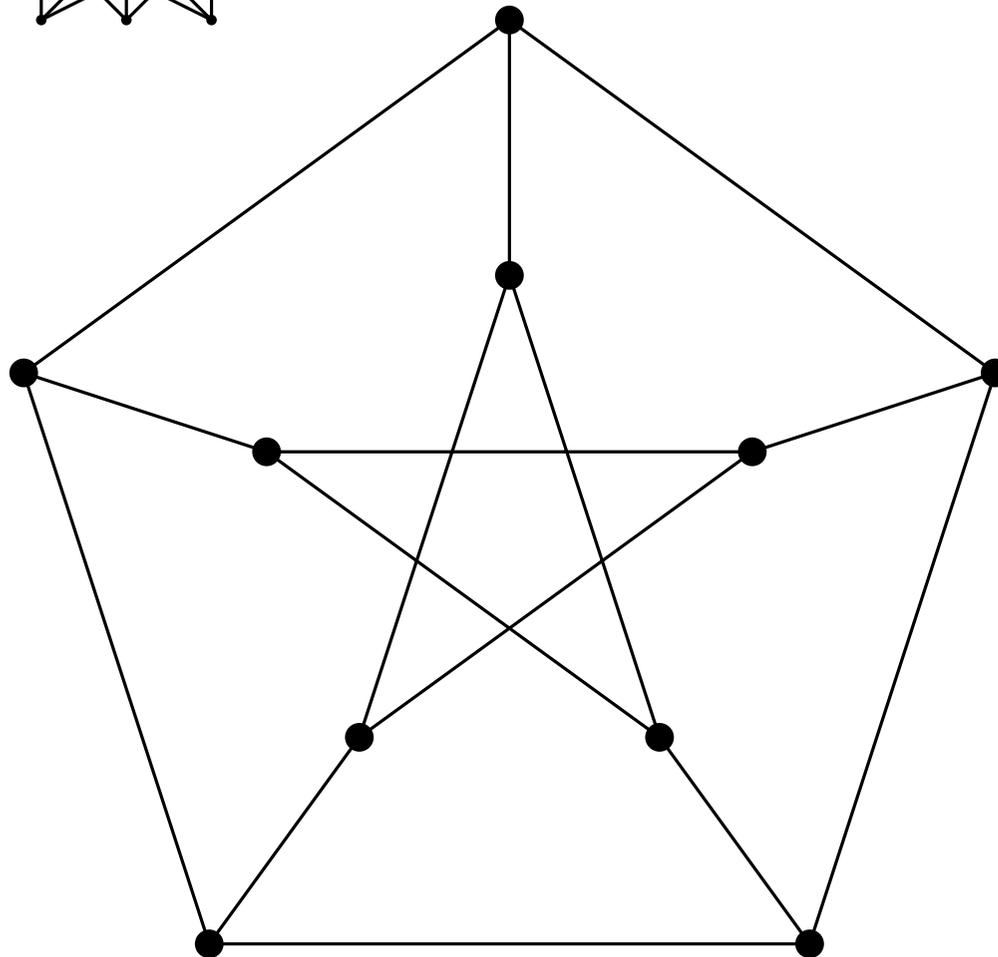
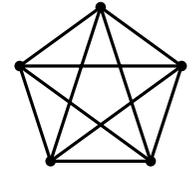
Planar Embedding

- 1: 2,14,10
- 2: 1,3,7
- 3: 2,11,4
- 4: 3,5,6
- 5: 4,8,13,7,6
- 6: 4,5
- 7: 2,5
- 8: 5,9,12
- 9: 8,11,10
- 10: 1,9,11
- 11: 3,10,9
- 12: 8,14,13
- 13: 5,12,14
- 14: 1,13,12

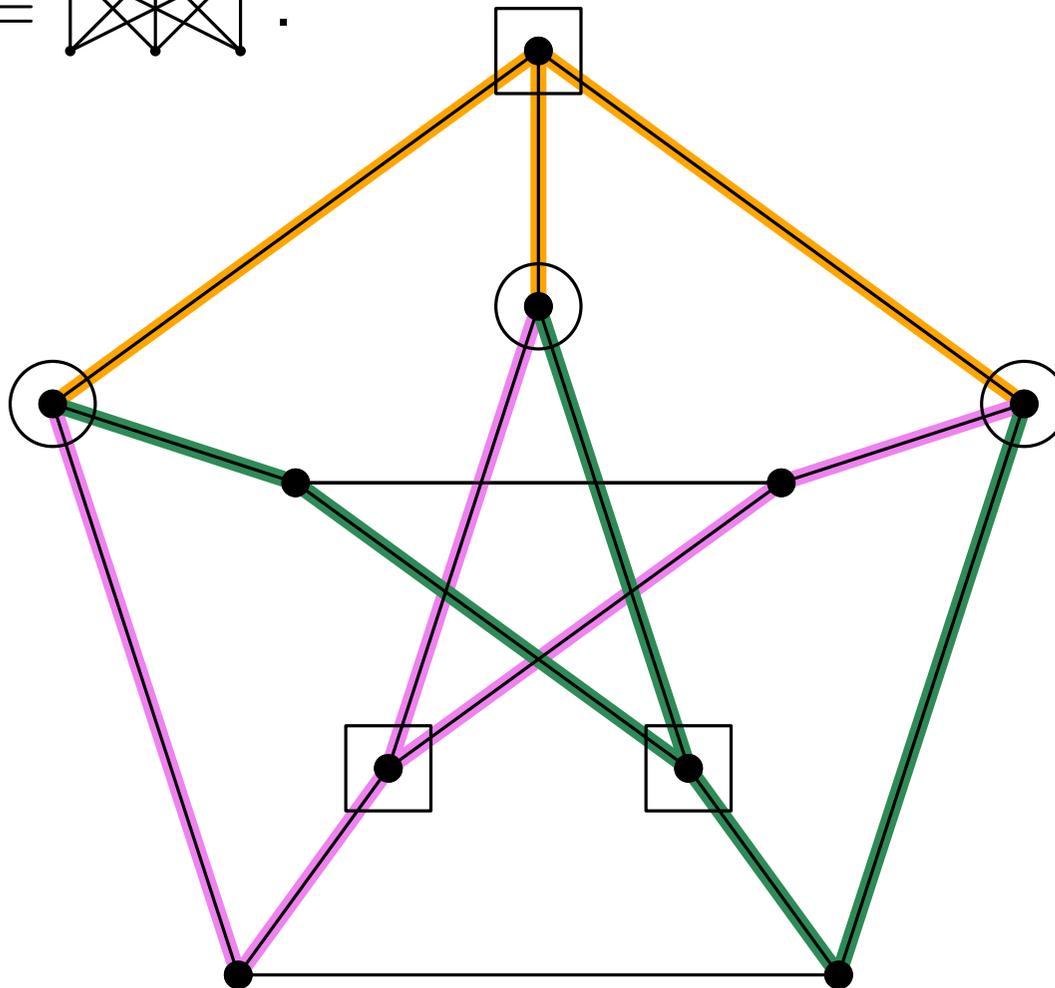
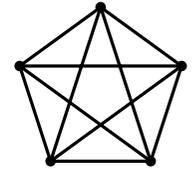
Outer face: 1,14,13,5,7,2

(Clockwise order of neighbors.)

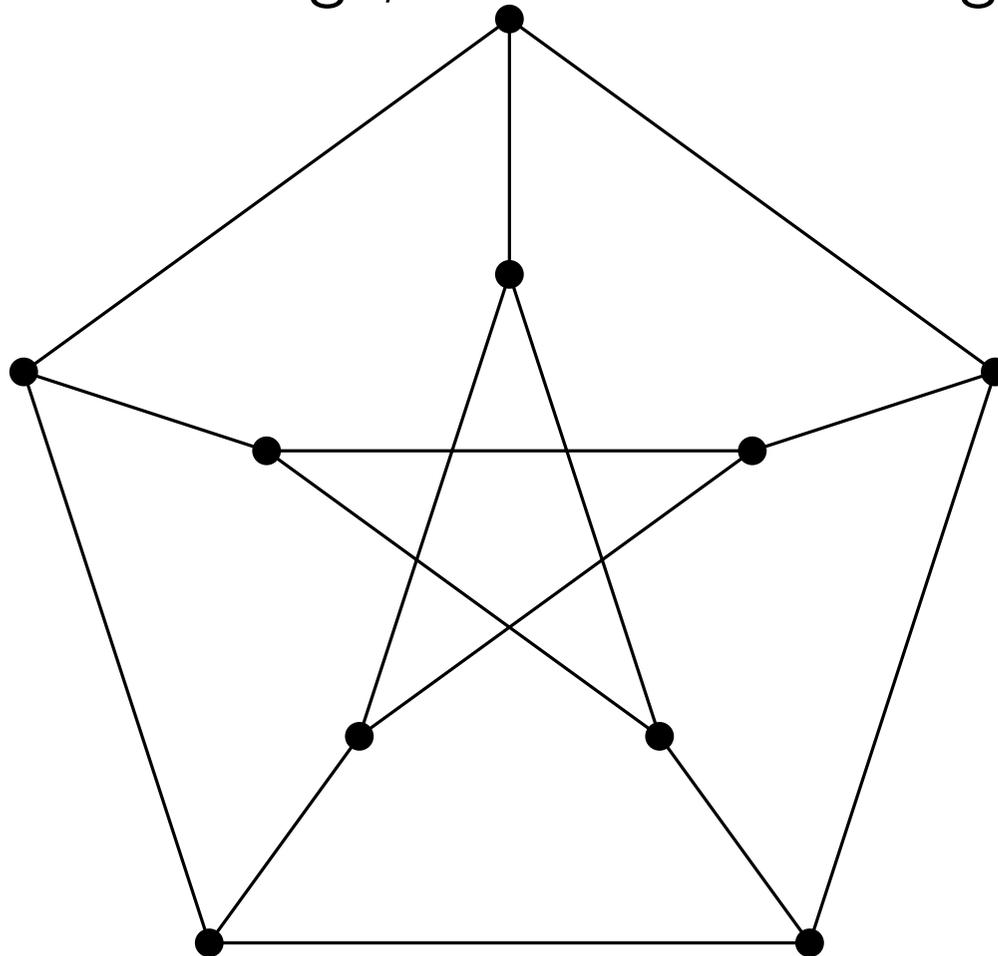
Kuratowski's Theorem G is planar iff it doesn't contain a subgraph that is a subdivision of $K_5 \equiv$ or $K_{3,3} \equiv$.



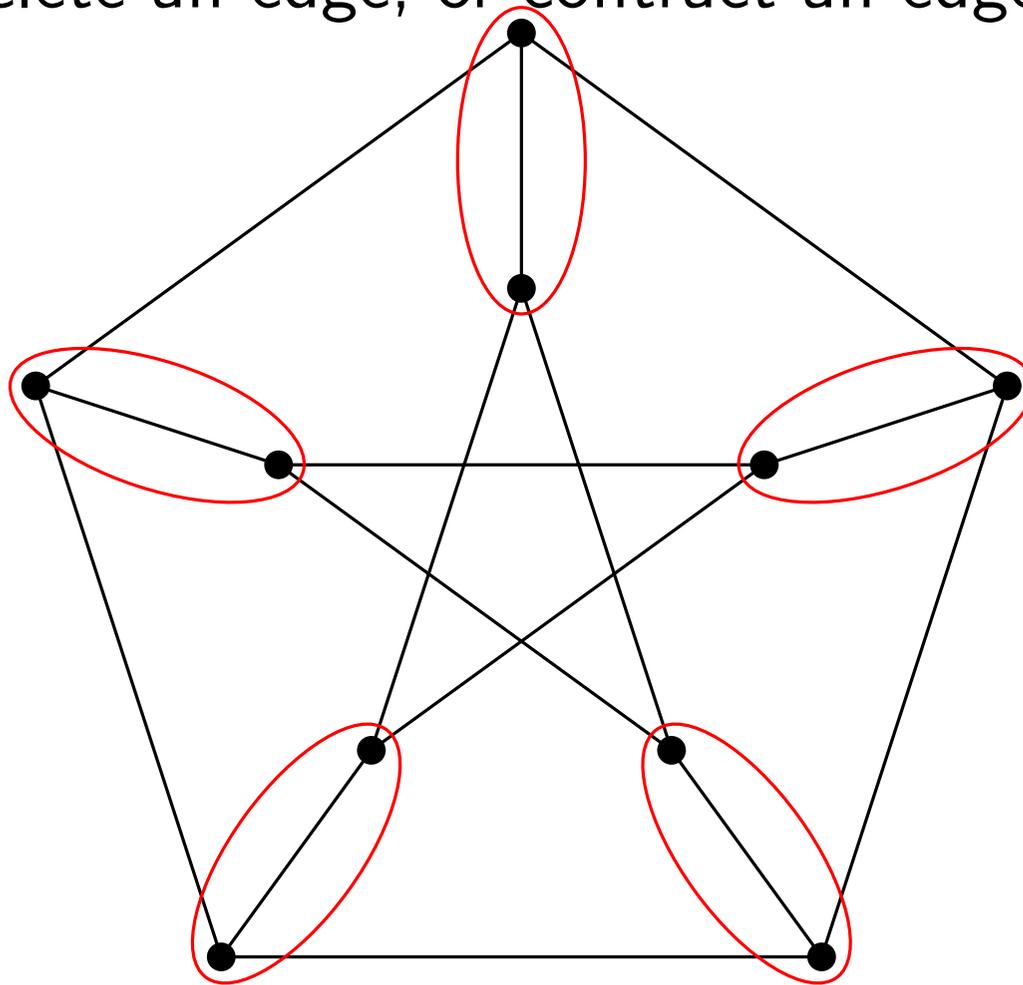
Kuratowski's Theorem G is planar iff it doesn't contain a subgraph that is a subdivision of $K_5 \equiv$ or $K_{3,3} \equiv$.



Wagner's Theorem G is planar iff it doesn't have K_5 or $K_{3,3}$ as a *minor* (to get a minor, repeatedly delete a vertex, delete an edge, or contract an edge).

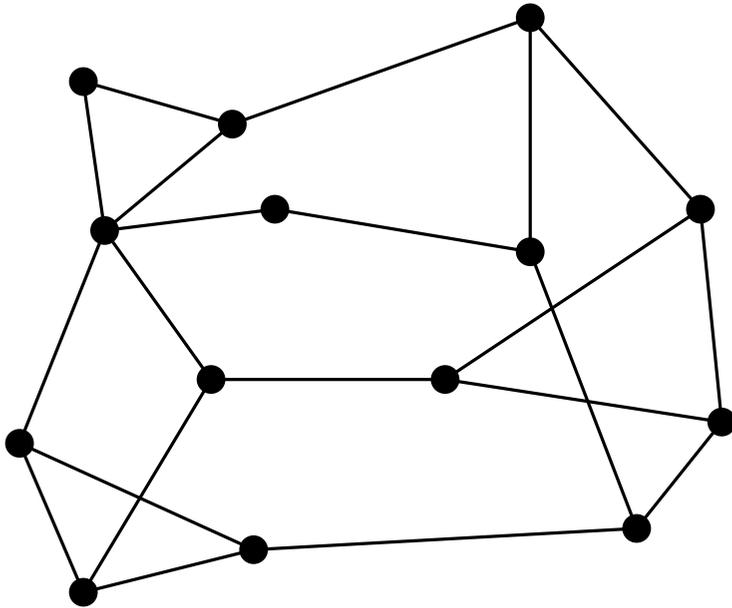


Wagner's Theorem G is planar iff it doesn't have K_5 or $K_{3,3}$ as a *minor* (to get a minor, repeatedly delete a vertex, delete an edge, or contract an edge).

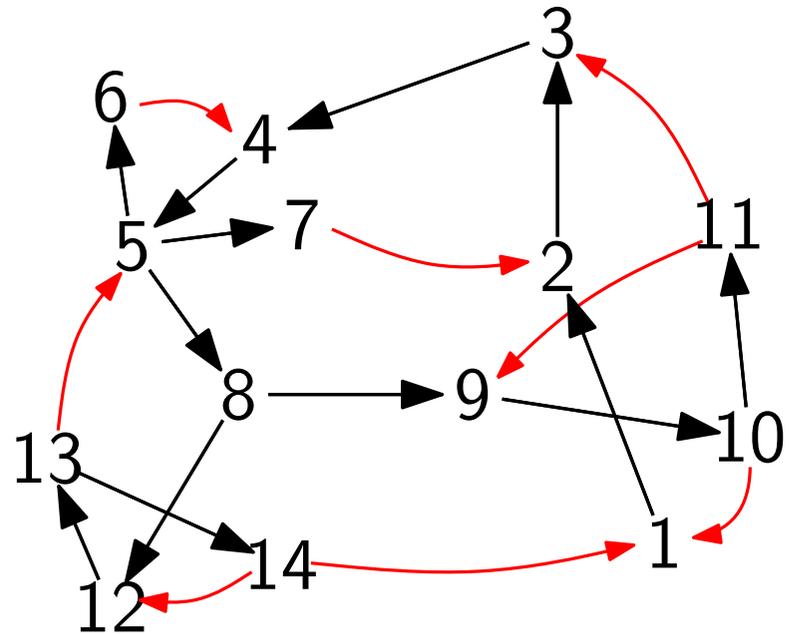


Kuratowski's and Wagner's conditions do not imply a fast algorithm to check planarity, but there are linear time planarity testing algorithms,

Planar graph

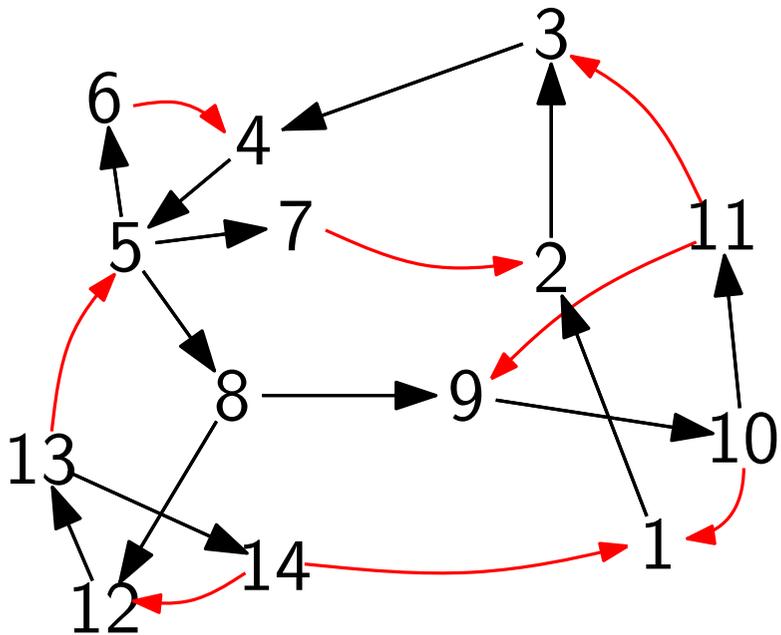


DFS orientation & numbering

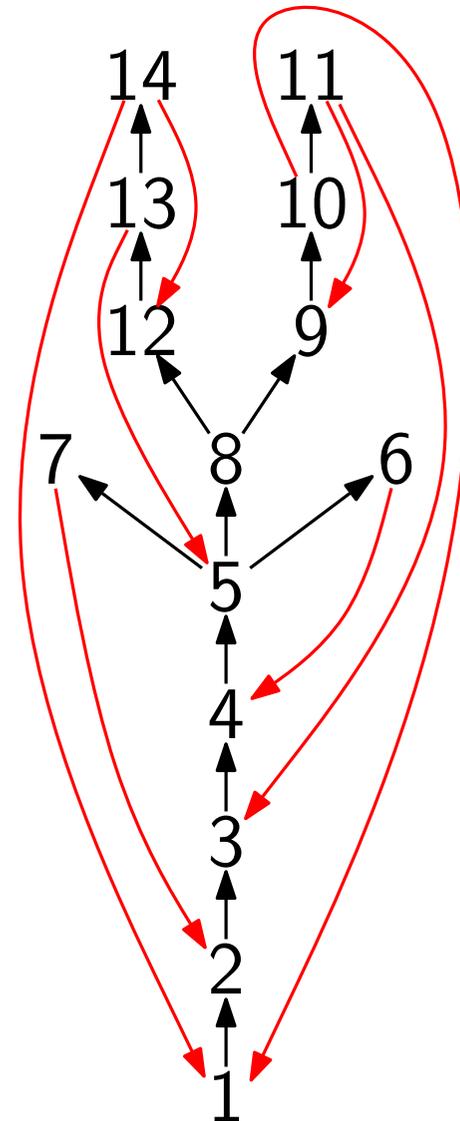


→ tree edge
↪ back edge

DFS orientation & numbering

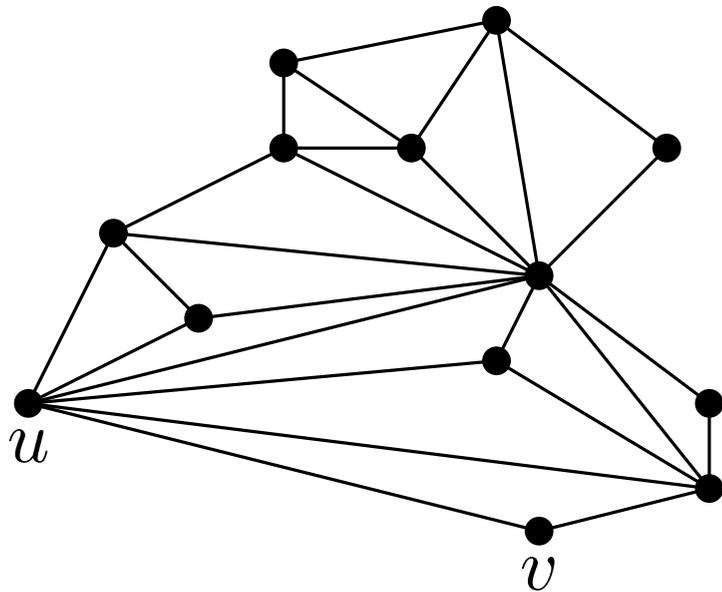


LR partition [Brandes 11]



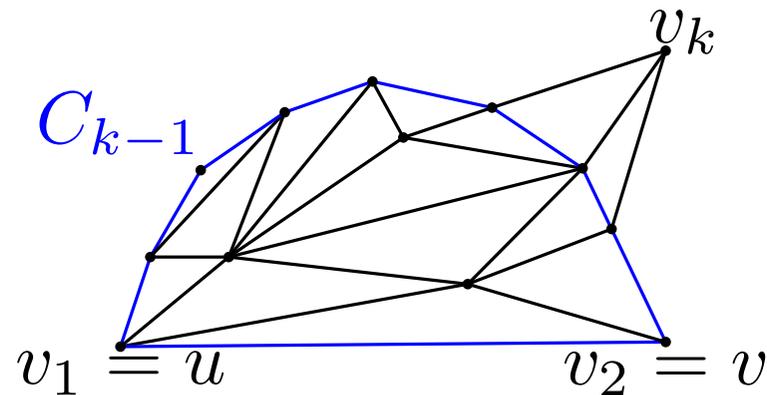
Drawing a Planar Graph on a Grid [de Fraysseix, Pach, Pollack 90]

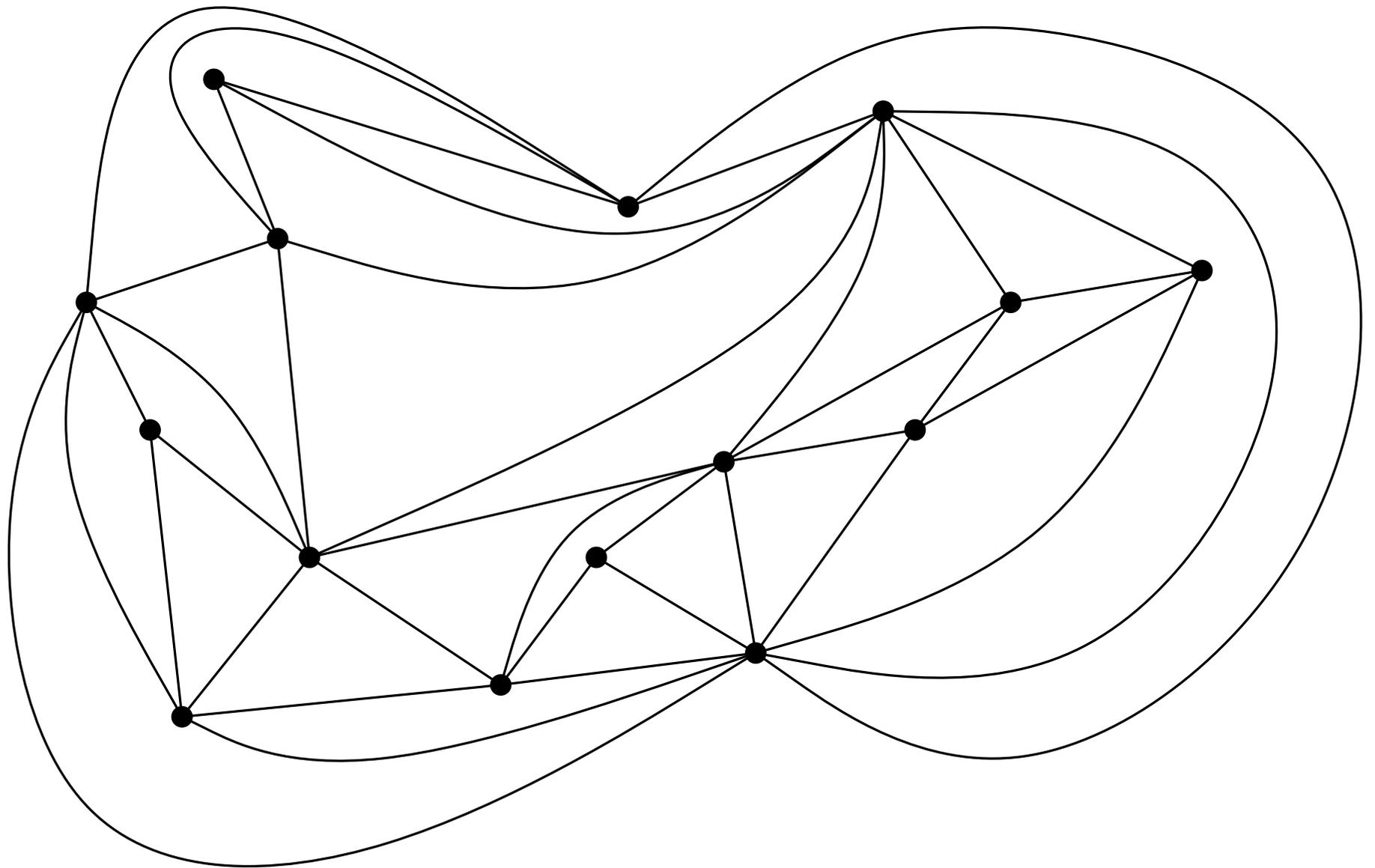
Lemma. Let G be a simple planar graph with an embedding and $u = u_1, u_2, \dots, u_k = v$ be a cycle of G . Then there exists a vertex w on the cycle, different from u and v and not adjacent to any inside chord.

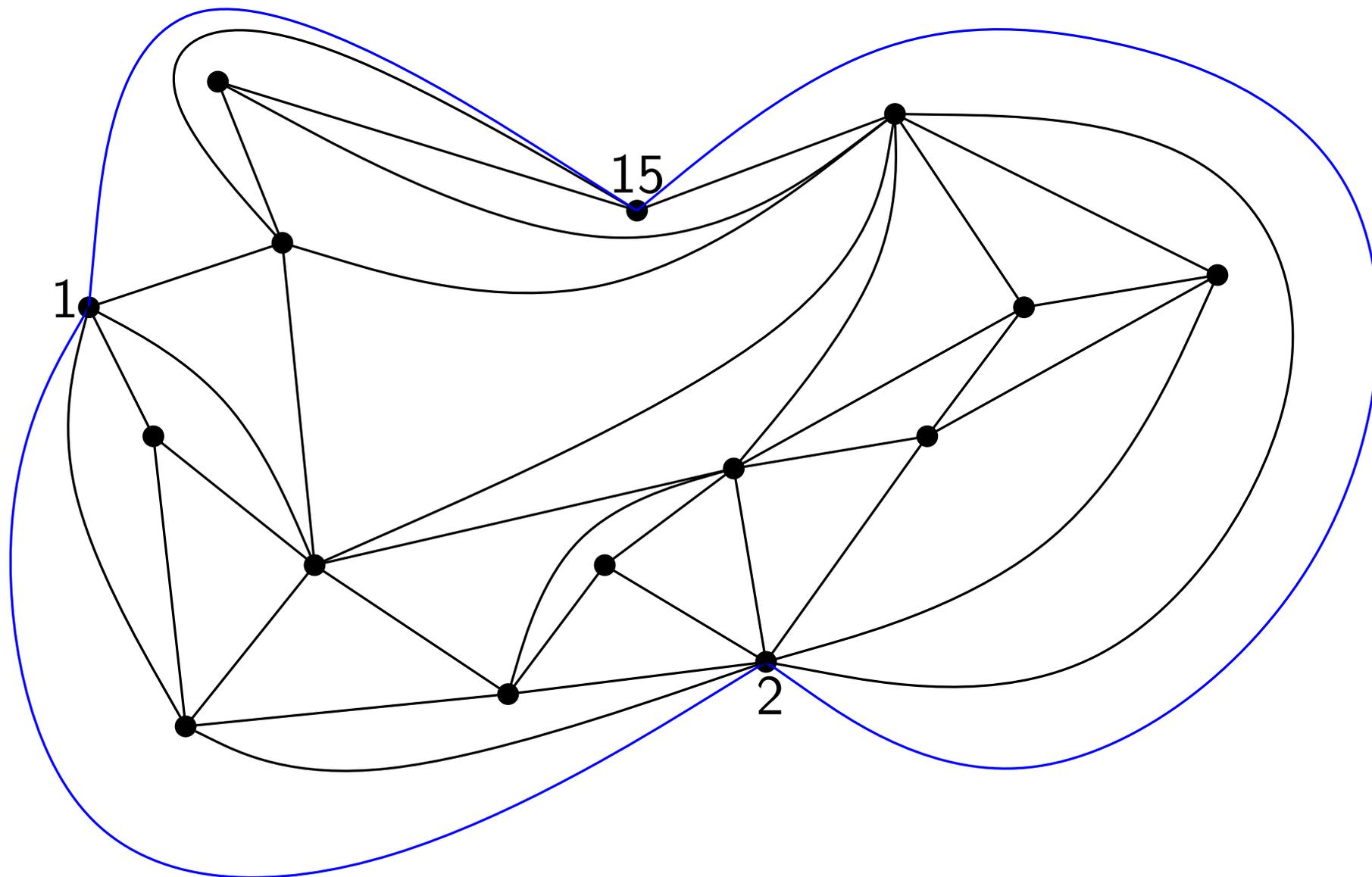


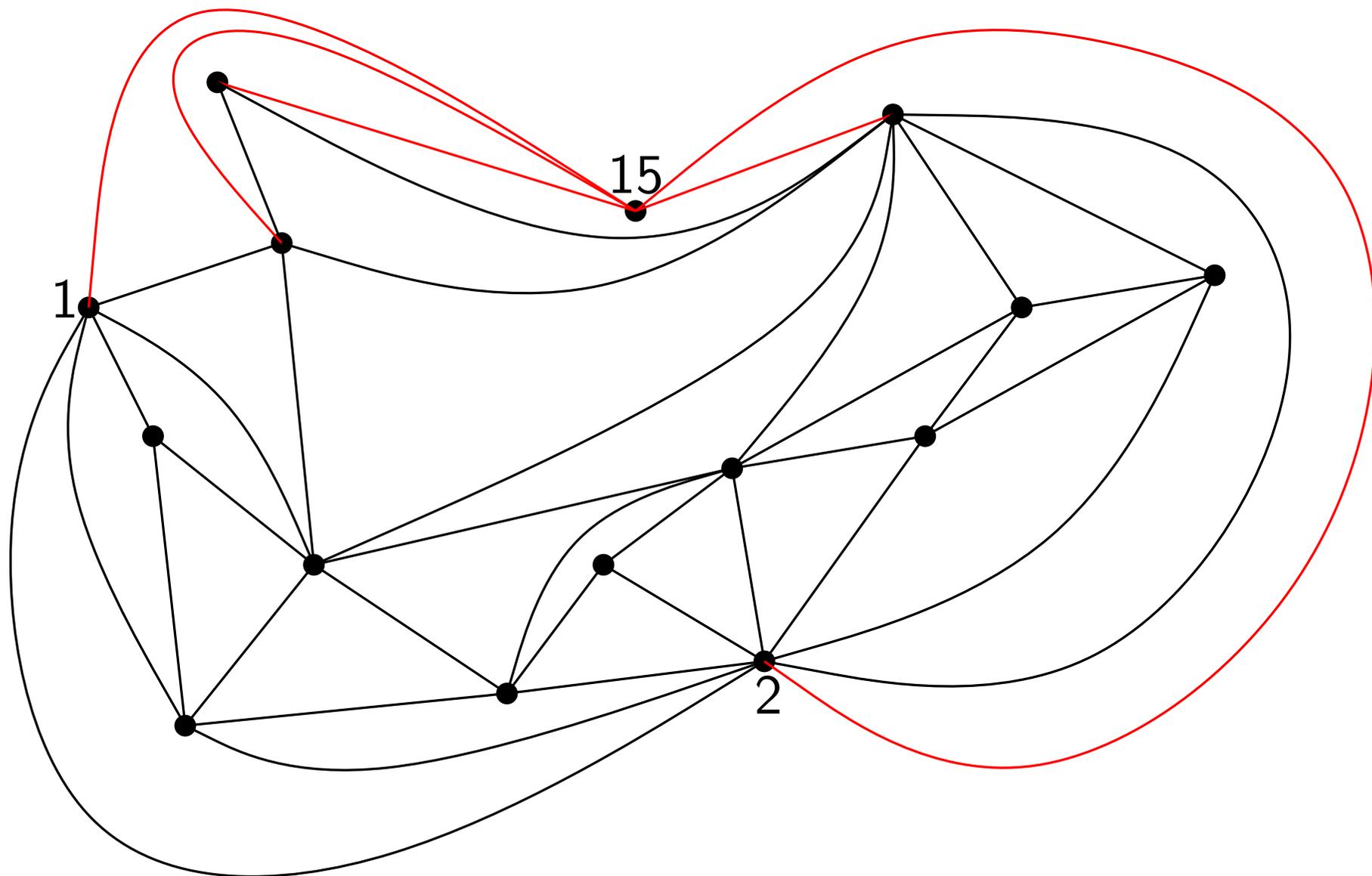
Canonical Ordering. If $G = (V, E, \mathcal{E})$ is a **maximal** plane graph with exterior face u, v, w then there exists an ordering of V , $v_1 = u, v_2 = v, v_3, \dots, v_n = w$ such that for all $4 \leq k \leq n$:

- (i) The subgraph G_{k-1} induced by v_1, v_2, \dots, v_{k-1} is **2-connected** with uv on its exterior cycle C_{k-1} ;
- (ii) v_k is in the exterior face of G_{k-1} , and its neighbors in G_{k-1} are a length ≥ 1 subpath of the path $C_{k-1} - uv$.

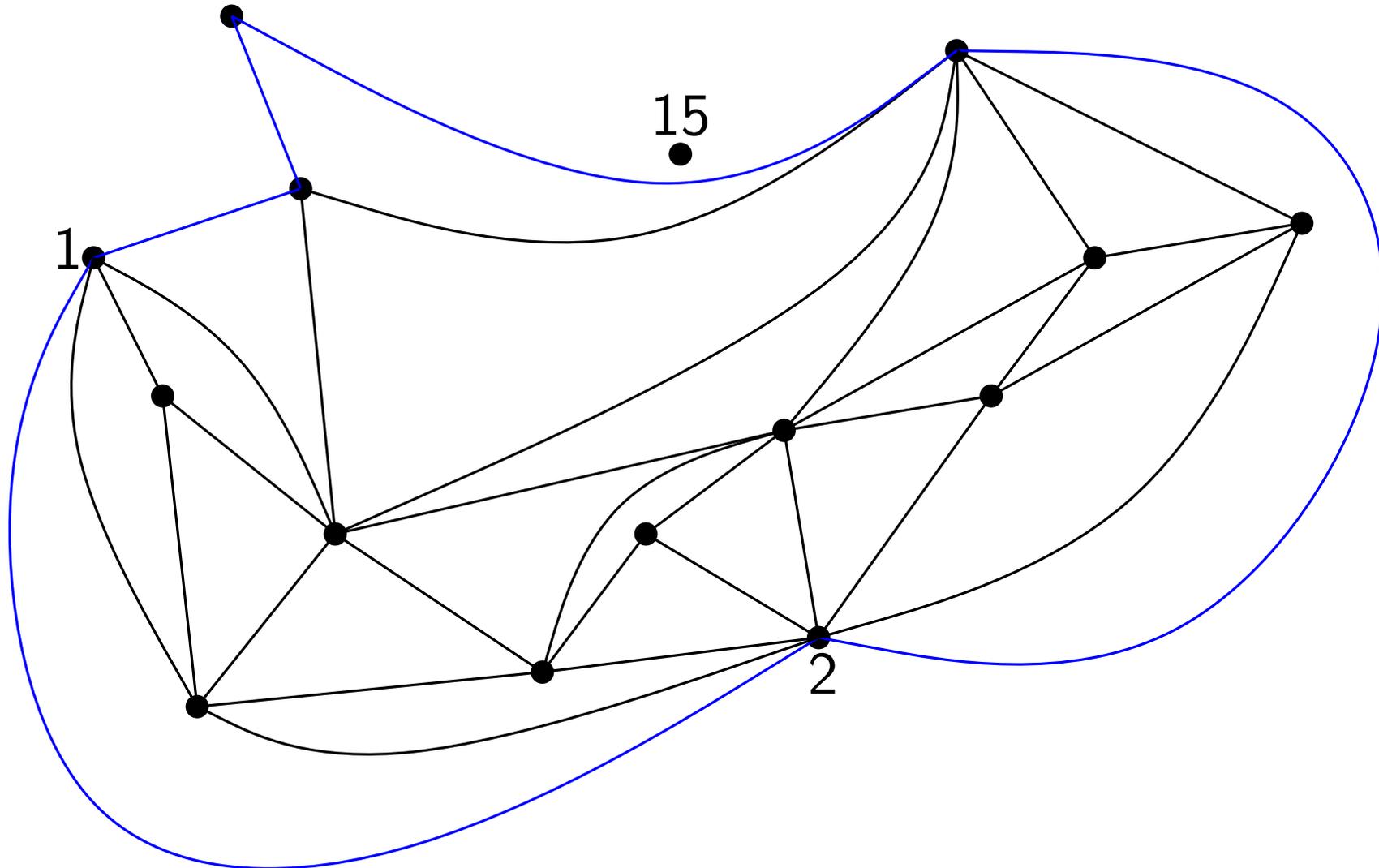




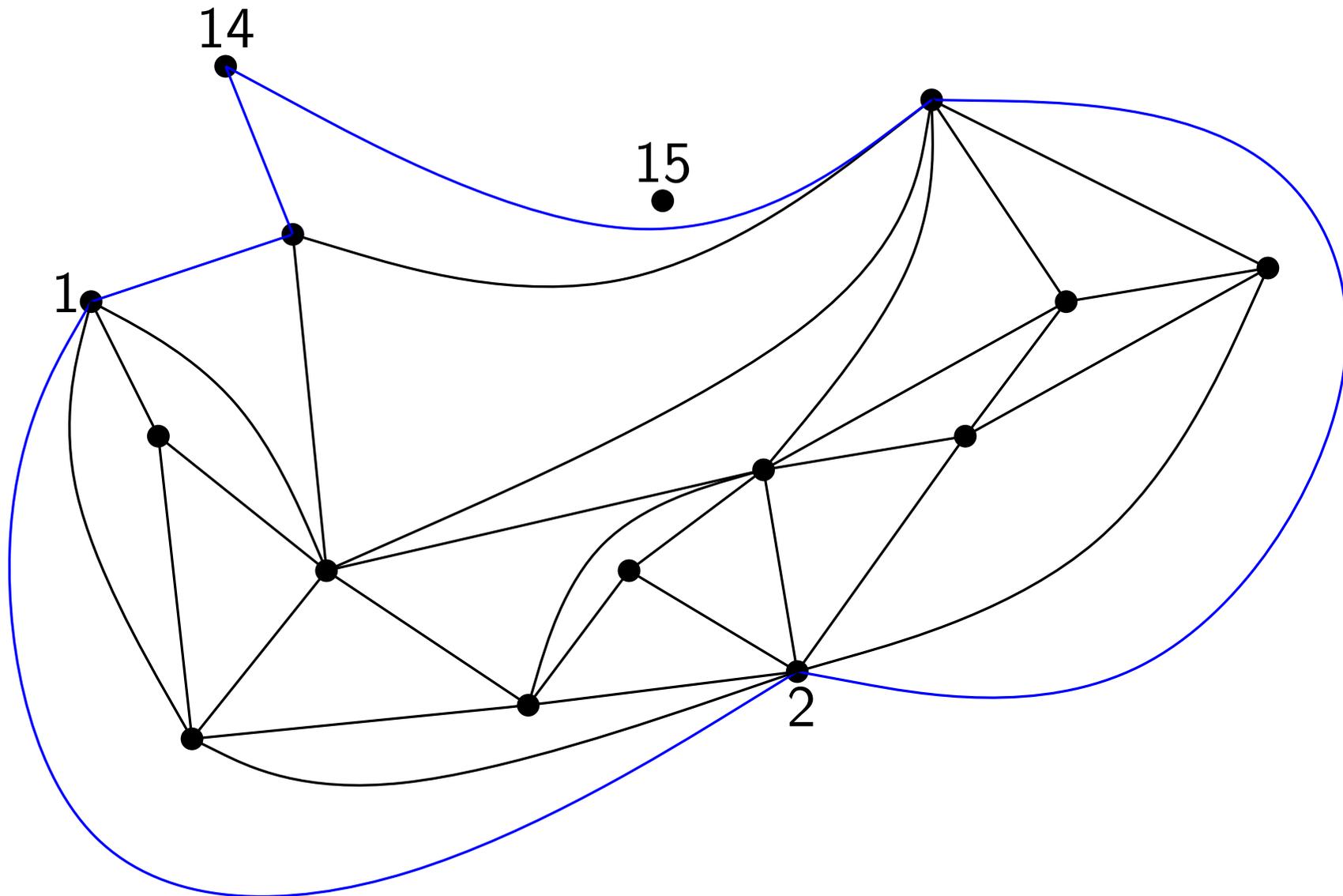




Use Lemma to find next vertex.



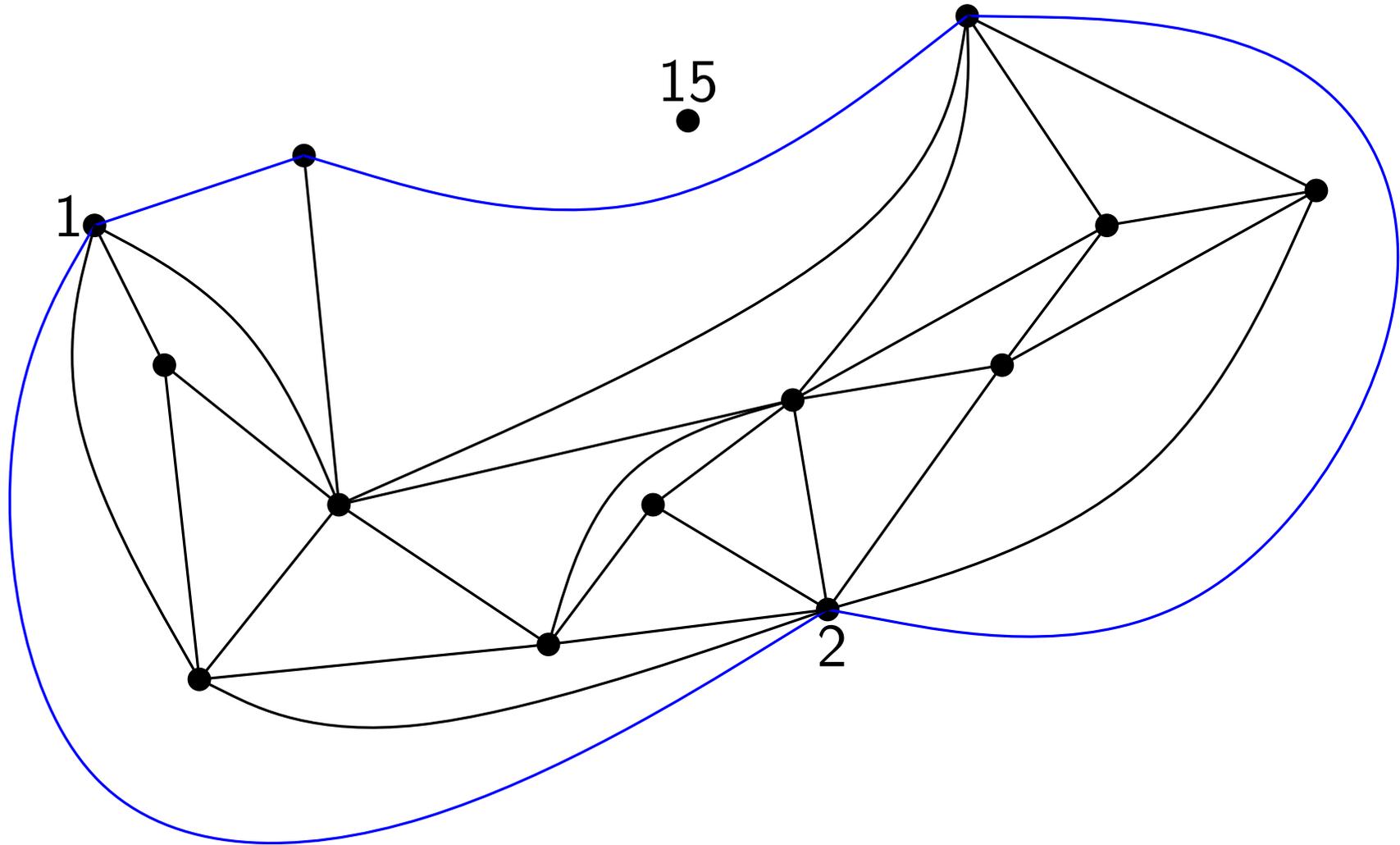
Use Lemma to find next vertex.

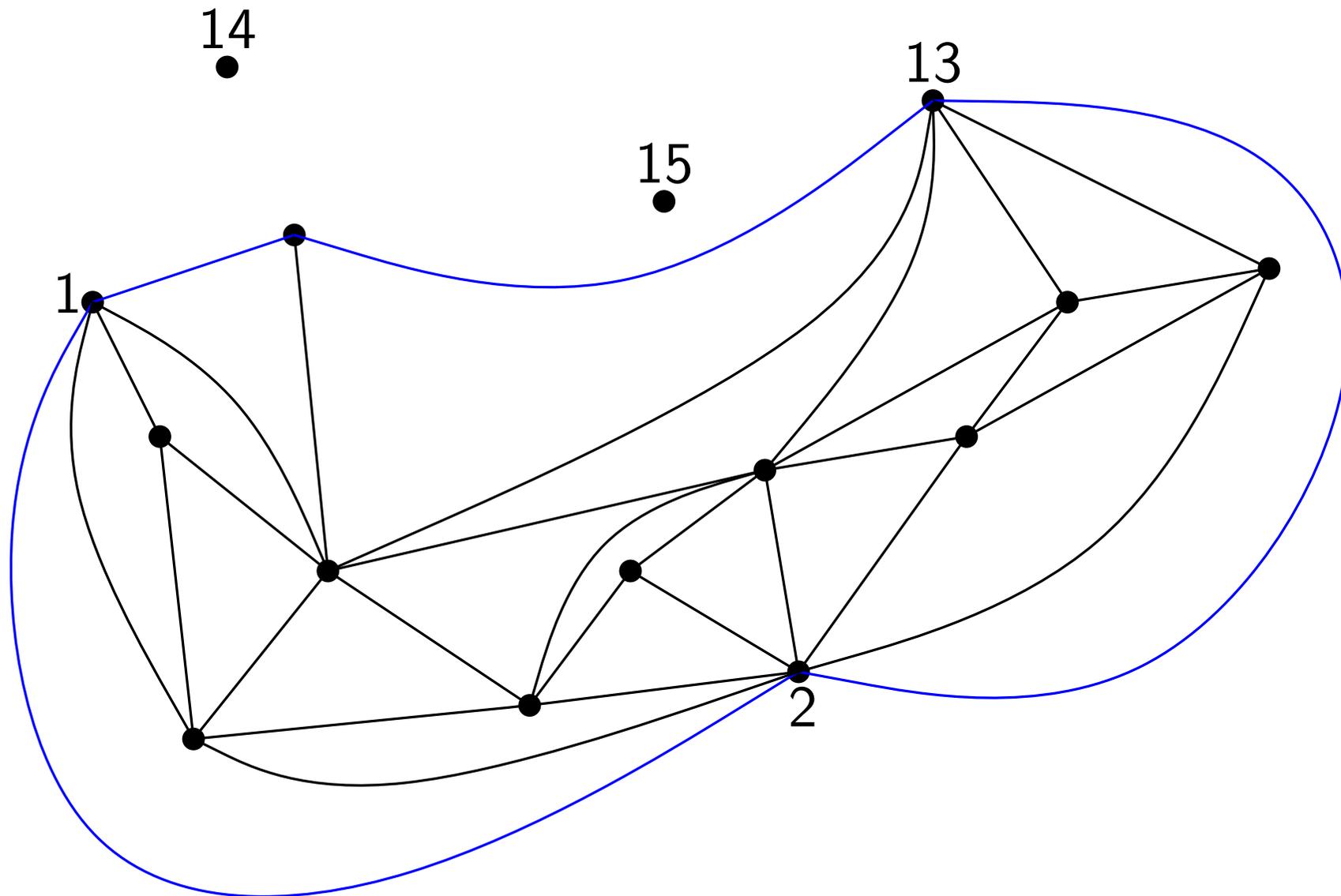


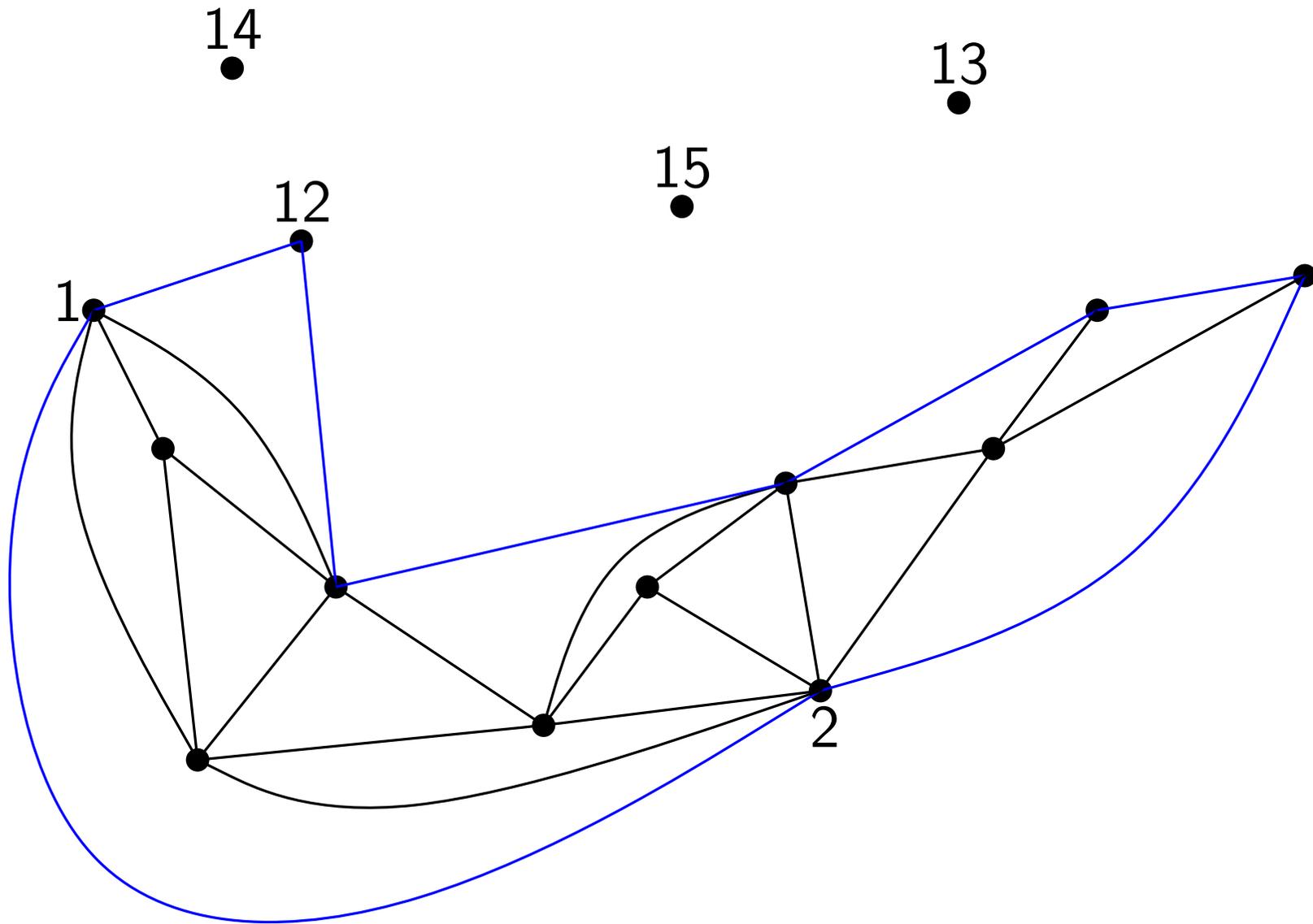
Repeat

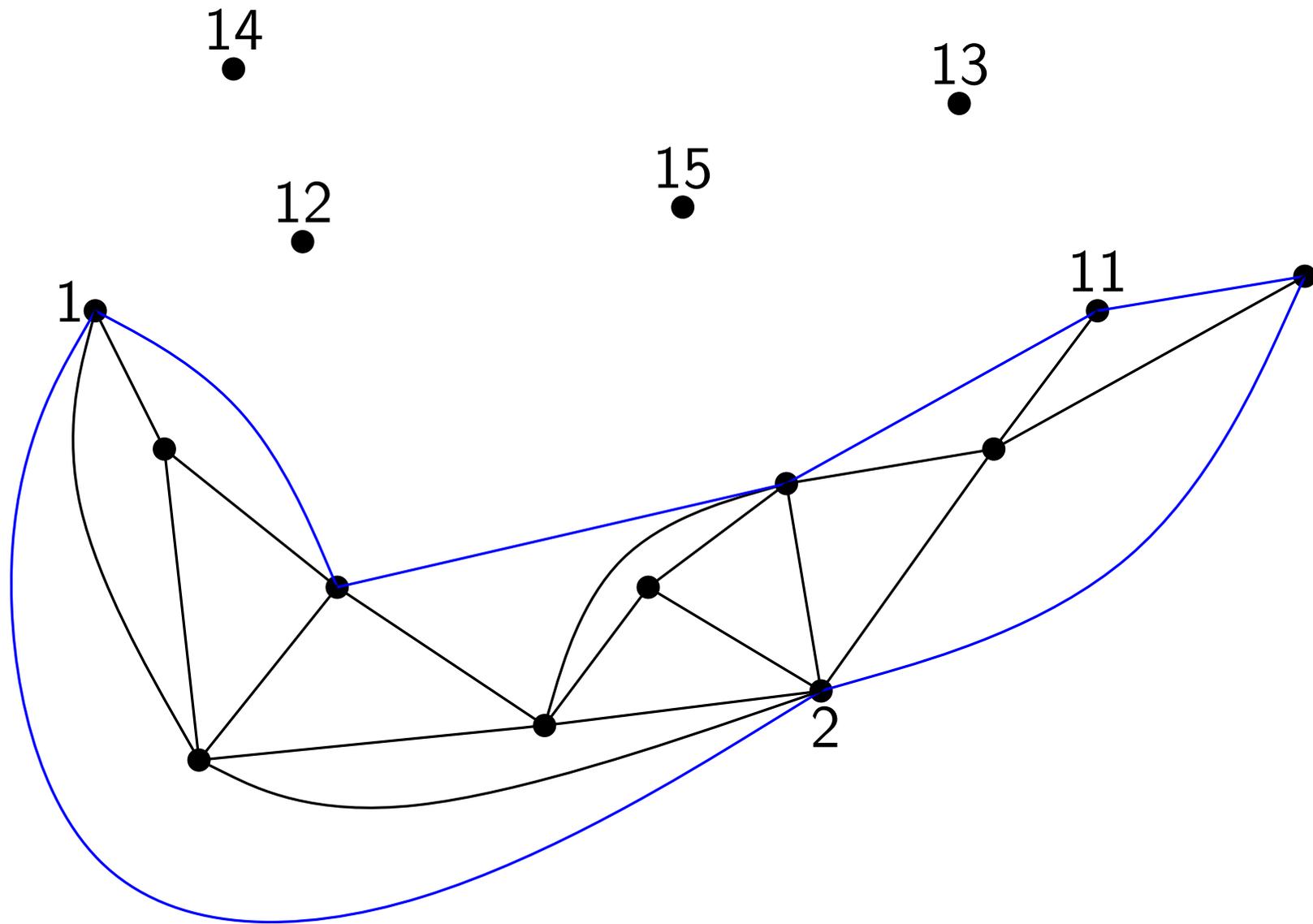
14
●

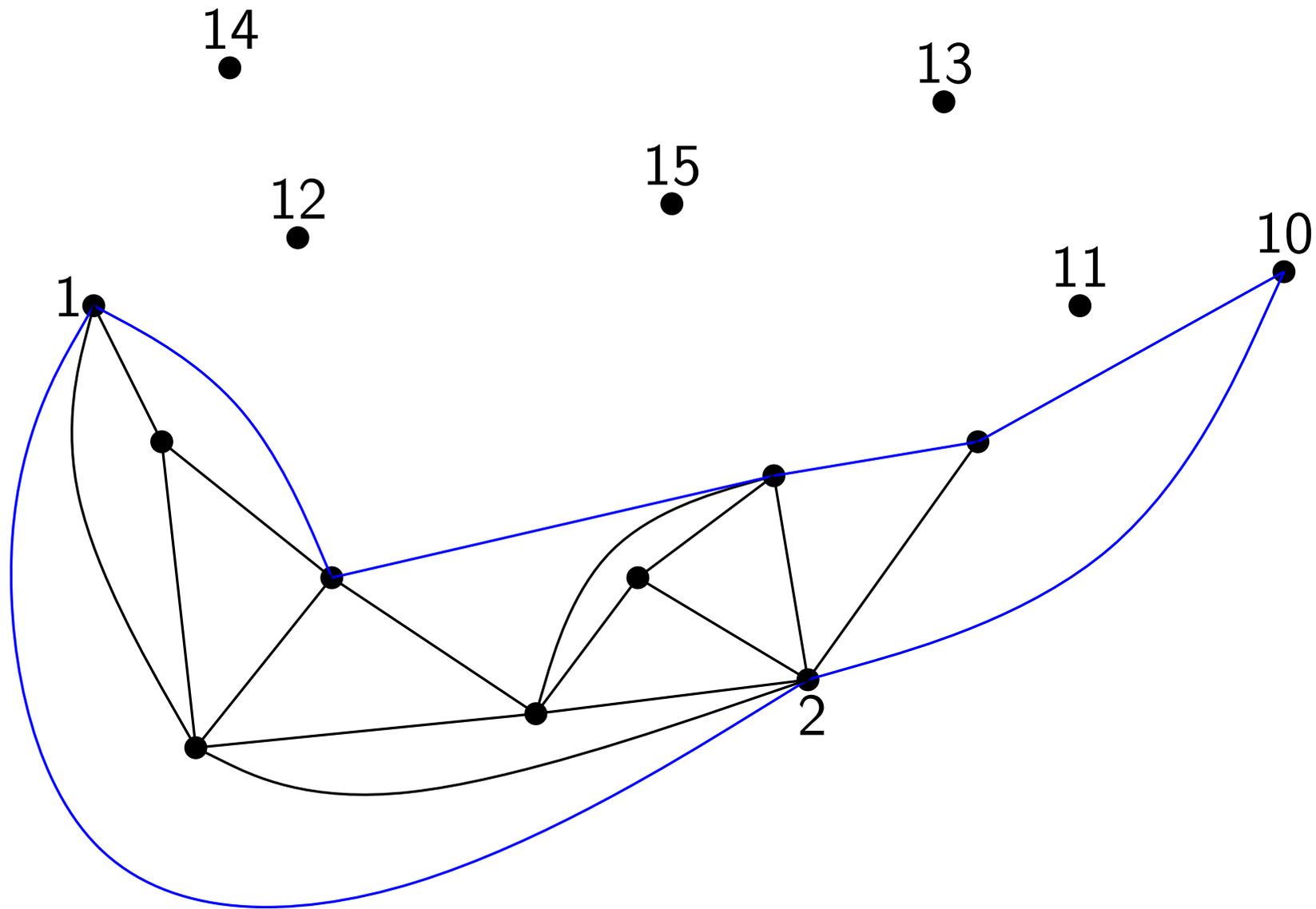
15
●

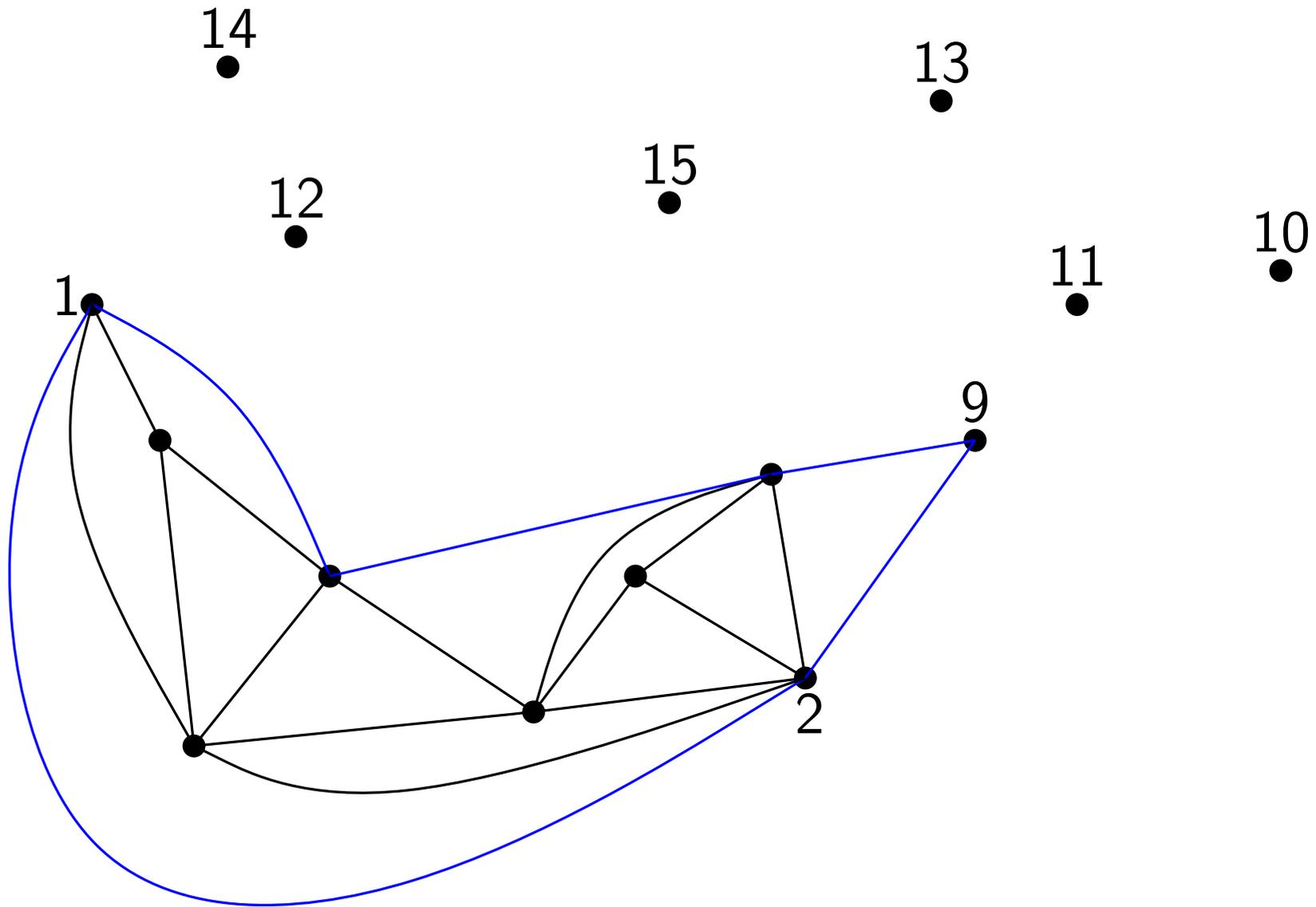


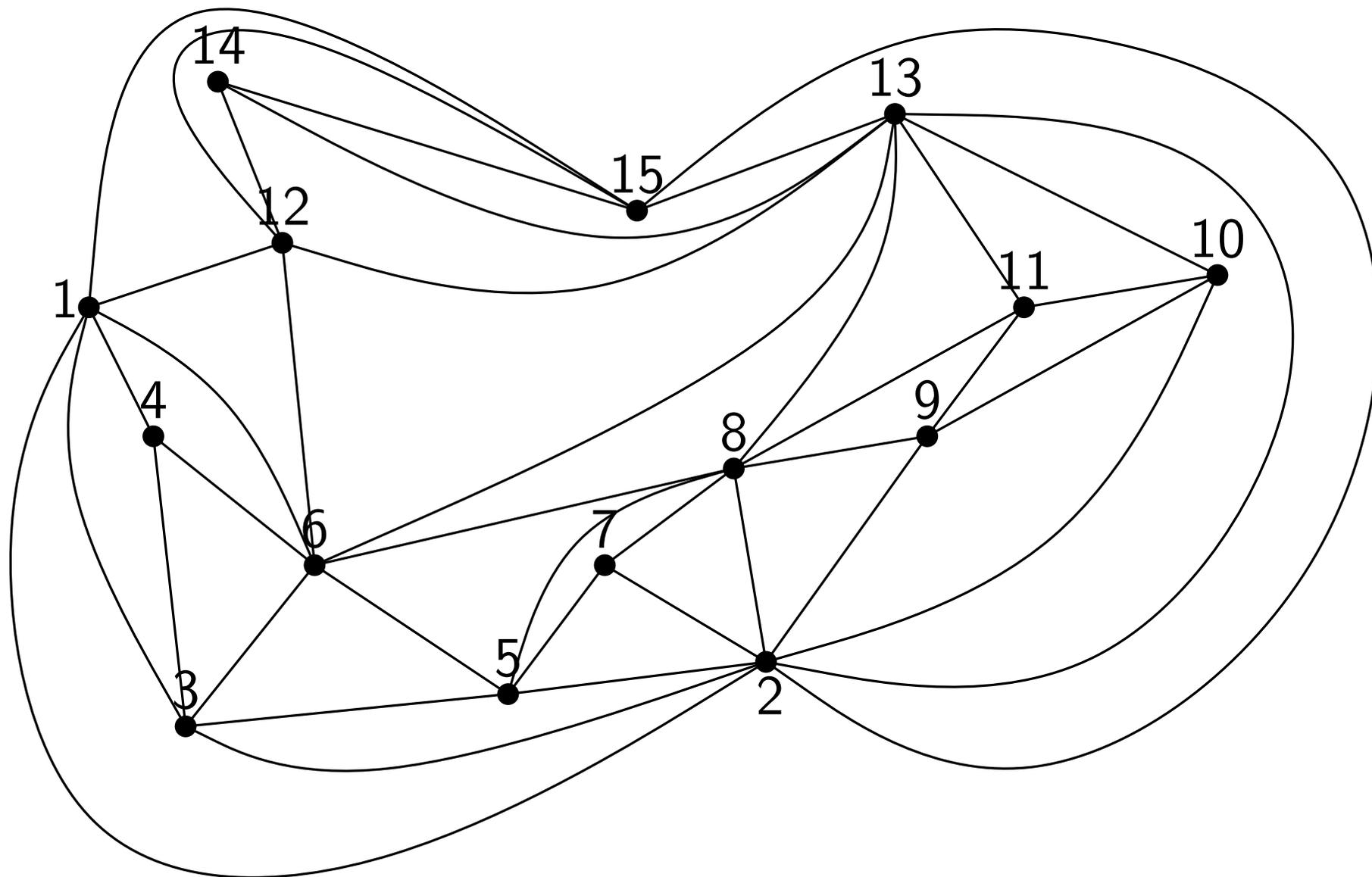


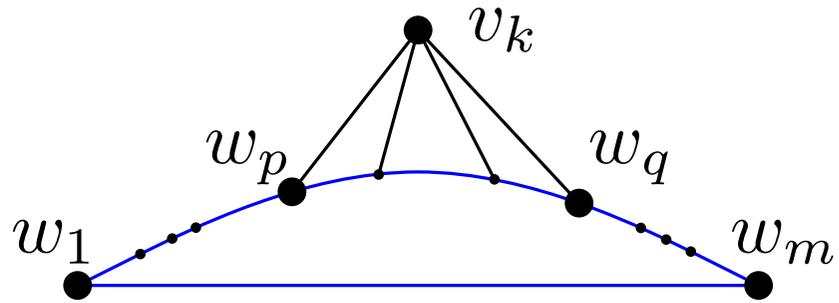








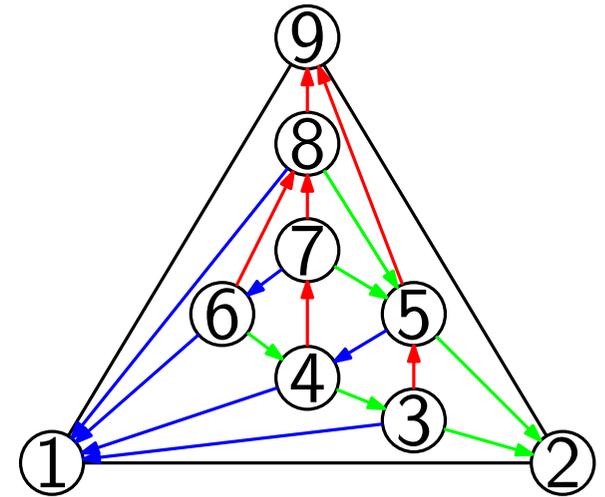


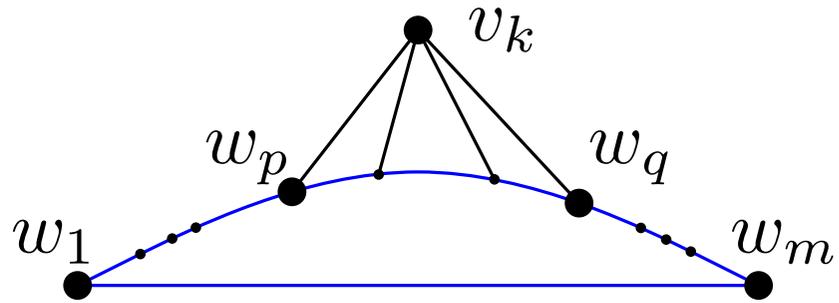


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

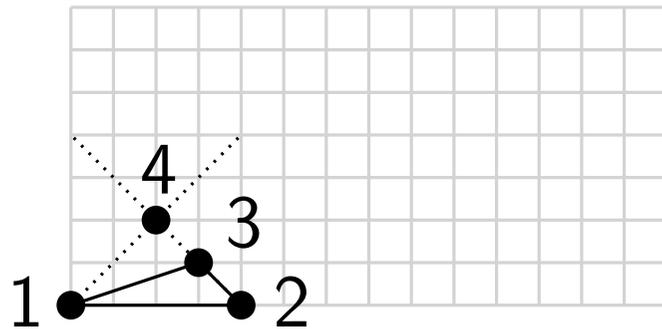
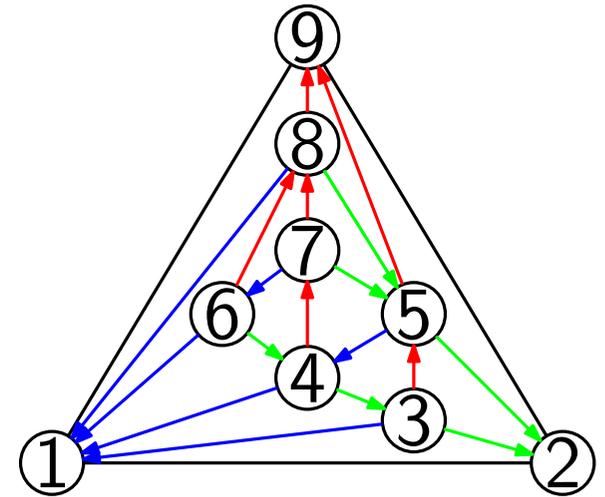


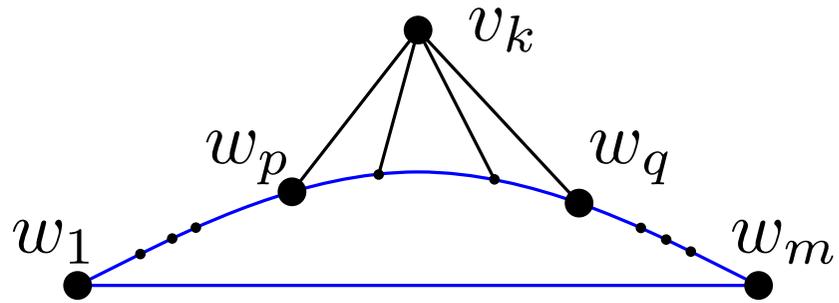


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

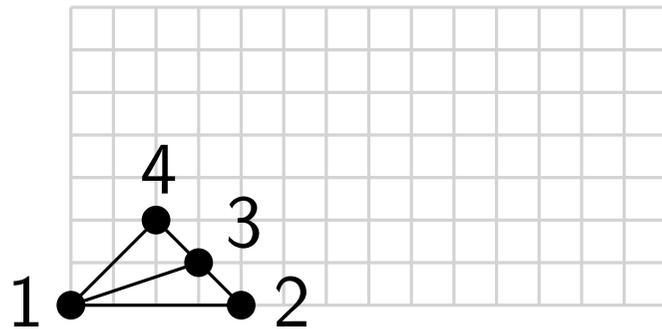
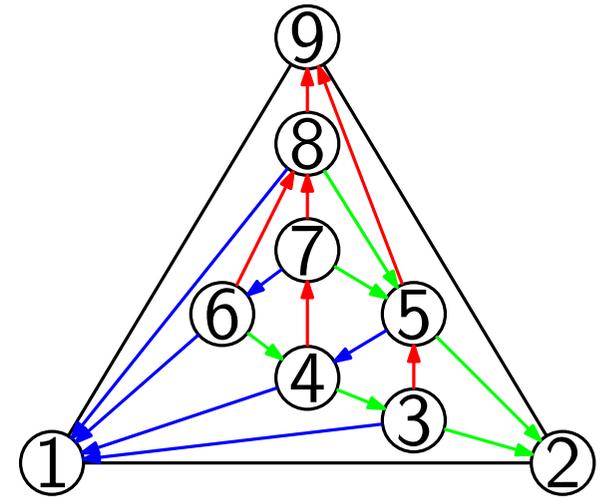


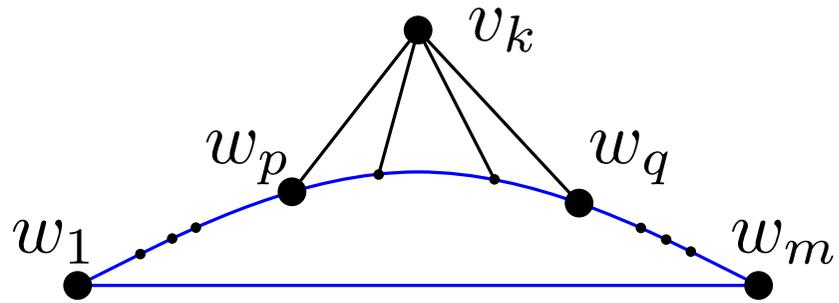


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

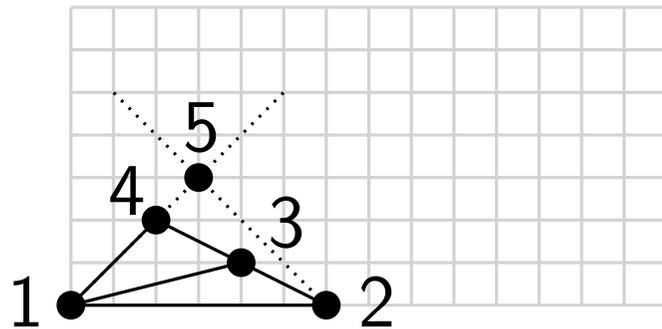
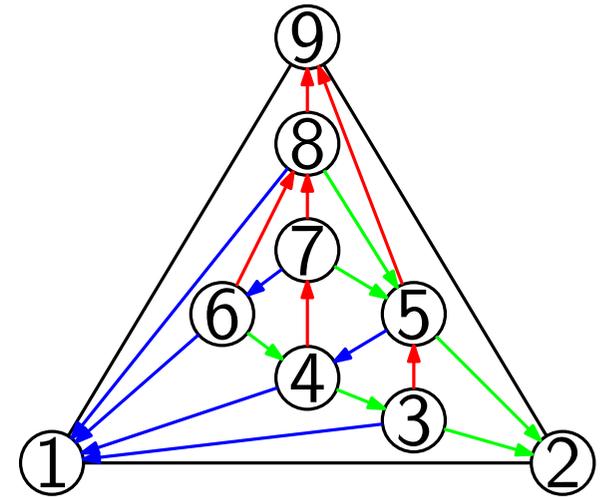


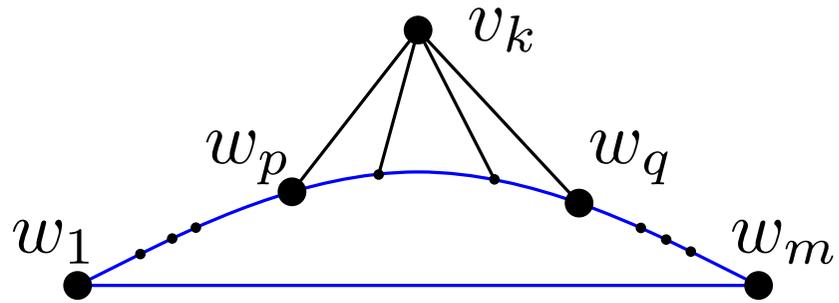


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

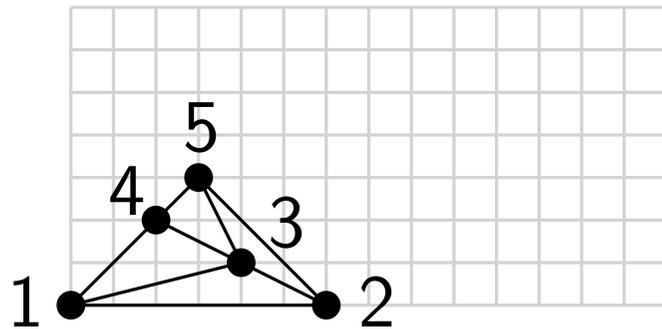
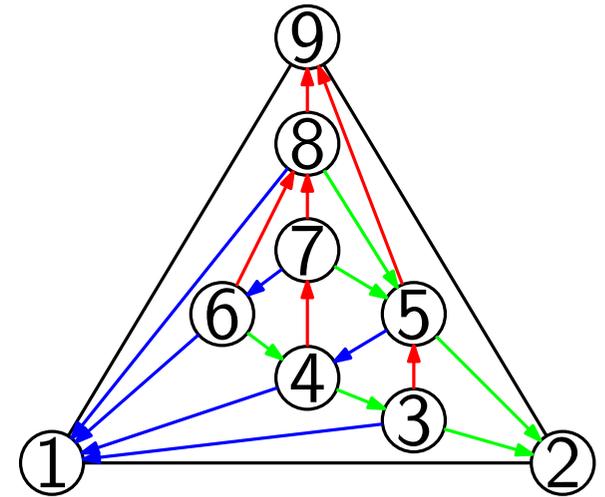


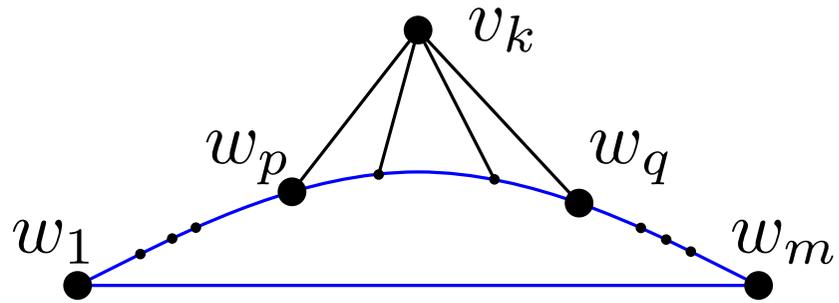


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

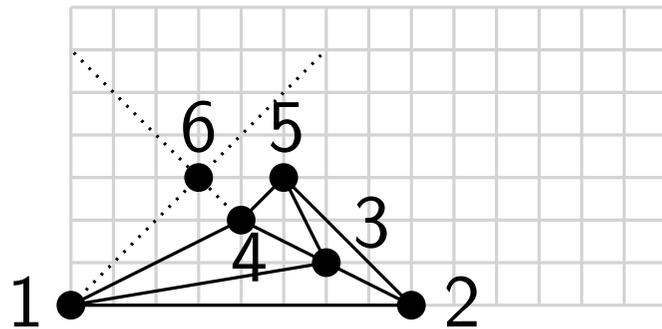
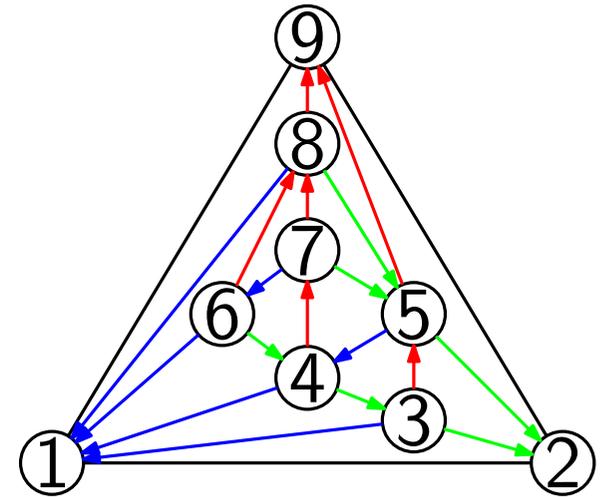


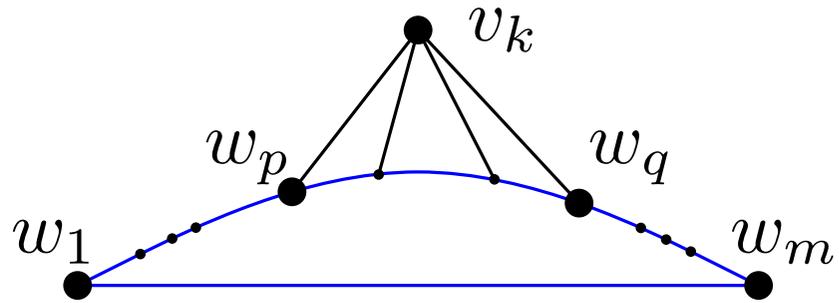


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

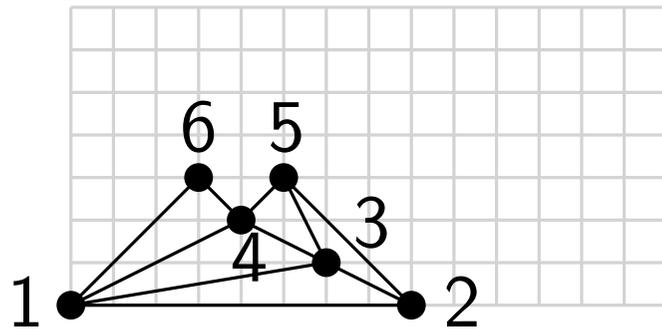
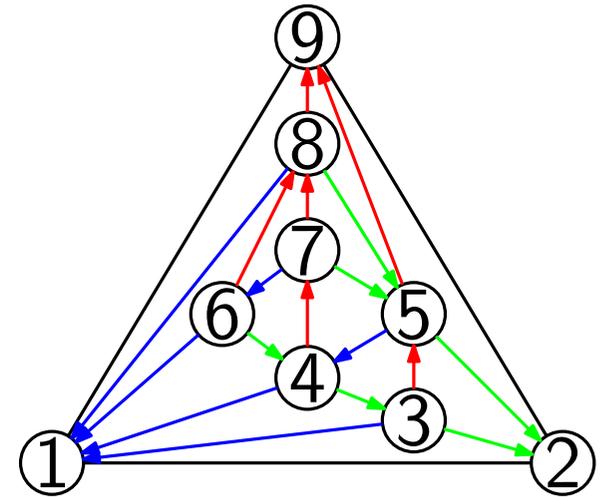


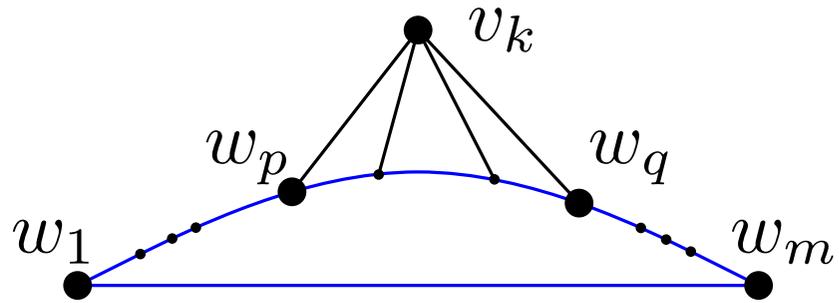


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

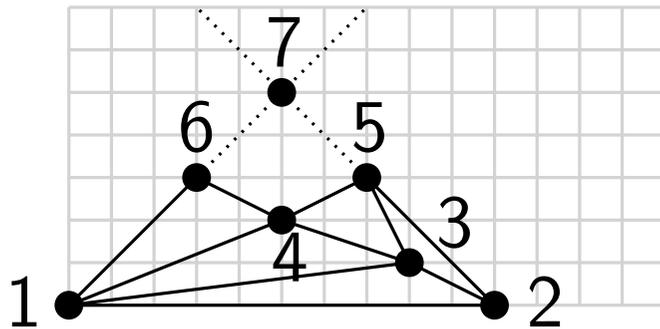
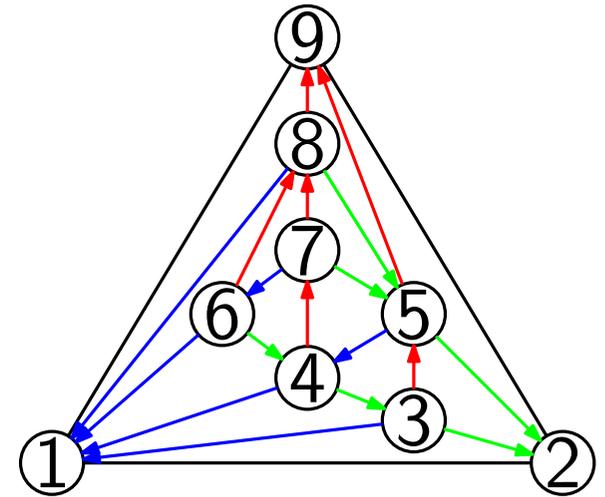


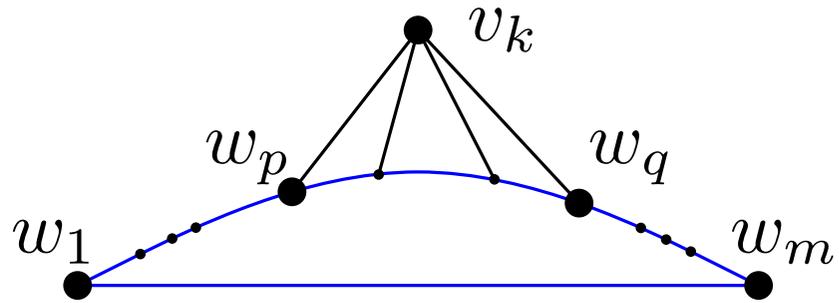


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

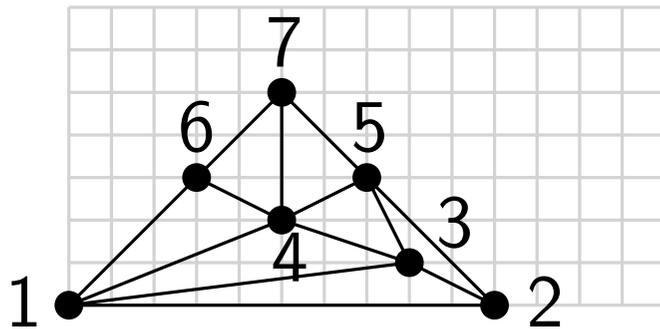
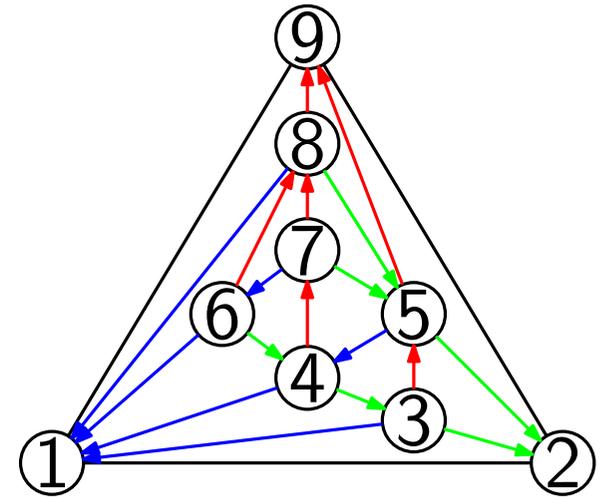


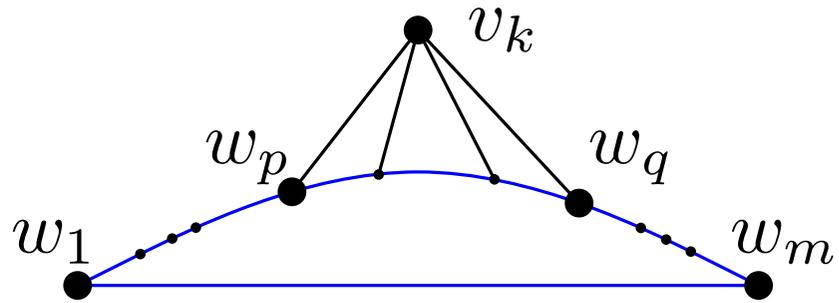


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

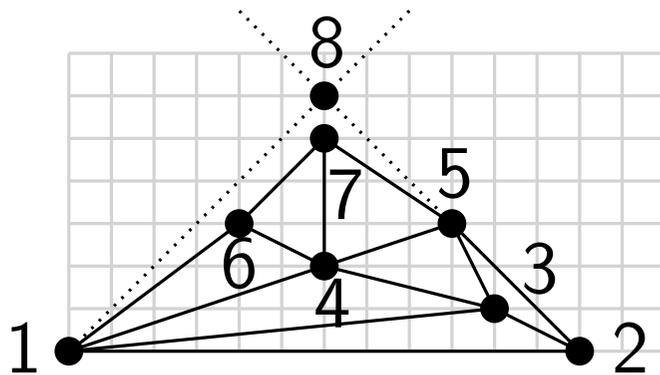
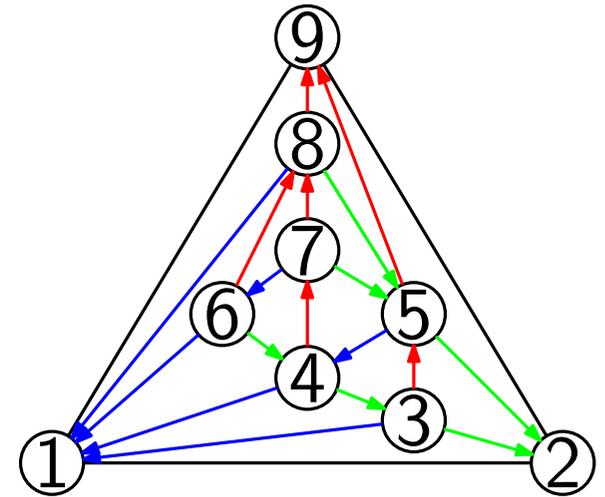


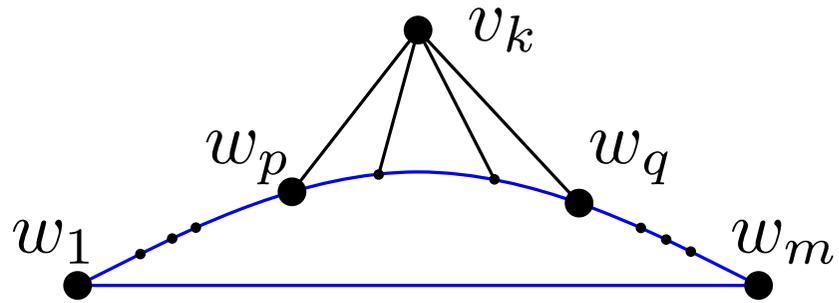


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

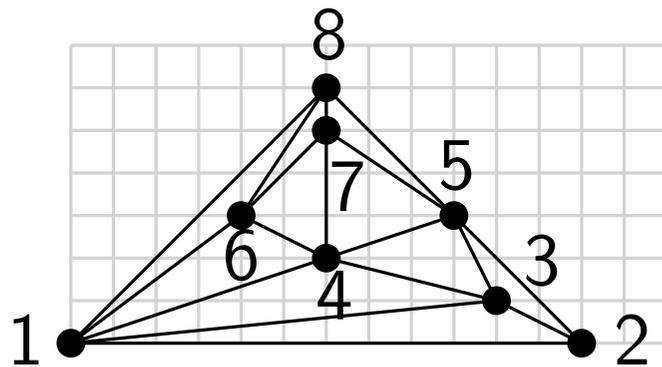
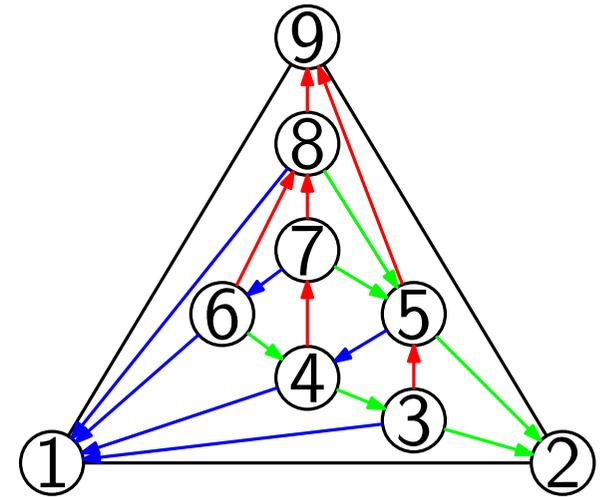


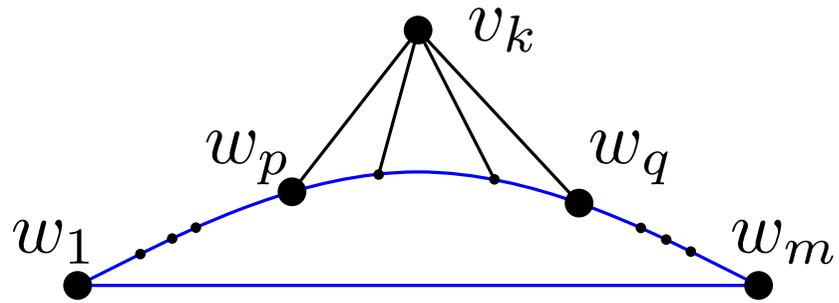


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

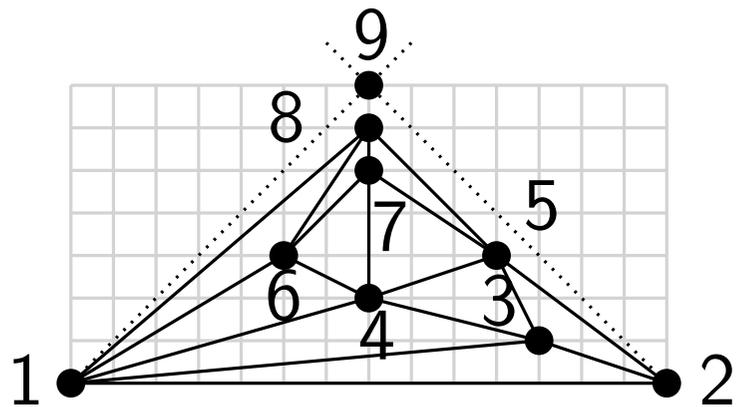
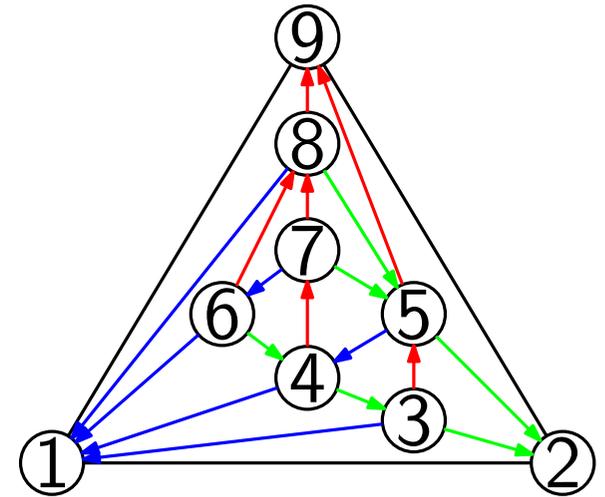


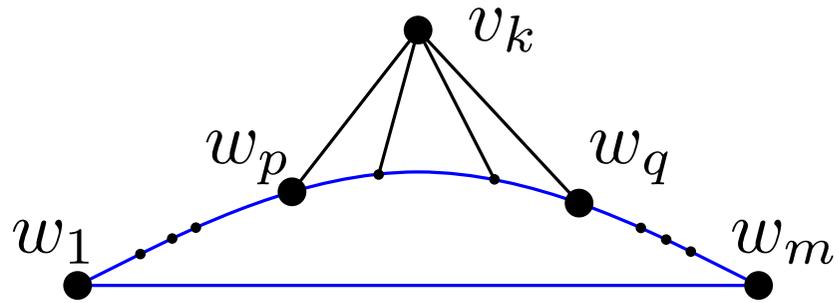


If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$





If $v \in \bigcup_{i=p+1}^{q-1} L(w_i)$ then $x(v) += 1$.

If $v \in \bigcup_{i=q}^m L(w_i)$ then $x(v) += 2$.

$$L(v_k) = \{v_k\} \cup \bigcup_{i=p+1}^{q-1} L(w_i).$$

