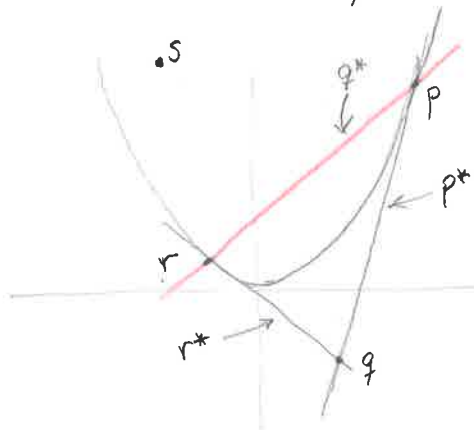


L15

Remember duality from last time...



Dual of point p on parabola?

Dual of point q below parabola?

Incidence property (twice)

$$q \in p^* \Rightarrow p \in q^*$$

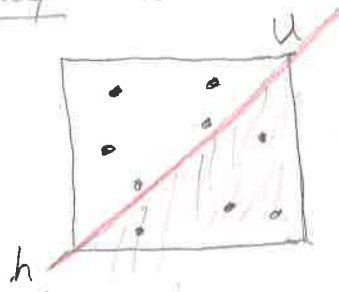
$$q \in r^* \Rightarrow r \in q^*$$

now we know two points on q^*

Dual of point s above parabola?

Discrepancy

[What's a good sample?]



Given a set S of n points in unit square U

what is the maximum (supremum) over all half planes h of

$$|\mu(h) - \mu_S(h)| = |1/2 - 4/9|$$

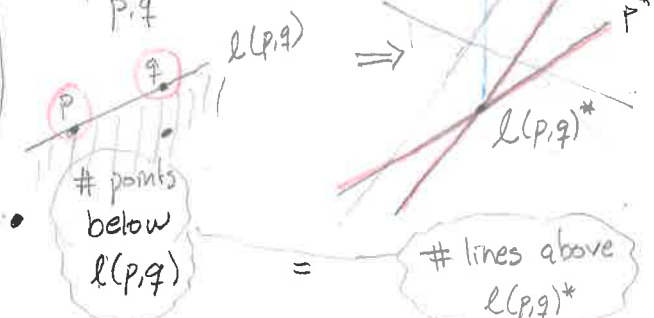
we could shift h slightly to make this difference larger until h touches $p \in S$.

$$|\mu(h) - \mu_S(h)|$$

fraction of U in h fraction of S in h

$$\textcircled{2} \max |\mu(h) - \mu_S(h)|$$

h touching two points of S
 p, q



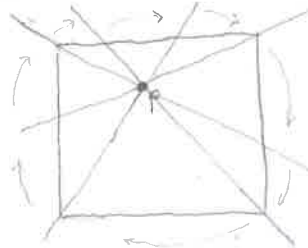
$$\# \text{ points below } l(p,q) = \# \text{ lines above } l(p,q)^*$$

(we also need #points above and on)

Using the level of every vertex in line arrangement we can calculate this

$$\text{level} = \# \text{ lines above a vertex in } A(L)$$

$\textcircled{1}$ If h touches only one point p of S , rotating h about p

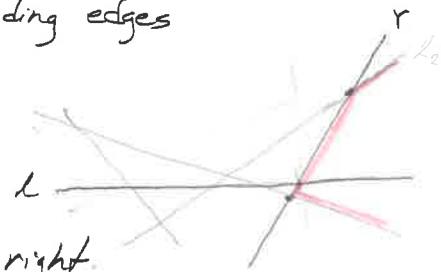


in each wedge, at most two extremes of $\mu(h)$ at most 8 wedges

Find Max in Linear time for all p

We prove that the number of left bounding edges of faces in zone of l is $\leq 5n$.

By induction on n . $n=1$ true



Let $r \in L$ be line intersecting l furthest to right.

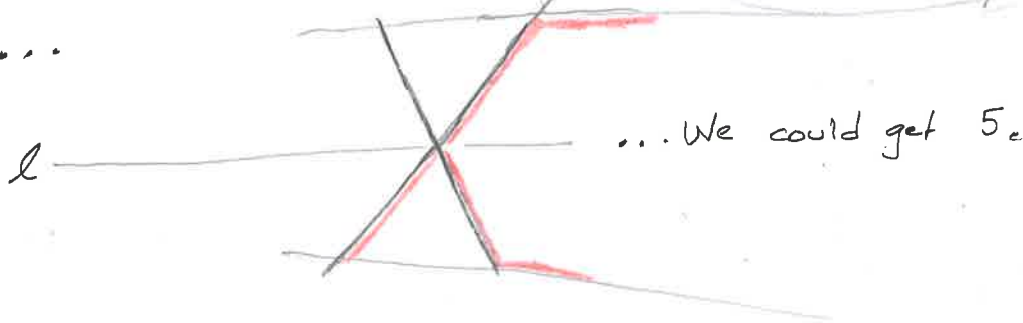
By induction the zone of l in $A(L \setminus \{r\})$ has $\leq 5(n-1)$ left bounding edges.

Adding r creates new left-bounding edges on r (exactly 1) and splits old left-bounding edges (at most 2).

any line intersects one existing cell in at most two edges

any line contributes at most one bounding edge to one cell why?

But ...



$$5(n-1) + 5 \leq 5n \quad \square$$

Incremental construction of arrangement takes

$$O\left(\sum_{i=1}^n i\right) = O(n^2) \text{ time.}$$

Can calculate level of each vertex in $O(n^2)$ time

[Walk along each line from left to right]
[Update level at vertices]

$\Rightarrow O(n^2)$ time discrepancy alg.

